

Elliptic curves, number theory and cryptography

3. handin – The Montgomery ladder

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1. THE MONTGOMERY LADDER IN SAGEMATH ON `CURVE25519`

You will implement the Montgomery ladder for constant-time scalar multiplication on an elliptic curve. The general Montgomery curve, for $B \neq 0$ and $A \neq \pm 2$, is

$$(1) \quad E^M : By^2 = x^3 + Ax^2 + x \text{ over a finite field } \mathbb{F}_q \text{ of characteristic } p \geq 5 .$$

For testing, we need a correspondance with a leading coefficient of y^2 to be 1. Let us neutralise the B coefficient of y^2 by dividing the curve equation by $B^3 \in \mathbb{F}_q$:

$$\begin{aligned} \frac{E^M}{B^3} : \frac{By^2}{B^3} &= \frac{x^3}{B^3} + \frac{Ax^2}{B^3} + \frac{x}{B^3} \\ \iff \left(\frac{y}{B}\right)^2 &= \left(\frac{x}{B}\right)^3 + \frac{A}{B} \left(\frac{x}{B}\right)^2 + \frac{1}{B^2} \frac{x}{B} \\ \iff y'^2 = x'^3 + \frac{A}{B}x'^2 + \frac{1}{B^2}x' &\text{ with } x' = x/B, y' = y/B \end{aligned}$$

therefore there is a \mathbb{F}_q -rational isomorphism between E^M and E'^M

$$E'^M : y'^2 = x'^3 + \frac{A}{B}x'^2 + \frac{1}{B^2}x'$$

$$\begin{aligned} i : E^M &\rightarrow E'^M \\ (x, y) &\mapsto (x/B, y/B) \end{aligned}$$

and the inverse is

$$\begin{aligned} i^{-1} : E'^M &\rightarrow E^M \\ (x', y') &\mapsto (x \cdot B, y \cdot B) \end{aligned}$$

We will use the representation E'^M for tests in SageMath as follows. Let E be the `curve25519` curve:

```
p = ZZ(2**255-19)
Fp = GF(p)
# Montgomery form is y^2 = x^3 + 486662*x^2 + x
A = Fp(486662)
B = Fp(1)
EM = EllipticCurve([0, A/B, 0, 1/B**2, 0])
```

The Montgomery form of elliptic curve is not competitive compared to the short Weierstrass form with the double-and-add algorithm. However Peter L. Montgomery observed that skipping the y -coordinate and using projective X, Z -coordinates, the scalar multiplication becomes competitive.

In this part, you will implement the group law in X, Z -coordinates, then the Montgomery ladder for scalar multiplication. The file `handin3.py` contains the

addition and doubling in affine and projective coordinates on E^M with test functions, to serve as an example. **You are expected to download the file `handin3.py` and write the answers to the questions as SageMath functions in this file. Upload this file with your code for the hand-in.**

Question 1. Implement the x -only addition and doubling in x -only affine coordinates according to the formulas of the 1st hand-in:

Let $P_1 = (x_1, y_1)$ and $P_2 = (x_2, y_2)$ be two points on E^M . Let $P_1 + P_2 = (x_3, y_3)$ and $P_1 - P_2 = (x_4, y_4)$. Assume that $x_1 \neq x_2$, $x_1 \neq 0$ or $x_2 \neq 0$, and $x_4 \neq 0$. From the 1st hand-in one has

$$(2) \quad x_3 x_4 (x_1 - x_2)^2 = (x_1 x_2 - 1)^2$$

and for $P_1 = P_2$, with $x_1 \neq 0$,

$$(3) \quad 4x_1 x_3 (x_1^2 + Ax_1 + 1) = (x_1^2 - 1)^2.$$

The functions whose header are given below are sketched in the PYTHON file of the hand-in, complete these functions in the file:

```
def add_affine_x_only(x1, x2, x4):
def double_affine_x_only(x1, A):
```

Question 2. Test your functions of the previous question. The functions assume that the inputs are not \mathcal{O} nor points of order 2, more precisely: $x_1 \neq x_2$, $x_1 \neq 0$ or $x_2 \neq 0$, and $x_4 \neq 0$.

Question 3. Implement the x -only addition and doubling in X, Z -projective coordinates, based on the affine coordinates. It means to avoid the divisions, you will have two coordinates (X, Z) such that the correspondance with the affine coordinates is $x = X/Z$ for non-zero Z , and if $P(X, Z = 0)$, then P corresponds to the point at infinity \mathcal{O} .

Remember that you can use the Elliptic Curve Formula Database at <http://www.hyperelliptic.org/EFD/> to check your answers.

Use these function names:

```
def add_proj_x_only(X1, Z1, X2, Z2, X4, Z4):
def double_proj_x_only(X1, Z1, A):
```

Question 4. Test your functions of the previous question.

Montgomery's binary scalar multiplication is given in Algorithm 1. Montgomery observed that at each step, the difference $R_0 - R_1$ is always equal to P .

Algorithm 1: Montgomery's binary scalar multiplication

Input: $m = \sum_{i=0}^{n-1} b_i 2^i$ with $b_{n-1} = 1$, and point $P \in E^M$ in affine coordinates

Output: $[m]P$

```
1  $(R_0, R_1) \leftarrow (P, [2]P)$ 
2 for  $i = n - 2$  down to 0 do
3   if  $b_i = 0$  then
4      $(R_0, R_1) \leftarrow ([2]R_0, R_0 + R_1)$ 
5   else
6      $(R_0, R_1) \leftarrow (R_0 + R_1, [2]R_1)$ 
7 return  $R_0$ 
```

This gives the Montgomery ladder in Algorithm 2

Algorithm 2: Montgomery's ladder for scalar multiplication

Input: $m = \sum_{i=0}^{n-1} b_i 2^i$ with $b_{n-1} = 1$, and x_P affine x -coordinate of point $P \in E^M$

Output: x -coordinate of $[m]P$

```

1  $(x_0, x_1) \leftarrow (x_P, \text{double\_affine\_x\_only}(x_P, A))$ 
2 for  $i = n - 2$  down to 0 do
3   if  $b_i = 0$  then
4      $(x_0, x_1) \leftarrow$ 
        $(\text{double\_affine\_x\_only}(x_0), \text{add\_affine\_x\_only}(x_0, x_1, x_P))$ 
5   else
6      $(x_0, x_1) \leftarrow$ 
        $(\text{add\_affine\_x\_only}(x_0, x_1, x_P), \text{double\_affine\_x\_only}(x_1, A))$ 
7 return  $R_0$ 
```

Question 5. Implement the Montgomery ladder, using the two affine functions of Question 1.

Question 6. Test your function of the previous question.

Question 7. Implement the Montgomery ladder, using the two projective functions of Question 3.

Question 8. Test your function of the previous question.

2. THE MONTGOMERY LADDER IN SAGEMATH IN CHARACTERISTIC 2

Now you will implement the Montgomery Ladder in characteristic 2 using x -only homogeneous projective coordinates (X, Z) in Weierstrass form.

Define the base field (actually an extension of F2)

`m = 233`

`R = PolynomialRing(GF(2), 'Z')`

`Z = R.gen()`

`F2m = GF(2**m, 'z', modulus = Z**233 + Z**74 + 1)`

Define curve NIST B-233 $y^2 + xy = x^3 + a_2x^2 + a_6$

`a_2 = F2m(1)`

`a_6 = F2m.fetch_int(0x0066647ede6c332c7f8c0923bb58213b333b20e9ce4281fe115f7d8f90ad)`

`E2 = EllipticCurve(F2m, [1, a_2, 0, 0, a_6])`

Question 9. Let $P_1 = (x_1, y_1)$ and $P_2 = (x_2, y_2)$ be two points not of order 2 on E^2 , with $P_1 \neq \pm P_2$. Let $P_1 + P_2 = (x_3, y_3)$ and $P_1 - P_2 = (x_4, y_4)$.

Start by implementing the x -only point addition and doubling in x -only affine coordinates according to the formulas below:

$$(4) \quad x_3 = x_4 + \frac{x_2}{x_1 + x_2} + \left(\frac{x_2}{x_1 + x_2} \right)^2$$

For the case where $P_1 = P_2$ and $x_1 \neq 0$, then $x_3 = x_1^2 + \frac{a_6}{x_1^2}$. You can complete the functions in the file:

`def add_affine_x_only_char2(x1, x2, x4):`

`def double_affine_x_only_char2(x1, a_6):`

Question 10. Now follow the same procedure above to finish the Montgomery Ladder implementation in affine and projective coordinates.