

Logics

I. Decide whether the following formula is valid or not (justify your answer).

$$(((a \rightarrow b) \rightarrow a) \wedge (a \rightarrow c)) \rightarrow c$$

II. Prove that the following formula is derivable.

$$\neg(\alpha \vee \beta) \rightarrow (\neg\alpha \wedge \neg\beta)$$

III. Consider a first-order language without functional symbol, and whose only relational symbol is R . Then, consider the formula ϕ , which is defined as follows:

$$\phi = \phi_0 \rightarrow (\phi_1 \rightarrow \phi_2)$$

where:

$$\phi_0 = \forall x. (\exists y. R(x, y))$$

$$\phi_1 = \forall x. (\forall y. (\forall z. R(x, y) \rightarrow (R(y, z) \rightarrow R(x, z))))$$

$$\phi_2 = \exists x. R(x, x)$$

Define a model $\mathcal{M} = \langle D, I \rangle$ as follows:

$$D = \mathbb{N}$$

$$I(R) = \{(m, n) \in \mathbb{N} \times \mathbb{N} \mid m < n\}$$

Show that \mathcal{M} is a counter-model of ϕ .

IV. Let \mathcal{F} be a ranked alphabet of function symbols, and let \mathcal{X} be an alphabet of variables. The set of terms \mathcal{T} , built over \mathcal{F} and \mathcal{X} , is inductively defined as follows:

1. If $x \in \mathcal{X}$, then $x \in \mathcal{T}$;
2. If $c \in \mathcal{F}$ is of arity 0, then $c \in \mathcal{T}$;
3. If $f \in \mathcal{F}$ is of arity n , and $t_1, \dots, t_n \in \mathcal{T}$, then $f(t_1, \dots, t_n) \in \mathcal{T}$.

Let $\mathcal{M} = \langle D, I \rangle$ be a model, and $\rho : D^{\mathcal{X}}$ be a valuation. The interpretation of a term t is inductively defined as follows:

1. $\llbracket x \rrbracket \rho = \rho(x)$, for $x \in \mathcal{X}$;
2. $\llbracket c \rrbracket \rho = I(c)$, for $c \in \mathcal{F}$ of arity 0;
3. $\llbracket f(t_1, \dots, t_n) \rrbracket \rho = I(f)(\llbracket t_1 \rrbracket \rho, \dots, \llbracket t_n \rrbracket \rho)$, for $f \in \mathcal{F}$ of arity n .

Let $\rho \in D^{\mathcal{X}}$ be a valuation, and let $a \in D$ and $x \in \mathcal{X}$. The valuation $\rho[x:=a]$ is defined as follows:

$$\rho[x:=a](y) = \begin{cases} a & \text{if } x = y \\ \rho(y) & \text{if } x \neq y \end{cases}$$

Finally, let t and u be terms, and x be a variable. The substitution of x by u in t , in notation $t[u/x]$, is inductively defined as follows:

1. $x[u/x] = u$;
2. $y[u/x] = y$, for $y \in \mathcal{X}$ and $y \neq x$;
3. $c[u/x] = c$, for $c \in \mathcal{F}$ of arity 0;
4. $f(t_1, \dots, t_n)[u/x] = f(t_1[u/x], \dots, t_n[u/x])$, for $f \in \mathcal{F}$ of arity n .

Let x be a variable, u be a term, and ρ be a valuation. Show that for every term t :

$$\llbracket t[u/x] \rrbracket \rho = \llbracket t \rrbracket \rho[x:=\llbracket u \rrbracket \rho]$$

Hint: proceed by induction on the structure of t .