Réplication et cohérence de données
(Data replication and consistency)

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B-022 LORIA
Organisation

• Lecture (16h)
  ▪ Last 4h dedicated to presentations of assigned papers (groups of 2 students)

• Exercises (8h)
  ▪ A project (groups of 2 students)

• Web site: https://members.loria.fr/CIIgnat/
Evaluation

• Paper presentation
  ▪ Groups of 2 students
  ▪ 20 min presentation + 10 min questions
• Project
  ▪ Groups of 2 students
• Exam
  ▪ If remotely, documents allowed, otherwise not
Course overview

• Introduction to replication
• Consistency models (*)
• Consistency protocols (*)
• Pessimistic replication vs. optimistic replication
• Optimistic replication approaches

(*) Andrew S. Tanenbaum, Maarten Van Steen, "Distributed Systems: Principles and Paradigms“, 2002
Replication is everywhere

• Facebook
  ▪ >2.7 billion active users

• Twitter
  ▪ 330 million monthly active users
  ▪ 500 million tweets / day
  ▪ 2.1 billion search queries /day

• Google
  ▪ >2.5 million servers (since July 2016)
  ▪ >5.5 billion search requests/day
  ▪ Synchronous data replication (Google Drive, Gmail, Google Sites, Calendar)

• Amazon
  ▪ >310 million active customer accounts
  ▪ ~10,000$ revenue/second
Introduction

• Distributed system
  ▪ Sites (processing unit + storage device)
  ▪ Communication links (bidirectional communication between 2 sites)
  ▪ Messages (no guarantee to be delivered within a maximum delay)

• Failures of sites and links
  ▪ An error leads to a faulty state
  ▪ Errors: human mistakes or physical damage
  ▪ Site failures: stopping, crash of a critical subsystem, malicious actions (byzantine failures)
  ▪ Link failures: messages are no longer transmitted or excessively delayed, unidirectional transmission, byzantine communication failure
  ▪ One consequence: network partition
Reliability (Fiabilité)

- Property of tolerating component failures for the longest time
- A system is perfectly reliable if it never fails
- A system is reliable if it fails rarely and almost always recovers from component failures and design faults s.t. its activity is resumed without perceptible interruption
- A system is reliable to the extent that it is able to successfully complete a service request once it accepts it
Availability (Disponibilité)

- Accessibility of system services to users
- A system is highly available if the fraction of its down-time is very small (failures are rare or it can restart very quickly after a failure)
- A system is highly available if denial of service request is rare
Reliability vs. Availability

- Reliability: duration of time a system is expected to remain in continuous operation
- Availability: fraction of time instants where the system is operational
- A reliable system is not necessarily highly available
- A highly available system is not necessarily reliable
- Dependability = Reliability \times Availability
Reasons for replication

• Reliability
  ▪ System continues to work if one replica crashes
  ▪ Better protection against corrupted data

• Performance
  ▪ Scaling in numbers - no overloading of a server (replicated web servers)
  ▪ Scaling with the size of a geographical area – reduced communication latency (web cache)

• Challenge: how to maintain consistency between replicated data?
CAP Theorem

• Properties of shared-data systems
  ▪ Data consistency
  ▪ System availability
  ▪ Tolerance to network partition

• Only two out of three properties can be achieved at a given time

CAP Theorem

• A system without network partition, can achieve consistency + availability
  ▪ Client and storage system are part of the same environment

• In a large-scale distributed systems, network partitions exist
  ▪ Consistency and availability cannot be both achieved
    ○ Relax consistency, maintain availability
    ○ Maintain consistency, tolerate unavailability under certain conditions
Data-Centric Consistency Models

• The general organization of a logical data store, physically distributed and replicated across multiple processes.
Data-Centric Consistency Models

• Write operation when it changes data, otherwise read operation

• Consistency model=contract between processes and data store
  ▪ If processes agree to obey certain rules, data store promises to work correctly

• A process that performs a read expects a return value that shows results of last write

• Too strict criteria as lack of global clock, need for other consistency models

• Each model restricts the values that a read can return
Strict Consistency

- **Any read on a data item** $x$ **returns a value corresponding to the result of the most recent write on** $x$

- Assumes existence of absolute global time

- **Two main issues:**
  - Definition of “most recent event”
  - Instantaneous execution of operations

- **Example:**
  - $X$ a data item stored at machine $B$
  - Machine $A$ reads $X$ at time $T_1$ and message sent to $B$ to read $X$
  - At $T_2$, a process on $B$ writes $X$
  - Read should return old value of $X$ regardless of where machines are and how close $T_2$ and $T_1$ is
Strict Consistency

- “most recent event” needs perfectly synchronised clocks (0 delay between any 2 sites)
- An operation that requires a remote access cannot be executed instantaneously (a local operation launched after a remote operation can be terminated before)
- Notation:
  - $W_i(x)a$ – process $P_i$ writes $x$ with value $a$
  - $R_i(x)b$ – process $P_i$ reads $x$ with value $b$
Strict Consistency

Behavior of two processes, operating on the same data item.

\[
\begin{align*}
&P1: \ W(x)a \\
&P2: \ R(x)a \\
&\text{(a)}
\end{align*}
\]

\[
\begin{align*}
&P1: \ W(x)a \\
&P2: \ R(x)\text{NIL} \quad R(x)a \\
&\text{(b)}
\end{align*}
\]

A strictly consistent store

A store that is not strictly consistent.

- Strict consistency is an ideal model
- Need for relaxed consistency models
Sequential Consistency (1)

- Sequential consistency (defined by Lamport)
  - *The result of any execution is the same as if the (read and write) operations by all processes on the data store were executed in some sequential order and the operations of each individual process appear in this sequence in the order specified by its program*
  - Any valid interleaving of operations is possible, but all processes see the same interleaving
  - There is no reference to the “most recent” write
  - A process sees writes from all processes but only its own reads
## Sequential Consistency (2)

<table>
<thead>
<tr>
<th></th>
<th>P1: W(x)a</th>
<th>P1: W(x)a</th>
</tr>
</thead>
<tbody>
<tr>
<td>P2:</td>
<td>W(x)b</td>
<td>W(x)b</td>
</tr>
<tr>
<td>P3:</td>
<td>R(x)b</td>
<td>R(x)b</td>
</tr>
<tr>
<td>P4:</td>
<td>R(x)b</td>
<td>R(x)a</td>
</tr>
<tr>
<td>(a)</td>
<td></td>
<td></td>
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</tbody>
</table>

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<td>R(x)b</td>
<td>R(x)b</td>
</tr>
<tr>
<td>P4:</td>
<td>R(x)a</td>
<td>R(x)b</td>
</tr>
<tr>
<td>(b)</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

**a)** A sequentially consistent data store.

**b)** A data store that is not sequentially consistent.
Linearizability (1)

- Linearizability weaker than strict consistency but stronger than sequential consistency
  - $ts_{OP}(x)$ timestamp assigned to operation OP performed on data item x
  - The result of any execution is the same as if the (read and write) operations by all processes on the data store were executed in some sequential order and the operations of each individual process appear in this sequence in the order specified by its program. In addition if $ts_{OP1}(x) < ts_{OP2}(y)$, then operation $OP_1(x)$ should precede $OP_2(y)$ in this sequence
  - A linearizable data store is also sequentially consistent
Example Sequential Consistency (1)

Process P1: 
- x = 1;
- print (y, z);

Process P2: 
- y = 1;
- print (x, z);

Process P3: 
- z = 1;
- print (x, y);
Example Sequential Consistency (2)

- Four valid execution sequences

<table>
<thead>
<tr>
<th>Sequence 1</th>
<th>Sequence 2</th>
<th>Sequence 3</th>
<th>Sequence 4</th>
</tr>
</thead>
<tbody>
<tr>
<td>x = 1;</td>
<td>y = 1;</td>
<td>x = 1;</td>
<td>y = 1;</td>
</tr>
<tr>
<td>print (y, z);</td>
<td>y = 1;</td>
<td>print (x,z);</td>
<td>z = 1;</td>
</tr>
<tr>
<td>y = 1;</td>
<td>print (x, z);</td>
<td>print (x, y);</td>
<td>z = 1;</td>
</tr>
<tr>
<td>print (x, y);</td>
<td>print(y, z);</td>
<td>print (x, z);</td>
<td>print (x, z);</td>
</tr>
<tr>
<td>z = 1;</td>
<td>z = 1;</td>
<td>x = 1;</td>
<td>print (y, z);</td>
</tr>
<tr>
<td>print (x, y);</td>
<td>print (x, y);</td>
<td>print (y, z);</td>
<td>print (x, y);</td>
</tr>
</tbody>
</table>

Prints: 001011  Prints: 101011  Prints: 010111  Prints: 111111
Signature: 001011  Signature: 101011  Signature: 110101  Signature: 111111

- 000000 not permitted
- 001001 not allowed
Formal Expression Sequential Consistency (1)

- Each process $P_i$ has an associated execution $E_i$ of read and write performed on data store $S$

- $E_1$: $W_1(x)a$
- $E_2$: $W_2(x)b$
- $E_3$: $R_3(x)b$, $R_3(x)a$
- $E_4$: $R_4(x)b$, $R_4(x)a$
- Merge $E_i$ into $H$ s.t. each operation in $E_i$ appears in $H$ once
Formal Expression Sequential Consistency (2)

• Legal values for H must obey the constraints
  ▪ Program order must be maintained: if OP₁ before OP₂ in Eᵢ, then OP₁ before OP₂ in H
  ▪ Data coherence must be respected: R(x) must return the value most recently written to x

• H=W₂(x)b, R₃(x)b, R₄(x)b, W₁(x)a, R₃(x)a, R₄(x)a
Linearizability and Sequential Consistency

Contrainte de temps : \( W(x)a \) précede \( W(x)b \)

\[
\begin{align*}
p_1 & \quad W(x)a \\
p_2 & \quad W(x)b \\
p_3 & \quad R(x)a \quad R(x)b \\
p_4 & \quad R(x)a \quad R(x)b
\end{align*}
\]

Notation :
\( R_i(x)a \) : par \( p_i \), lecture de \( x \), résultat \( a \)
\( W_i(x)a \) : écriture de \( x \), valeur \( a \)

\( S = W(x)a \quad R_3(x)a \quad R_4(x)a \quad W(x)b \quad R_3(x)b \quad R_4(x)b \)

linéarisable

\[
\begin{align*}
p_1 & \quad W(x)a \\
p_2 & \quad W(x)b \\
p_3 & \quad R(x)a \quad R(x)b \\
p_4 & \quad R(x)a \quad R(x)b
\end{align*}
\]

\( S = W(x)a \quad R_3(x)a \quad R_4(x)a \quad W(x)b \quad R_4(x)b \quad R_3(x)b \)

séquentiel, non linéarisable

\[
\begin{align*}
p_1 & \quad W(x)a \\
p_2 & \quad W(x)b \\
p_3 & \quad R(x)b \quad R(x)a \\
p_4 & \quad R(x)a \quad R(x)b
\end{align*}
\]

non séquentiel
Sequential Consistency

• It is costly to be realised
  ▪ $t$ the minimal transfer time of a message between two sites
  ▪ $r$ the reading period
  ▪ $w$ the writing period
  ▪ $r+w \geq t$
  ▪ Gain on reading results in a lose of writing time
Causal Consistency (1)

• If event B is caused or influenced by an earlier event A, causality requires that everyone sees A and then B

• Concurrent operations = operations not causally related

• Examples:
  - If \(E_1\): p writes x, then \(E_2\): q reads x, then \(E_1 \rightarrow E_2\)
  - If \(E_1\): p reads x, then \(E_2\): p writes y, then \(E_1 \rightarrow E_2\) (value of y depends on x)
  - If \(E_1\): p writes x, then \(E_2\): q writes y (independently), then \(E_1 \parallel E_2\)
Causal Consistency (2)

- **Writes that are potentially causally related must be seen by all processes in the same order.** Concurrent writes may be seen in a different order on different machines.
### Causal Consistency (3)

<table>
<thead>
<tr>
<th></th>
<th>W(x)a</th>
<th>R(x)a</th>
<th>W(x)b</th>
<th>R(x)c</th>
<th>R(x)b</th>
<th>R(x)c</th>
</tr>
</thead>
<tbody>
<tr>
<td>P1:</td>
<td></td>
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<td>P2:</td>
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<td>P3:</td>
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<tr>
<td>P4:</td>
<td></td>
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<td></td>
</tr>
</tbody>
</table>

- Allowed with a causally-consistent store
- Not allowed with sequentially or strictly consistent store.
Causal Consistency (4)

P1: \( W(x)a \)

P2: \( R(x)a \quad W(x)b \)

P3: \( R(x)b \quad R(x)a \)

P4: \( R(x)a \quad R(x)b \)

(a) \( W_2(x)b \) depends on \( W_1(x)a \)

P1: \( W(x)a \)

P2: \( W(x)b \)

P3: \( R(x)b \quad R(x)a \)

P4: \( R(x)a \quad R(x)b \)

(b)

a) A violation of a casually-consistent store.

b) A correct sequence of events in a casually-consistent store.
Causal Consistency (5)

- Implementation needs keeping track of which processes have seen which writes
- Construction of a dependency graph
- One solution based on vector timestamps
FIFO Consistency (1)

- **Writes done by a single process are seen by all other processes in the order in which they were issued, but writes from different processes may be seen in a different order by different processes.**
FIFO Consistency (2)

<table>
<thead>
<tr>
<th>P1: W(x)a</th>
</tr>
</thead>
<tbody>
<tr>
<td>P2: R(x)a</td>
</tr>
<tr>
<td>P3:</td>
</tr>
<tr>
<td>P4:</td>
</tr>
</tbody>
</table>

- valid FIFO consistency
FIFO Consistency (3)

- Result impossible to obtain with sequential consistency

\[
\begin{align*}
&\text{x = 1;} \\
&\text{print (y, z);} \\
&\text{y = 1;} \\
&\text{print(x, z);} \\
&\text{z = 1;} \\
&\text{print (x, y);} \\
\text{Prints: 00}
\end{align*}
\]

\[
\begin{align*}
&\text{x = 1;} \\
&\text{y = 1;} \\
&\text{print(x, z);} \\
&\text{print (y, z);} \\
&\text{z = 1;} \\
&\text{print (x, y);} \\
\text{Prints: 10}
\end{align*}
\]

\[
\begin{align*}
&\text{y = 1;} \\
&\text{print (x, z);} \\
&\text{z = 1;} \\
&\text{print (x, y);} \\
&\text{x = 1;} \\
&\text{print (y, z);} \\
\text{Prints: 01}
\end{align*}
\]
FIFO Consistency (4)

With FIFO consistency both processes $P_1$ and $P_2$ can be killed

<table>
<thead>
<tr>
<th>Process P1</th>
<th>Process P2</th>
</tr>
</thead>
<tbody>
<tr>
<td>$x = 1;$</td>
<td>$y = 1;$</td>
</tr>
<tr>
<td>if ($y == 0$) kill (P2);</td>
<td>if ($x == 0$) kill (P1);</td>
</tr>
</tbody>
</table>
## Summary of Consistency Models

<table>
<thead>
<tr>
<th>Consistency</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Strict</td>
<td>Absolute time ordering of all shared accesses matters.</td>
</tr>
<tr>
<td>Linearizability</td>
<td>All processes must see all shared accesses in the same order. Accesses are furthermore ordered according to a (nonunique) global timestamp</td>
</tr>
<tr>
<td>Sequential</td>
<td>All processes see all shared accesses in the same order. Accesses are not ordered in time</td>
</tr>
<tr>
<td>Causal</td>
<td>All processes see causally-related shared accesses in the same order.</td>
</tr>
<tr>
<td>FIFO</td>
<td>All processes see writes from each other in the order they were used. Writes from different processes may not always be seen in that order</td>
</tr>
</tbody>
</table>
Data Centric vs. Client Centric Consistency Models

• Data-centric consistency model
  ▪ Provides a systemwide consistent view on a data store
  ▪ Assumption: concurrent processes may simultaneously update the data store

• Client-centric consistency model
  ▪ Assumption: lack of simultaneous updates or when updates happen they can easily be resolved
  ▪ Most operations involve reading data
  ▪ Data stores offer a very weak consistency model called eventual consistency
Eventual Consistency (1)

- Concurrency appears only in a restricted form
  - Database systems
    - most processes hardly perform updates, but only read data
  - DNS
    - domains assigned to naming authorities that are allowed to update their part of the name space
    - no write-write conflicts, only read-write conflicts
    - Acceptable to propagate an update in a lazy fashion
  - Web pages
    - Usually updated by a single authority
    - No write-write conflicts
    - Design of local caches for efficiency
- If no updates take place for a long time, all replicas become consistent
Eventual Consistency (2)

- The principle of a mobile user accessing different replicas of a distributed database.

- Client moves to other location and (transparently) connects to other replica.

- Replicas need to maintain client-centric consistency.

- Portable computer

- Distributed and replicated database

- Read and write operations

- Wide-area network
Client centric consistency

- Provides guarantees for a single client concerning consistency of accesses to a data store by that client
- A client connects to different replicas during a period of time and the differences should be made transparent
- Whenever a client connects to a new replica, that replica is brought up to date with the data that was manipulated by that client before and can reside at other replica sites

Notations
- $x_i[t]$ version of $x$ at local copy $L_i$ at time $t$
- $WS(x_i[t])$ series of write operations at $L_i$ that took place since initialization
- $WS(x_i[t_1];x_j[t_2])$ if operations in $WS(x_i[t_1])$ have been performed at the local copy $L_j$ at time $t_2$
Monotonic Reads (1)

• If a process reads the value of a data item \( x \), any successive read operation on \( x \) by that process will always return that same value or a more recent value

• Example:
  - A distributed email database
  - Each user’s mailbox may be distributed and replicated across multiple machines
  - Mail can be inserted at any location
  - Updates propagated in a lazy manner
  - Emails read (no remove, etc.) at location X are present when they are read later at location Y
Monotonic Reads (2)

\[
\begin{align*}
L1: & \quad WS(x_1) & \quad R(x_1) \\
L2: & \quad WS(x_1;x_2) & \quad R(x_2)
\end{align*}
\]

(a)

\[
\begin{align*}
L1: & \quad WS(x_1) & \quad R(x_1) \\
L2: & \quad WS(x_2) & \quad R(x_2) & \quad WS(x_1;x_2)
\end{align*}
\]

(b)

a) A monotonic-read consistent data store

b) A data store that does not provide monotonic reads.
Monotonic Writes (1)

- A write operation by a process on a data item $x$ is completed before any successive write operation on $x$ by the same process
- Similarity with FIFO consistency
- Example: updating a software library
Monotonic Writes (2)

\[
\begin{align*}
L1: & \quad W(x_1) \\
L2: & \quad W(x_1) \quad W(x_2) \\
\text{(a)} \\
\end{align*}
\]

\[
\begin{align*}
L1: & \quad W(x_1) \\
L2: & \quad \quad \quad W(x_2) \\
\text{(b)} \\
\end{align*}
\]

\begin{itemize}
  \item [a)] A monotonic-write consistent data store.
  \item [b)] A data store that does not provide monotonic-write consistency.
\end{itemize}
Read Your Writes (1)

• The effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same process

• Examples
  ▪ Updating passwords
  ▪ Updating web pages
Read Your Writes (2)

\[
\begin{array}{c}
\text{L1:} \quad W(x_1) \\
\text{L2:} \quad WS(x_1,x_2) \quad R(x_2)
\end{array}
\]

(a)

\[
\begin{array}{c}
\text{L1:} \quad W(x_1) \\
\text{L2:} \quad WS(x_2) \quad R(x_2)
\end{array}
\]

(b)

a) A data store that provides read-your-writes consistency.

b) A data store that does not.
Writes Follow Reads (1)

- A write operation by a process on a data item x following a previous read operation on x by the same process, is guaranteed to take place on the same or a more recent value of x that was read.

- Example:
  - Network newsgroup that see a posting of a reaction to an article only after they have seen the original article.
Writes Follow Reads (1)

\[
\begin{array}{c|c|c}
\text{L1:} & WS(x_1) & R(x_1) \\
\hline
\text{L2:} & WS(x_1;x_2) & W(x_2) \\
\hline
\end{array}
\]

(a)

\[
\begin{array}{c|c|c}
\text{L1:} & WS(x_1) & R(x_1) \\
\hline
\text{L2:} & WS(x_2) & W(x_2) \\
\hline
\end{array}
\]

(b)

a) A writes-follow-reads consistent data store

b) A data store that does not provide writes-follow-reads consistency
Examples Systems Consistency

- Microsoft’s Windows Azure – strongly consistent storage
  - Clients see the latest written value for a data object

- Amazon Simple Storage Service (S3) - eventual consistency
  - Relaxed consistency for better performance and availability
  - Returned value for reads=object value at some past point in time but not necessarily the latest value

- Amazon’s DynamoDB and Google App Engine Datastore – both eventually consistent reads and strongly consistent reads
Consistency protocols

• Describe implementation of a specific consistency model

• Primary-based protocols
  ▪ Each data item x has an associated primary responsible for coordinating write operations on x
  ▪ Remote-write protocols
  ▪ Local-write protocols

• Replicated-write protocols
  ▪ Write operations can be carried out at multiple replicas
  ▪ Active replication
  ▪ Quorum-based protocols
Remote-Write Protocols (1)

- Primary-based remote-write protocol with a fixed server to which all read and write operations are forwarded.

W1. Write request
W2. Forward request to server for x
W3. Acknowledge write completed
W4. Acknowledge write completed

R1. Read request
R2. Forward request to server for x
R3. Return response
R4. Return response
Remote-Write Protocols (2)

W1. Write request
W2. Forward request to primary
W3. Tell backups to update
W4. Acknowledge update
W5. Acknowledge write completed
R1. Read request
R2. Response to read

- The principle of primary-backup protocol.
Remote-Write Protocols (3)

- Primary-backup protocol implements update as a blocking operation
- Alternative solution: non-blocking protocol
  - As soon as primary updated the local copy of x, it returns an acknowledgement
  - After that ask backup servers to perform the update
  - Fault tolerance concerns
- Implementation of sequential consistency
Local-Write Protocols (1)

1. Read or write request
2. Forward request to current server for \( x \)
3. Move item \( x \) to client's server
4. Return result of operation on client's server

• Primary-based local-write protocol in which a single copy is migrated between processes
• Disadvantage: keeping track where each data item currently is
Local-Write Protocols (2)

- **Primary-backup protocol** in which the primary migrates to the process wanting to perform an update.
- **Advantage** if nonblocking protocol: write operations carried locally, while reading can access local copies.
- **Protocol** suitable for mobile computers.

**Diagram:***

- Client
- Old primary for item x
- New primary for item x
- Backup server
- Data store

**Transactions:**
- **W1.** Write request
- **W2.** Move item x to new primary
- **W3.** Acknowledge write completed
- **W4.** Tell backups to update
- **W5.** Acknowledge update

- **R1.** Read request
- **R2.** Response to read

- **W4**
- **W5**

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Réplication et cohérence de données

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Replicated write protocols

Active Replication (1)

• Operations sent to each replica
• Operations have to be carried out in the same order everywhere
  - Need of totally-ordered multicast
    o Using Lamport timestamps
    o Using a central coordinator called sequencer
• Deal with replicated invocations
Active Replication (2)

- Client replicates invocation request
- All replicas see the same invocation
- Replicated object
- Object receives the same invocation three times
Active Replication (3)

(a) Client replicates invocation request

(b) Coordinator of object C

Result
Quorum-Based Protocols (1)

• Use voting: clients request and acquire permission of multiple servers before reading/writing a replicated object

• Example distributed file system
  - File replicated on N servers
  - For an update a client must contact a majority of servers (half +1)
  - If agreement file changed and version number updated
  - For a read a client must contact at least half of servers +1 and ask them to send version numbers of the file
  - Choose the most recent version
Quorum-Based Protocols (2) - Gifford scheme

- A file with N replicas
- A read quorum \((N_R\) servers) for reading the file
- A write quorum \((N_W\) servers) for modifying the file
- \(N_R + N_W > N\)
- \(N_W > N/2\)
Quorum-Based Protocols (3)

a) A correct choice of read and write set
b) A choice that may lead to write-write conflicts
c) A correct choice, known as ROWA (read one, write all)
Pessimistic vs. optimistic replication (1)

• Pessimistic replication
  - Give the illusion of one replica (no divergence)
  - Block access to a replica unless it is up-to-date
  - Example: primary-copy algorithms
    - Elect a primary replica
    - After an update primary writes the change to secondary replicas
    - If primary crashes elect a new replica
  - Bad performance and availability
Pessimistic vs. optimistic replication (2)

- Optimistic replication
  - Allows replicas to diverge
    - Commit modifications immediately and propagate later
    - Observers can see different values on different sites
  - Eventual consistency
  - Mandatory for offline access
  - Better scaling
Eventual Consistency

- A history is eventually consistent (EC) when for every object $x$ if there is a bounded amount of write operations on $x$ in $h$, then eventually all the read operations observe the same state.
Strong Eventual Consistency

- **Eventual delivery**: An update executed at some correct replica eventually executes at all correct replicas.

- **Strong convergence**: Correct replicas that have executed the same updates have equivalent states.

- No consensus in background, no need to rollback.
Pessimistic vs. optimistic replication (3)

- Basic principles of optimistic replication
  - N sites replicate an object
  - An object is modified by applying an operation
  - Local operations applied immediately
  - Operations broadcast to the other sites
  - Remote operations integrated and executed
  - System is correct if when it is idle all replicas are identical