# Revisiting cycles of pairing-friendly elliptic curves 

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Joint work with Jorge Jimenez-Urroz and Javier Silva


## 1. THE PROBLEM

## THE 2-CYCLE PROBLEM

Find two (ordinary) elliptic curves $E / \mathbb{F}_{q}$ and $E^{\prime} / \mathbb{F}_{p}$ such that


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EASY

## THE PAIRING-FRIENDLY 2-CYCLE PROBLEM

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with low embedding degree (pairing-friendly).

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$$

with low embedding degree (pairing-friendly).

## Embedding degree

- E: smallest $k$ such that $p \mid q k-1$.
- $\mathrm{E}^{\prime}$ : smallest $l$ such that $\mathrm{q} \mid \mathrm{p} l$ - .


## Pairing

- Small embedding degree: DL attacks
- Large embedding degree: inefficiency


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with low embedding degree (pairing-friendly).

## DIFFICULT

## 2. MOTIVATION

## VERIFIABLE COMPUTATION



## VERIFIABLE COMPUTATION


powerful computer
data $\mathbf{x}$, operation $\mathbf{F}$

$F(x)=?$
small computer

## VERIFIABLE COMPUTATION


powerful computer
data $\mathbf{x}$, operation $\mathbf{F}$

$$
y=F(x)
$$

how do we know
$y$ is really the result??

small computer

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## PAIRING-BASED SNARKS PROOF SYSTEMS



Prover


- proof is small
- fast verification


Verifier

## $\mathbb{F}_{\mathrm{p}}$-ARITHMETIC CIRCUIT SATISFIABILITY



A proof asserts that a set of inputs and outputs satisfy the relations defined in the circuit.

## PAIRING-BASED SNARKS PROOF SYSTEMS

$\mathbb{F}_{\mathrm{p}}$-arithmetic circuit



## RECURSIVE PROOF COMPOSITION



By verifying one single proof, we can verify that all computations (and proofs) are correct.

## PAIRING-BASED SNARKS PROOF SYSTEMS



Prover
Elliptic curve $E / \mathbb{F}_{q}$ such that $\left|E\left(\mathbb{F}_{q}\right)\right|=p$.


Pairs $(x, y)$ in $\mathbb{E}_{q}^{2}$.


Verifier

It needs to be Pairing operation on points of E
efficient! (arithmetic over $\mathbb{F}_{q}$ ).
(1)A SNARK instantiated with $E / \mathbb{F}_{q}$ such that $\left|E\left(\mathbb{F}_{q}\right)\right|=p$.

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A SNARK instantiated with $E^{\prime} / \mathbb{E}$ such that $\left|E^{\prime}\left(\mathbb{F}_{\mathrm{r}}\right)\right|=\mathrm{q}$.

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A SNARK instantiated with $E^{\prime} / \mathbb{E}$ such that $\left|E^{\prime}\left(\mathbb{F}_{\mathrm{p}}\right)\right|=q$.

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with low embedding degree (pairing-friendly).

## 3. WHAT WAS KNOWN

## CONDITIONS

- Curves involved in a cycle must be of prime order.
- The only known method to produce prime-order curves is via families of curves parameterized by polynomials $q(X), p(X)$, and $t(X)$ with embedding degree $k$ and discriminant $d$. This means:

1. $p(X)=q(X)+1-t(X)$.
2. $p(X)$ is integer-valued.
3. $p(X)$ and $q(X)$ represent primes.
4. $p(X) \mid \Phi_{k}(t(X)-1)$.
5. The equation $4 q(X)=t(X)^{2}+|d| Y^{2}$ has infinitely many integer solutions $(x, y)$.
$\xrightarrow{(1-3)}$ infinitely many parameters compatible with elliptic curves
$\xrightarrow{(4)}$ the embedding degree is at most $k$
$\xrightarrow{(5)}$ infinitely many curves in the family with same discriminant

## FAMILIES OF PAIRING-FRIENDLY CURVES OF PRIME ORDER

- There are no elliptic curves with prime order and embedding degree $\mathbf{k}<3$.
- For $\mathbf{k}=3,4,6$ we have the families of curves Miyaji-Nakabayashi-Takano (MNT).
(exhaustive)
- For $\mathbf{k}=10$ we have the Freeman family of curves.
- For $\mathbf{k}=\mathbf{1 2}$ we have the Barreto-Naehrig (BN) family of curves.


## DO THEY FORM CYCLES?

- MNT4 and MNT6 curves do form cycles.

But: Low embedding degree -> large parameters. Unbalanced embedding degrees.

- Freeman and BN curves do not form cycles with curves from their own family.

- Can they form cycles with other curves?



## 4. MAIN CONTRIBUTION

## THEOREM

Consider a family of elliptic curves with embedding degree $k$ parameterized by polynomials $p(X), q(X)$. Let $l$ be a natural number. Then either:

- $q(X) \mid p(X)^{l}-1$, or
- there are at most finitely many 2-cycles formed by a curve form the family and a curve with embedding degree $l$.


In particular, we did an exhaustive search for the known families of curves.

## COROLLARY

Except for the few cases described in the table below, we have that:

- An MNT3 curve cannot form 2-cycles with a curve of embedding degree $l<23$.
- A Freeman curve cannot form 2-cycles with a curve of embedding degree $l<26$.
- A BN curve cannot form 2-cycles with a curve of embedding degree $l<33$.

Exceptions

|  | $k$ | $l$ | $q$ | $p$ |
| :--- | :---: | :---: | :---: | :---: |
| MNT3 | 3 | 10 | 11 | 19 |
| MNT3 | 3 | 10 | 11 | 7 |
| BN | 12 | 18 | 19 | 13 |

## 5. FUTURE WORK

## FUTURE WORK

- Improve our bounds (code) to all $k<56$.
- Generalize our result to $s$-cycles with $s>2$.
- Do there exist cycles consisting of elliptic curves with the same embedding degree? It is already known that this is not the case for $k=4,6,8,12$.

You can find more open problems in:
A. Chiesa, L. Chua, M. Weidner, On cycles of pairing-friendly elliptic curves, arXiv: 1803.02067.

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