

Co-factor clearing and subgroup membership testing in pairing groups

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September 14, 2022

<https://members.loria.fr/AGuillevic/files/talks/22-Nancy-JAV.pdf>

Outline



Introduction: GLV on elliptic curves

Subgroup membership testing with GLV on \mathbf{G}_1

Faster co-factor clearing

Ensuring correct subgroup membership testing in \mathbf{G}_2 and \mathbf{G}_T

References

-  Youssef El Housni and Aurore Guillevic.
Families of SNARK-friendly 2-chains of elliptic curves.
In Orr Dunkelman and Stefan Dziembowski, editors, *EUROCRYPT 2022*, volume 13276 of *LNCS*, pages 367–396. Springer, 2022.
ePrint 2021/1359.
-  Youssef El Housni, Aurore Guillevic, and Thomas Piellard.
Co-factor clearing and subgroup membership testing on pairing-friendly curves.
In Lejla Batina and Joan Daemen, editors, *AFRICACRYPT'2022*, volume 13503 of *LNCS*, Fes, Morocco, July 2022. Springer.
to appear, ePrint 2022/352.

Definitions

Elliptic curve E defined over a field \mathbb{F}_q
with a *subgroup* \mathbb{G} of prime order r
 \mathcal{O} neutral element (point at infinity)

P, Q points of $E(\mathbb{F}_q)$

- Co-factor **clearing**: $[c]P \in \mathbb{G}$
- Subgroup **membership testing**: $[r]Q \stackrel{?}{=} \mathcal{O}$

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Scalar multiplication on elliptic curves (Double-and-Add)

Input: Elliptic curve E over \mathbb{F}_q , point $P \in E(\mathbb{F}_q)$, scalar $m \in \mathbb{Z}$

Output: $[m]P$

```
1 if  $m = 0$  then
2   return  $\mathcal{O}$ 
3 if  $m < 0$  then
4    $m \leftarrow -m$ ;  $P \leftarrow -P$ 
5 write  $m$  in binary expansion  $m = \sum_{i=0}^{n-1} b_i 2^i$ , where  $b_i \in \{0, 1\}$ 
6  $R \leftarrow P$ 
7 for  $i = n - 2$  downto 0 do
8    $R \leftarrow [2]R$ 
9   if  $b_i = 1$  then
10      $R \leftarrow R + P$ 
11 return  $R$ 
```

Scalar multiplication on elliptic curves (Double-and-Add)

Input: Elliptic curve E over \mathbb{F}_q , point $P \in E(\mathbb{F}_q)$, scalar $m \in \mathbb{Z}$

Output: $[m]P$

1 **if** $m = 0$ **then**

2 **return** \mathcal{O}

3 **if** $m < 0$ **then**

4 $m \leftarrow -m; P \leftarrow -P$

5 write m in binary expansion $m = \sum_{i=0}^{n-1} b_i 2^i$, where $b_i \in \{0, 1\}$

6 $R \leftarrow P$

7 **for** $i = n - 2$ *downto* 0 **do**

8 $R \leftarrow [2]R$

9 **if** $b_i = 1$ **then**

10 $R \leftarrow R + P$

11 **return** R

$\log_2 m$ (Dbl + $\frac{1}{2}$ Add) in average

Multi-scalar multiplication

Input: Elliptic curve E over \mathbb{F}_q , points $P, Q \in E(\mathbb{F}_q)$, scalars $m \geq m' > 0 \in \mathbb{Z}^{+*}$

Output: $[m]P + [m']Q$

- 1 write $m = \sum_{i=0}^{n-1} b_i 2^i$, $m' = \sum_{i=0}^{n'-1} b'_i 2^i$, where $b_i, b'_i \in \{0, 1\}$
 - 2 $S \leftarrow P + Q$
 - 3 **if** $n > n'$ **then** $R \leftarrow P$
 - 4 **else** $R \leftarrow S$ ($n = n'$)
 - 5 **for** $i = n - 2$ **downto** 0 **do**
 - 6 $R \leftarrow [2]R$
 - 7 **if** $b_i = 1$ **and** $n' \geq i$ **and** $b'_i = 1$ **then**
 - 8 $R \leftarrow R + S$
 - 9 **else if** $b_i = 1$ **and** ($n' < i$ **or** $b'_i = 0$) **then**
 - 10 $R \leftarrow R + P$
 - 11 **else if** $n' \geq i$ **and** $b'_i = 1$ **then**
 - 12 $R \leftarrow R + Q$
 - 13 **return** R
-

Multi-scalar multiplication

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$\log_2 m$ (Dbl + $\frac{3}{4}$ Add) in average

Gallant–Lambert–Vanstone (GLV) with endomorphism

An example: $j = 0$

Let $E: y^2 = x^3 + b$ defined over a prime field \mathbb{F}_q where $q \equiv 1 \pmod{3}$.

There exists $\omega \in \mathbb{F}_q$ such that $\omega^3 = 1$, $\omega \neq 1$

$$\omega^3 - 1 = \underbrace{(\omega - 1)}_{\neq 0} \underbrace{(1 + \omega + \omega^2)}_{=0} = 0$$

$$\begin{aligned} \phi: E(\mathbb{F}_q) &\rightarrow E(\mathbb{F}_q) \\ P(x, y) &\mapsto (\omega x, y), \text{ where } \omega \in \mathbb{F}_q, \omega^2 + \omega + 1 = 0 \end{aligned}$$

ϕ is an **endomorphism**,

$\phi^2: (x, y) \mapsto (\omega^2 x, y)$, $\phi^3 = \text{Id}$ because $\omega^3 = 1$, but $\phi \neq \text{Id} \implies \phi^2 + \phi + 1 = 0$

ℓ -torsion points

Let $E: y^2 = x^3 + ax + b/\mathbb{F}_q$

$$E[\ell] = \{P \in E: [\ell]P = \mathcal{O}\}$$

and $\mathcal{O} \in E[\ell]$

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and $\mathcal{O} \in E[\ell]$

Example

$\ell = 2$, $q \geq 5$: points of order 2 have $y = 0 \iff x^3 + ax + b = 0$.

Factor $x^3 + ax + b$ in \mathbb{F}_q :

- $x^3 + ax + b$ has no root in \mathbb{F}_q : $E(\mathbb{F}_q)[2] = \{\mathcal{O}\}$
- $(x - e_0)(x^2 + ux + v)$ over \mathbb{F}_q : $E(\mathbb{F}_q)[2] = \{(e_0, 0), \mathcal{O}\}$
- $(x - e_0)(x - e_1)(x - e_2)$ over \mathbb{F}_q : $E(\mathbb{F}_q)[2] = \{(e_0, 0), (e_1, 0), (e_2, 0), \mathcal{O}\}$

There exists an extension \mathbb{F}_{q^i} such that $E(\mathbb{F}_{q^i})[2] = \{(x_0, 0), (x_1, 0), (x_2, 0), \mathcal{O}\}$

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$$\ell \text{ coprime to } q, \#E[\ell] = \ell^2$$

ℓ -torsion points

Let ℓ coprime to q , the structure of the points of ℓ -torsion is

$$E[\ell] = \mathbb{Z}/\ell\mathbb{Z} \times \mathbb{Z}/\ell\mathbb{Z}$$

a $\mathbb{Z}/\ell\mathbb{Z}$ two-dimensional vector space.

→ there exists a basis $\{P, Q\}$, with P, Q of order ℓ and “independent”.

Endomorphism ϕ with basis $\{P, Q\}$

$$\phi(P) = [a]P + [c]Q$$

$$\phi(Q) = [b]P + [d]Q$$

$$\phi \leftrightarrow \begin{bmatrix} a & b \\ c & d \end{bmatrix} \in \mathcal{M}_{2,2}(\mathbb{Z}/\ell\mathbb{Z}) \text{ w.r.t. basis } \{P, Q\}$$

Gallant–Lambert–Vanstone (GLV)

$$E: y^2 = x^3 + b$$

ℓ is prime, $\ell \mid \#E(\mathbb{F}_q)$, $\ell^2 \nmid \#E(\mathbb{F}_q)$:

$P \in E(\mathbb{F}_q)[\ell]$, $Q \notin E(\mathbb{F}_q)$ but over an extension of \mathbb{F}_q

$$\implies \phi(P) = [a]P + [0]Q = [\lambda]P$$

where $\lambda \bmod \ell$ is the **eigenvalue** of ϕ : $\lambda^2 + \lambda + 1 = 0 \bmod \ell$, $\approx \sqrt{\ell} \leq |\lambda| \leq \ell - 1$.

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To speed-up $[m]P$, decompose $m = m_0 + m_1\lambda$ with $|m_0|, |m_1| \approx \sqrt{\ell}$ and use $[m]P = [m_0]P + [m_1\lambda]P = [m_0]P + [m_1] \underbrace{\phi(P)}_{\text{cheap}}$ with **multi-scalar** multiplication

$$\frac{1}{2} \log_2 \ell \left(\text{Dbl} + \frac{3}{4} \text{Add} \right)$$

instead of $\log_2 |m| \left(\text{Dbl} + \frac{1}{2} \text{Add} \right) \implies$ **factor ≈ 2 speed-up** but cost of decomposition

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Bilinear pairing

$(\mathbf{G}_1, +), (\mathbf{G}_2, +), (\mathbf{G}_T, \cdot)$ three cyclic groups of large prime order ℓ

Pairing: map $e : \mathbf{G}_1 \times \mathbf{G}_2 \rightarrow \mathbf{G}_T$

1. bilinear: $e(P_1 + P_2, Q) = e(P_1, Q) \cdot e(P_2, Q)$, $e(P, Q_1 + Q_2) = e(P, Q_1) \cdot e(P, Q_2)$
2. non-degenerate: $e(G_1, G_2) \neq 1$ for $\langle G_1 \rangle = \mathbf{G}_1$, $\langle G_2 \rangle = \mathbf{G}_2$
3. efficiently computable.

Most often used in practice:

$$e([a]P, [b]Q) = e([b]P, [a]Q) = e(P, Q)^{ab} .$$

Focus on \mathbf{G}_1 : Endomorphism on an elliptic curve

$$E: y^2 = x^3 + b/\mathbb{F}_q, \quad q \equiv 1 \pmod{3}, \quad j(E) = 0$$

$\mathbf{G}_1 \subset E(\mathbb{F}_q)$ subgroup of prime order

- $r = \#\mathbf{G}_1$ is prime
- $r \mid \#E(\mathbb{F}_q)$
- $r^2 \nmid \#E(\mathbb{F}_q)$

$\implies \phi$ acts as $[\lambda]$ in \mathbf{G}_1 , and $\lambda^2 + \lambda + 1 = 0 \pmod{r}$

Given $m \in \mathbb{Z}/r\mathbb{Z}$, decompose $m = m_0 + m_1\lambda \pmod{r}$ with $|m_0|, |m_1| \approx \sqrt{r}$

Focus on \mathbf{G}_1 : Endomorphism on an elliptic curve

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Given $m \in \mathbb{Z}/r\mathbb{Z}$, decompose $m = m_0 + m_1\lambda \pmod{r}$ with $|m_0|, |m_1| \approx \sqrt{r}$

No computable endomorphism on most of standard curves (NIST, Edwards 25519...)

Exception: Four- \mathbb{Q} , ordinary characteristic 2 \mathbb{F}_{2^n} (GLS-254 Aardal–Aranha SAC'22 ePrint 2022/748)

Barreto, Lynn, Scott method to get pairing-friendly curves.

Becomes more and more popular, replacing BN curves

$$E_{BLS} : y^2 = x^3 + b/\mathbb{F}_q, \quad q \equiv 1 \pmod{3}, \quad j(E) = 0, \quad D = -3 \text{ (ordinary)}$$

$$q = (u - 1)^2/3(u^4 - u^2 + 1) + u$$

$$t = u + 1$$

$$r = (u^4 - u^2 + 1) = \Phi_{12}(u)$$

$$q + 1 - t = (u - 1)^2/3(u^4 - u^2 + 1)$$

$$t^2 - 4q = -3y(u)^2 \rightarrow \text{no CM method needed}$$

BLS12-381 with seed $u_0 = -0xd201000000010000$

BLS12 curves, testing if $P \in \mathbf{G}_1$ for $P \in E(\mathbb{F}_q)$

Well-known GLV trick: write $r_0 + r_1\lambda = 0 \pmod r$
with λ the eigenvalue of $\phi \pmod r$, $\lambda = -u^2$.

$$\underbrace{1}_{r_0} + \underbrace{(1 - u^2)}_{r_1}\lambda = r = u^4 - u^2 + 1$$

Compute $P + [1 - u^2]\phi(P) = ? \mathcal{O}$

Works because ϕ is a distortion map on the cofactor subgroup

$$P \in E(\mathbb{F}_q)[r] \implies \phi(P) = [\lambda]P$$

but no \iff in the general case unless r prime and $\gcd(r, \#E(\mathbb{F}_q)/r) = 1$.

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Order $\#E(\mathbb{F}_q) = 3\ell^2 r$ where $\ell = (u - 1)/3$, $r = u^4 - u^2 + 1$

Co-factor clearing

Given $P \in E(\mathbb{F}_q)$ (e.g. result of a hash map $\{0, 1\}^* \rightarrow E(\mathbb{F}_q)$), compute $[c]P$ where $c = \#E(\mathbb{F}_q)/\#\mathbf{G}_1$

Wahby–Boneh, CHES'2019: $c = 3\ell^2$ but no point of order ℓ^2 , only points of order dividing ℓ

\implies compute only $[\ell]P$

Luck or generic pattern?

Schoof's theorem 3.7 (1987), simplified

 René Schoof.

Nonsingular plane cubic curves over finite fields.

Journal of Combinatorial Theory, Series A, 46(2):183–211, 1987.

$$E[\ell] \subset E(\mathbb{F}_q) \iff \begin{cases} \ell^2 \mid \#E(\mathbb{F}_q), \\ \ell \mid q - 1 \text{ and} \\ \pi_q \in \mathbb{Z} \text{ or } \mathcal{O}\left(\frac{t^2 - 4q}{\ell^2}\right) \subset \text{End}_{\mathbb{F}_q}(E) \\ \rightarrow \text{simplifies to } \ell \mid y \text{ where } t^2 - 4q = -Dy^2 \end{cases}$$

Generic pattern for all BLS curves

BLS- k curves, $3 \mid k$

- $c = (x - 1)^2 / 3(x^{2k/3} + x^{k/3} + 1) / \Phi_k(x)$, $k = 3 \pmod{6}$
- $c = (x - 1)^2 / 3(x^{k/3} - x^{k/6} + 1) / \Phi_k(x)$, $k = 0 \pmod{6}$

and $E(\mathbb{F}_q)[\ell] = \mathbb{Z}/\ell\mathbb{Z} \times \mathbb{Z}/\ell\mathbb{Z}$ where $\ell = (x - 1)/3$.

Other pairing-friendly curves

For all curves in the Taxonomy paper of Freeman, Scott, Teske,

- we identify the families such that the cofactor has a square factor
- we check the conditions of Schoof's theorem
- we list the curves with faster co-factor clearing: all but KSS and 6.6 where $k \equiv 2, 3 \pmod{6}$.

SageMath verification script at

`gitlab.inria.fr/zk-curves/cofactor`

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\mathbf{G}_2 technicalities

\mathbf{G}_2 is more tricky and the endomorphism is ψ , of characteristic polynomial

$$X^2 - tX + q$$

where t is the trace of E over \mathbb{F}_q .

GLV on $\mathbf{G}_1 \rightarrow$ GLS (Galbraith Lin Scott) on \mathbf{G}_2

A point $Q \in E'(\mathbb{F}_{q^i})$ has some eigenvalue μ under ψ is a *consequence* of Q having order r

- flaw in Scott's proof identified
 - and fixed
 - corner cases under control
- all safe as long as r is prime

\mathbf{G}_T membership testing

$$\mathbf{G}_T = \mu_r = \{z \in \mathbb{F}_{q^k}^*, z^r = 1\}$$

Proposition

- $E: y^2 = x^3 + ax + b/\mathbb{F}_q$
- prime $r \mid \#E(\mathbb{F}_q)$, $r^2 \nmid \#E(\mathbb{F}_q)$
- $E[r] \subset E(\mathbb{F}_{q^k})$ and k is minimal $\iff \mathbf{G}_T \subset \mathbb{F}_{q^k}^*$

Let $z \in \mathbb{F}_{q^k}^*$.

$$z^{\Phi_k(q)} = 1 \text{ and } z^q = z^{t-1} \text{ and } \gcd(q+1-t, \Phi_k(q)) = r \implies z^r = 1 (z \in \mathbf{G}_T)$$

Future work

- fix the problem of $m_0 + m_1\lambda = h \cdot r$ and h is not coprime to the cofactor
hint of the fix in ePrint 2022/348 and hal-00874925 § 4 *Shrinking the basis*
- alternative def of \mathbf{G}_2 : trace-zero subgroup, $\ker \xi \circ (1 + \pi_q + \pi_{q^2} + \dots + \pi_{q^{k-1}}) \circ \xi^{-1}$
early abort test?
- Apply to other curves, e.g. BW6 for 2-chain SNARKs