SIAM AG 2023 · Elliptic curves and pairings in cryptography

Revisiting cycles of pairing-friendly elliptic curves

Marta Bellés

marta@dusk.network

Joint work with Jorge Jimenez-Urroz and Javier Silva







1. THE PROBLEM

THE 2-CYCLE PROBLEM



THE 2-CYCLE PROBLEM

Find two elliptic curves E/\mathbb{F}_q and E'/\mathbb{F}_p such that

 $|E(\mathbb{F}_q)| = p \text{ and } |E'(\mathbb{F}_p)| = q.$

EASY

Find two elliptic curves E/\mathbb{F}_q and E'/\mathbb{F}_p such that $|E(\mathbb{F}_q)| = p \text{ and } |E'(\mathbb{F}_p)| = q,$ with low embedding degree (pairing-friendly).

Find two elliptic curves E/\mathbb{F}_q and E'/\mathbb{F}_p such that $|E(\mathbb{F}_q)| = p \text{ and } |E'(\mathbb{F}_p)| = q,$ with low embedding degree (pairing-friendly).

Embedding degree

- E: smallest k such that $p | q^k 1$.
- E': smallest l such that $q | p^l 1$.

Pairing

• *Small* embedding degree: DL attacks • *Large* embedding degree: inefficiency

Find two elliptic curves E/\mathbb{F}_q and E'/\mathbb{F}_p such that $|E(\mathbb{F}_q)| = p \text{ and } |E'(\mathbb{F}_p)| = q,$

with low embedding degree (pairing-friendly).

DIFFICULT

2. MOTIVATION





data **x**, operation **F**

$$\mathbf{y} = \mathbf{F}(\mathbf{x})$$

powerful computer









data **x**, operation **F**



how do we know y is really the result??







F(**x**) = **?**



• fast verification



Verifier



circuit (arithmetic modulo p)

A **proof** asserts that a set of **inputs** and **outputs** satisfy the relations defined in the **circuit**.



PAIRING-BASED SNARKS PROOF SYSTEMS

F_p-arithmetic circuit



Prover





Verifier

RECURSIVE PROOF COMPOSITION



By verifying one single proof, we can verify that all computations (and proofs) are correct.

PAIRING-BASED SNARKS PROOF SYSTEMS





Verifier
<u>Pairing</u> operation on points of E (arithmetic over Fq).











- Find two elliptic curves E/\mathbb{F}_q and E'/\mathbb{F}_p such that
 - $|E(\mathbb{F}_q)| = p \text{ and } |E'(\mathbb{F}_p)| = q,$
 - with low embedding degree (pairing-friendly).

3. WHAT WAS KNOWN



CONDITIONS

- Curves involved in a cycle must be of prime order.
- The only known method to produce prime-order curves is via **families of curves parameterized by polynomials** q(X), p(X), and t(X) with embedding degree k and discriminant d. This means:

1.
$$\mathbf{p}(X) = \mathbf{q}(X) + 1 - \mathbf{t}(X)$$
.

- 3. p(X) and q(X) represent primes.
- 4. **p**(X) $| \Phi_k(\mathbf{t}(X) 1).$
- 5. The equation $4\mathbf{q}(\mathbf{X}) = \mathbf{t}(\mathbf{X})^2 + |\mathbf{d}|\mathbf{Y}^2$ has infinitely many integer solutions (x, y).

 $\stackrel{(4)}{\longrightarrow}$ the embedding degree is at most *k*

- (1-3) infinitely many parameters compatible with elliptic curves
- (5) infinitely many curves in the family with same discriminant

FAMILIES OF PAIRING-FRIENDLY CURVES OF PRIME ORDER

• There are **no elliptic curves with prime order** and embedding degree **k < 3**.

For k = 3, 4, 6 we have the families of curves	MNT3
Miyaji-Nakabayashi-Takano (MNT).	MNT4
(exhaustive)	MNT6
	For k = 3, 4, 6 we have the families of curves Miyaji-Nakabayashi-Takano (MNT). (<i>exhaustive</i>)

- For **k = 10** we have the **Freeman** family of Freeman curves.
- BN • For k = 12 we have the Barreto-Naehrig (**BN**) family of curves.

- $\mathbf{p}(X) = 12X^2 6X + 1$ $\mathbf{q}(X) = 12X^2 1$
- $\mathbf{p}(X) = X^2 + 2X + 2$ $\mathbf{q}(X) = X^2 + X + 1$
- $p(X) = 4X^2 2X + 1$ $q(X) = 4X^2 + 1$
- $\mathbf{p}(\mathbf{X}) = 25\mathbf{X}^4 + 25\mathbf{X}^3 + 15\mathbf{X}^2 + 5\mathbf{X} + 1$ $q(X) = 25X^4 + 25X^3 + 25X^2 + 10X + 3$
- $\mathbf{p}(\mathbf{X}) = 36\mathbf{X}^4 + 36\mathbf{X}^3 + 24\mathbf{X}^2 + 6\mathbf{X} + 1$ $q(X) = 36X^4 + 36X^3 + 18X^2 + 6X + 1$

DO THEY FORM CYCLES?

- MNT4 and MNT6 curves **do** form cycles.
 - <u>But:</u> Low embedding degree -> large parameters. Unbalanced embedding degrees.
- Freeman and BN curves **do not** form cycles with curves from their own family.

• Can they form cycles with other curves?







4. MAIN CONTRIBUTION



THEOREM

Consider a family of elliptic curves with embedding degree k parameterized by polynomials p(X), q(X). Let *l* be a natural number. Then either:

- $q(X) | p(X)^{l} 1$, or
- there are at most finitely many 2-cycles formed by a curve form the family and a curve with embedding degree *l*.

In particular, we did an exhaustive search for the known families of curves.

COROLLARY

Except for the few cases described in the table below, we have that:

- An MNT3 curve cannot form 2-cycles with a curve of embedding degree *l* < 23.
- A **Freeman** curve **cannot form 2-cycles** with a curve of embedding degree *l* < 26.
- A **BN** curve **cannot form 2-cycles** with a curve of embedding degree *l* < 33.

<u>Exceptions</u>		k	1	q	p
	MNT3	3	10	11	19
	MNT3	3	10	11	7
	BN	12	18	19	13



5. FUTURE WORK



FUTURE WORK

- Improve our bounds (code) to all k < 56.
- Generalize our result to s-cycles with s > 2.
- Do there exist cycles consisting of elliptic curves with the same embedding degree? It is already known that this is not the case for k = 4, 6, 8, 12.

You can find more open problems in: A. Chiesa, L. Chua, M. Weidner, On cycles of pairing-friendly elliptic curves, arXiv: 1803.02067.

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