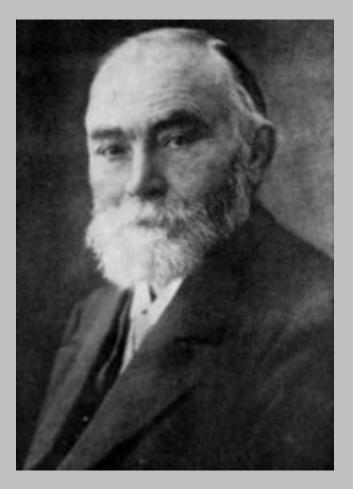
Formal Semantics of Natural Language

Philippe de Groote and Yoad Winter

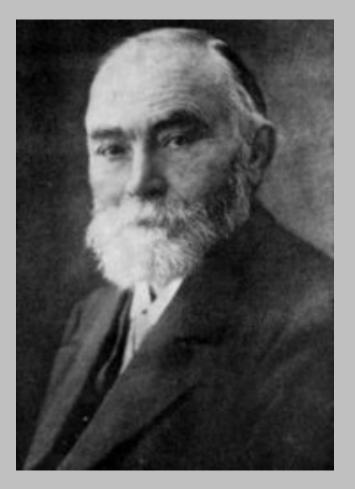
ESSLLI 2021, Online

Additional Topics:

Intensionality

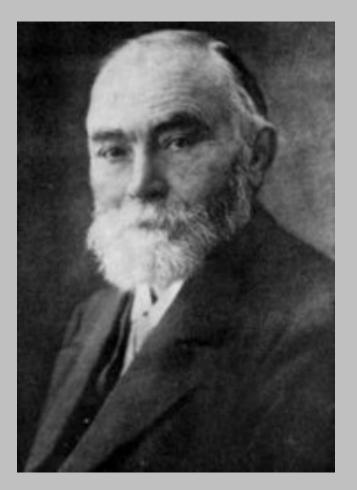


Gottlob Frege (1848-1925)



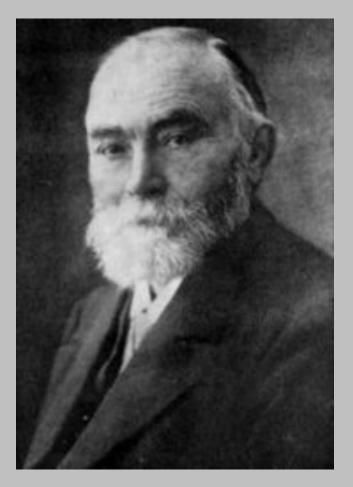
Sinn (sense)/Bedeutung (reference)
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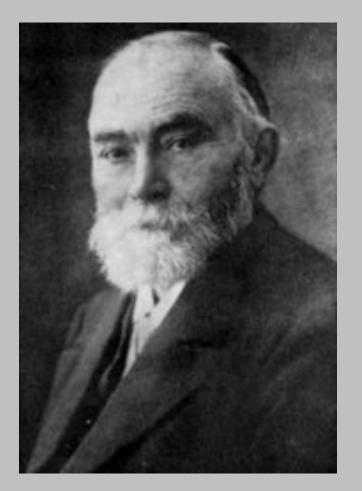
- Sinn (sense)/Bedeutung (reference) — Frege
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- According to Frege, the sense of an expression is its "mode of presentation", while the reference or denotation of an expression is the object it refers to.
- For instance, both expressions "1 + 1" and "2" have the same denotation but not the same sense.

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- Frege gives the example of *"the morning star"* and *"the evening star"* which both refer to the planet Venus.
- Compare "the morning star is the evening star" with "the ancients did not know that the morning star is the evening star".



G.W. von Leibniz (1646–1716)



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Pangloss enseignait la métaphysicothéologo-cosmolo-nigologie. Il prouvait admirablement qu'il n'y a point d'effet sans cause, et que, dans ce meilleur des mondes possibles, le château de monseigneur le baron

était le plus beau des châteaux et madame la meilleure des baronnes possibles.

Voltaire (Candide)

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 - ▶ Epistemic logic: *Bob knows that...* Bob ignores that...
 - Temporal logic: It will always be the case that... It will eventually be the case that...

Modal logic

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Syntax:

$F ::= a \mid \neg F \mid F \lor F \mid \Box F$

Define the other connectives in the usual way. Define $\Diamond A$ as $\neg \Box \neg A$.

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Validity:

let $\mathcal{M} = \langle W, P \rangle$, where W is a set of "possible worlds", and P is a function that asigns to each atomic proposition a subset of W.

 $\blacktriangleright \mathcal{M}, s \models a \text{ iff } s \in P(a).$

$$\blacktriangleright \mathcal{M}, s \models \neg A \text{ iff not } \mathcal{M}, s \models A.$$

▶ $\mathcal{M}, s \models A \lor B$ iff either $\mathcal{M}, s \models A$ or $\mathcal{M}, s \models B$, or both.

 $\blacktriangleright \mathcal{M}, s \models \Box A \text{ iff for every } t \in W, \ \mathcal{M}, t \models A.$

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necessarily := $\lambda A w. \forall v. (A v)$

This red car is a Ferrari

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This skillful surgeon is Dr Johnson

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This skillful surgeon is Dr Johnson

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Solution:

surgeon : e (st)driver : e (st)skillful : (e (st)) e (st)