# Rectangular Surface Parameterization-Supplemental Information

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### **B** Discrete Poisson Equation

For completeness, we describe the standard *cotan* discretization of the Poisson equation [MacNeal 1949, Section 3.2], used to recover our final parameterization coordinates in Section 5.2.

# B.1 Discrete Laplacian

The *cotan-Laplace* matrix L is then a sparse  $|V| \times |V|$  matrix with nonzero entries

$$L_{i,j} = \begin{cases} -w_{ij}, & ij \in E, \\ \sum_{ik} w_{ik} & i = j. \end{cases}$$
(1)

# B.2 Discrete Divergence

The discrete divergence operator is a sparse  $|V|\times |E|$  matrix with nonzero entries

$$\operatorname{div}_{i,ij} = s_{ij} w_{ij}, \ \forall i \in V, \ ij \in E,$$

$$(2)$$

where  $s_{ij} = +1$  when edge *ij* is oriented from *i* to *j*, and equals -1 otherwise.

# B.3 Discrete Poisson Equation

Finally, we discretize the Poisson equation in Equation 13 as a scalar Poisson equation for each coordinate  $f^1, f^2$  of the final rectangular parameterization  $f: V \to \mathbb{R}^2$ , namely

$$Lf^{p} = \operatorname{div}\mu^{p}, \quad p = 1, 2.$$
(3)

Here *L*, div, and  $\mu$  are the discrete Laplacian, divergence, and differential given in Equations 1, 2, and 18, *resp.*.

# C Underdetermined Newton Solver

To solve our main optimization problem (Equation 19), we need a method that can deal with underdetermined nonlinear constraints. We use the method of Polyak and Tremba [2020], detailed here. Similar to Newton's method, this method finds a zero of the first order optimality conditions by repeatedly solving linear systems, using a carefully-chosen line search.

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Fig. 32. Here we show boundary conditions like those from Figure 18, but without any cone singularities. This example illustrates both the robustness of our method to extreme constraints—and also emphasizes the utility of singularities for obtaining low distortion/good element quality.

More explicitly, let  $x := (u, v, \theta) \in \mathbb{R}^{2|V|+|F|}$  be the primal variables and  $\lambda \in \mathbb{R}^{|E|}$  be the Lagrange multiplier, the critical points of Equation 19 are solutions of the system of equations:

 $\nabla$ 

$$\Phi(x) + J_F(x)^\top \lambda = 0$$
  

$$F(x) = 0,$$
(4)

where  $\nabla \Phi$  denote gradient of the objective chosen in Section 5.1.1, and  $J_F$  is the Jacobian of the constraint equation given in Section 4.1.

The nonlinear system of equations in Equation 4 is solved by taking Newton steps, using a line search that directly encourages satisfaction of this system. More explicitly, let  $x_k$  and  $\lambda_k$  denote the current guess at the *k*th step of optimization. Then we solve the linear system

$$\begin{bmatrix} \nabla^2 \Phi(x_k) + \nabla (J_F(x_k)^\top \lambda_k) & J_F(x_k)^\top \\ J_F(x_k) & 0_{|E| \times |E|} \end{bmatrix} \begin{bmatrix} y_k \\ \delta_k \end{bmatrix} = \begin{bmatrix} \nabla \Phi(x_k) + J_F(x_k)^\top \lambda_k \\ F(x_k) \end{bmatrix}$$
(5)

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for the descent directions  $y_k$  and  $\delta_k$ , where  $\nabla^2 \Phi$  denote the Hessian the objective function. The updated guesses are  $x_{k+1} = x_k - \tau y_k$  and  $\lambda_{k+1} = \lambda_k - \tau \delta_k$ , where  $\tau > 0$  is a step size. The step size is computed by performing backtracking line search until

$$\begin{aligned} \|\nabla\Phi(x_k - \tau y_k) + J_F(x_k - \tau y_k)^\top (\lambda_k - \tau \delta_k)\| + \|F(x_k - \tau y_k)\| \\ &\leq (1 - \frac{\tau}{2}) \left( \|\nabla\Phi(x_k) + J_F(x_k)^\top \lambda_k\| + \|F(x_k)\| \right). \end{aligned}$$

This strategy guarantees that we converge to a local minimizer of the original optimization problem (Equation 19), at essentially the usual rate of Newton's method—see Polyak and Tremba [2020, Theorem 4.1]. In practice, we use a backtracking ratio of 0.9 for line search. Crucially, we find that Equation 5 always admits a solution and is only overconstrained when inadmissible boundary constraints are imposed (Section 5.1.3).

### D Singularities and Cuts

In the case where either (i) the field has singularities and/or (ii) the domain M is not a topological disk, we must effectively cut along the curve  $\Gamma$ , plus additional topological cuts, before flattening the mesh. In practice, we compute topological cuts by considering the traversal tree used to construct  $\Gamma$ , and taking the complementary graph of primal edges not crossed by any edge in this dual tree. We then iteratively remove any degree-1 vertex from the primal graph, except for singular vertices, until no more vertices can be removed. The remaining primal edges define the cut. (This approach is similar in spirit to the *tree-cotree* strategy of Eppstein [2002].)

This problem of solving for a parameterization on a cut mesh comes up frequently in geometry processing, and is a perennial nuisance. Hence, we will first introduce a perspective that simplifies the situation both conceptually and in terms of implementation, before giving the details of our particular algorithm. We make just one simple change to our basic setup: rather than store values at vertices, we will store values at *all* triangle corners, using linear constraints to identify equivalent values as needed. In the absence of any cuts, our formulation ultimately yields the same solution as before, but with the additional flexibility to easily incorporate cuts.

In more detail, for any function  $f_i$  at vertices, we now store values  $f_i^{jk}$  at all triangle corners. A sparse matrix  $U \in \mathbb{R}^{m \times |C|}$  encodes equivalence of values at different corners, where m is the number of identifications, and |C| is the number of corners. Two corners around a common vertex are equivalent if and only if they are not separated by a cut edge. For the *n*th equivalence  $f_i^{jk} = f_i^{kl}$  of consecutive corners, U will have nonzero entries  $U_{n,i}^{jk} = 1$  and  $U_{n,i}^{kl} = -1$ , *i.e.*, the equation number determines the row index, and the corners are specified via the column indices.

Next, suppose we want to solve any linear equation Af = b, which in general may be under- or over-determined. We can express this system in terms of values on corners, plus the identification matrix U, as a least-squares problem subject to linear constraints:

$$\min_{f \in \mathbb{R}^{|C|}} \quad \|Af - b\|_M^2$$
s.t.  $Uf = 0.$ 
(6)

Here  $|A|_M^2 := A^\top M A$  denotes the  $\ell$ -2 norm with a mass matrix  $M \in \mathbb{R}^{|C| \times |C|}$ . We use in particular a diagonal matrix with nonzero

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entries

$$M_{jk,jk}_{i,jk} = \frac{1}{2} \cot \alpha_i^{jk},$$

*i.e.*, we use the cotan weight associated with each corner. In the case where there is no cut, our problem amounts to solving the exact same cotan-Poisson problem in Equation 3.

Equation 6 can be solved using any available method—for instance, the method of Lagrange multipliers, we find the solution to this problem is given by the block linear system

$$\begin{bmatrix} A^{\mathsf{T}}MA & U^{\mathsf{T}} \\ U & 0 \end{bmatrix} \begin{bmatrix} f \\ \lambda \end{bmatrix} = \begin{bmatrix} A^{\mathsf{T}}Mb \\ 0 \end{bmatrix}$$

where  $\lambda \in \mathbb{R}^m$  are the Lagrange multipliers. By formulating the problem this way, rather than directly applying least-squares to the combined system  $[AU]^{\top}f = [b0]$ , we ensure that corner values are identified *exactly*, and only the original equations are approximated. Note that in general there may be more efficient ways to solve this kind of problem—we find that this approach is simple, and helps avoid implementation errors. For instance, one could directly eliminate variables by hand (though modern direct solvers are often quite good at effectively performing this same elimination); a general-purpose QP solver might also work fine depending on the use case.

In the case of our particular problem, we also have to modify Equation 3, splitting it into two different equations on the two sides of every edge. In particular, for each edge ij in  $\Gamma$  we modify the right-hand side so that both vectors  $\mu_{ij}^k$ ,  $\mu_{ji}^l$  are in the same coordinate system. In particular, consider the angles  $\beta_{ij}$  from Section 5.3.1, which give the rotation between frames (after parallel transport) across any cut edge ij. Hence, to evaluate Equation 3 for an edge from corner  $_i^{jk}$  to corner  $_i^{ki}$ , we use

$$f_j^{ki} - f_i^{jk} = \frac{1}{2} (\mu_{ij}^k + R_{-\zeta_{ij}} \mu_{ji}^l),$$
(7)

where  $\zeta_{ij} := \frac{\pi}{2} [(\beta - \pi/4)/(\pi/2)]$  is the closest quarter-rotation taking the frame in triangle *ijk* to the frame in triangle *jil*. In other words, we account for the jump in the frame across the cut. Likewise, on the other side of the edge, with endpoints at corners  $_{i}^{lj}$  and  $_{j}^{il}$ , we have

$$f_j^{il} - f_i^{lj} = \frac{1}{2} (R_{\zeta_{ij}} \mu_{ij}^k + \mu_{ji}^l).$$
(8)

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# E Pseudocode

Here we give detailed pseudocode for our rectangular parameterization algorithm. We omit pre-processing (namely, generation of the reference field) and post-processing (namely, quantization and contouring of the parameterization in the case of meshing), since these steps are handled by existing algorithms. All other methods are given explicitly in terms of basic operations on vectors and a standard halfedge mesh data structure, and should be easily translatable to real code. The only subroutines not spelled out here are:

- UNIT(v) given a nonzero vector  $v \in \mathbb{R}^3$ , returns the unit vector in the same direction.
- OBJECTIVE( $M, A, u, v, \theta$ ), GRADOBJECTIVE( $M, A, u, v, \theta$ ), HESSOBJECTIVE( $M, A, u, v, \theta$ ) – implements one of the objective functions from Section 5.1.1, and its gradient/Hessian, given the mesh connectivity M, the triangle areas A, log scale factors u, v per vertex, frame rotation angles  $\theta$  per face.
- SOLVE(A, b) returns x solution of Ax = b
- SOLVEQP(H, g, J, F) solves a quadratic program of the form  $\min_y \frac{1}{2}y^T Hy + y^T g$  subject to Jy + F = 0. (We use *OSQP* [Stellato et al. 2020].)
- SPARSEFROMTRIPLETS((I, J, S), n, m) returns a sparse matrix A of size  $n \times m$  whose coefficient  $A_{I_i J_i}$  is equal to  $S_i$  for all i. If a coefficient appears several times in the lists (I, J), their corresponding values in S are summed.
- APPEND(*L*, *a*) returns a new list appending the element *a* to the list *L*.
- ANGLE(x, y, z) returns the angle between the vector  $y x \in \mathbb{R}^3$  and  $z x \in \mathbb{R}^3$ .

For simplicity we assume M is without boundary.

# Algorithm 1 FLATTENMESH(M, x, W)

- **Input:** A manifold, orientable, simply-connected triangle mesh M = (V, E, F) with corners *C*, coordinates  $x_i \in \mathbb{R}^3$  at each vertex  $i \in V$ , and a cross field given by any one of four representative unit vectors  $W_{ijk} \in \mathbb{R}^3$  in each triangle  $ijk \in F$
- **Output:** Coordinates  $f : C \to \mathbb{R}^2$  describing the flattening of each triangle into the 2D plane. Note that these values are stored at triangle corners rather than vertices to accommodate possible cuts.

1:	for each $ij \in E$ do $\ell_{ij} \leftarrow  x_i - x_j $	<i>⊳edge lengths</i>	
2:	for each $j_i^{jk} \in C$ do $\alpha_i^{jk} \leftarrow \text{Angle}(x_i, x_i)$	j, x <sub>k</sub> ) ⊳corner angles	
3:	for each $ijk \in F$ do	⊳cross field angles	
	▷ compute a basis $(N, T_1, T_2)$ for triang	le ijk adapted to edge ij	
4:	$N_1 \leftarrow (x_j - x_i) \times (x_k - x_i)$	<i>▶unnormalized normal</i>	
5:	$A_{ijk} \leftarrow  N_1 /2$	⊳triangle area	
6:	$N_1 \leftarrow N_1/(2A_{ijk})$	⊳normalize	
7:	$T_1 \leftarrow \text{UNIT}(x_j - x_i)$		
8:	$T_2 \leftarrow N_1 \times T_1$		
9:	$\xi_{ijk} \leftarrow \operatorname{atan2}(\langle T_2, W_{ijk} \rangle, \langle T_1, W_{ijk} \rangle)$	⊳angle relative to ij	
10:	$\Gamma, \zeta, s, \varphi, \omega^0 \leftarrow \text{ComputeCutJumpData}(M, \alpha, \xi)$		
11:	$u, v, \theta \leftarrow \text{SolveOptimizationProblem}(M, A, \alpha, \varphi, \omega^0, s)$		
12:	$f \leftarrow \text{RecoverParameterization}(M, \Gamma, \zeta, \ell, \alpha, \varphi, \theta, u, v)$		
13:	return <i>f</i>		

### Algorithm 2 COMPUTECUTJUMPDATA( $M, \alpha, \xi$ )

- **Input:** A manifold, orientable triangle mesh M = (V, E, F) with angles  $\alpha_i^{jk}$  at each corner *i* of each triangle  $ijk \in F$ , angles  $\xi : F \to \mathbb{R}$  describing one of the four directions of the cross field in each triangle  $ijk \in F$  as a rotation relative to the direction of the first edge ij.
- **Output:** The curve  $\Gamma : E \to \{0, 1\}$  cutting *M* to a topological disk, the closest quarter-rotation  $\zeta : E \to \mathbb{R}$  between two frame in adjacent triangle, sign bits  $s : C \to \pm 1$  giving the sign of *v* at each corner relative to the values stored at vertices, angles  $\varphi : H \to \mathbb{R}$  which give the reference frame relative to each halfedge, and angles  $\omega^0 : E \to \mathbb{R}$  describing the rotation of the reference frame  $X^0$  along each oriented dual edge.
- 1: for each ijk ∈ F do ⊳ for convenience, express cross directions relative to halfedges
- $$\begin{split} \xi_{\vec{i}\vec{j}} &\leftarrow \xi_{ijk} \\ \xi_{\vec{j}\vec{k}} &\leftarrow \xi_{\vec{i}\vec{j}} (\pi \alpha_j^{ki}) \\ \xi_{\vec{k}\vec{i}} &\leftarrow \xi_{\vec{j}\vec{k}} (\pi \alpha_k^{ij}) \end{split}$$
  3: 4: ⊳Perform breadth-first traversal (Section 5.3.1) 5:  $ijk \leftarrow First(F)$ ▶ get the first triangle in the mesh 6: for each  $ijk \in F$  do  $\varphi_{ijk} \leftarrow \infty$ ▷ frame not yet set 7:  $\varphi_{ijk} \leftarrow \xi_{ijk}$  > reference frame gets one of the four cross directions 8:  $PUSH(Q, \overline{ij})$ ▶ push any halfedge of ijk onto a queue 9: while NOTEMPTY(Q) do ⊳go from ijk to jil 10:  $ij \leftarrow \operatorname{Pop}(Q)$ *⊳parallel transport*  $\varphi_{ij \to ji} \leftarrow \varphi_{ij} + \pi$ 11:  $n^{*} \leftarrow \operatorname{MoD}(\operatorname{Round}(2\frac{\varphi_{ij} \rightarrow ji - \xi_{ji}}{\pi}), 4) \qquad \triangleright closest \ cross \ index \\ \zeta_{ij} \leftarrow \frac{\pi}{2}n^{*} \qquad \triangleright jump \ angle \ across \ ij \\ \xi^{*} \leftarrow \zeta_{ij} + \xi_{ji} \qquad \triangleright closest \ cross \\ \omega_{ij}^{0} \leftarrow \varphi_{ij} - \xi^{*} + \pi \quad \triangleright smallest \ rotation \ to \ neighboring \ cross \\ if \ \varphi_{ji} == \infty \ then \qquad \triangleright neighboring \ frame \ not \ yet \ set \\ \Gamma \leftarrow c \ 0$ 12: 13: 14:
- 15: 16: ▶ primal edge not in cut  $\Gamma_{ij} \leftarrow 0$ 17:  $\begin{aligned} r_{ij} \leftarrow 0 \\ s_{ij} \leftarrow +1 \\ \varphi_{ji} \leftarrow \xi^* \\ \varphi_{j\vec{k}} \leftarrow \varphi_{i\vec{j}} - (\pi - \alpha_j^{ki}) \\ \varphi_{\vec{k}i} \leftarrow \varphi_{j\vec{k}} - (\pi - \alpha_k^{ij}) \\ \text{PUSH}(Q, i\vec{l}) \end{aligned}$ ⊳no sign change 18: ▶ set reference frame in triangle 19: 20: 21: ▶visit other two neighbors 22:  $PUSH(Q, \overline{lj})$ 23: ▷neighboring frame already set else 24:  $\Gamma_{ij} = 0$ ⊳primal edge could be in cut 25: if  $IsODD(n^*)$  then  $s_{ii} = -1$ 26: **else**  $s_{ij} = +1$ 27:

# 28: end while

```
Compute relative signs at corners

29: for each <math>i \in V do

30: S = 1 >cumulative product (Equation 20)

31: for p \leftarrow 0, \dots, DEGREE(M, i) - 1 do

32: s_i^{j_p j_{p+1}} \leftarrow S

33: S \leftarrow s_{i j_{p+1}} S >compute angle defect
```

34:  $K_i \leftarrow 0$ 35: for each  $ijk \in F$  do  $K_i \leftarrow K_i + \alpha_i^{jk}$  $K_j \leftarrow K_j + \alpha_j^{ki}$ 36: 37:  $K_k \leftarrow K_k + \alpha_k^{ij}$ 38: 39:  $K \leftarrow 2\pi - K$ *⊳compute cones indices at vertices* 40:  $c_i \leftarrow K$ 41: for each  $i \in V$  do **for each**  $j \in V$  such that  $ij \in E$  **do** *▶edge incident to i* 42:  $w \leftarrow \omega_{ij}^0$ 43 if i > j then  $w \leftarrow -w$ 44:  $c_i \leftarrow c_i + w$ ▷accumulate all rotations 45: *▶ compute the cut* 46:  $r \leftarrow \text{True}$ 47: while r do  $d \leftarrow 0$ *⊳initialize degree of unvisited edges* 48 for each  $\Gamma_{ii} == 1$  do *▶iterate over unvisited edges* 49:  $d_i \leftarrow d_i + 1$ 50:  $d_i \leftarrow d_i + 1$ 51:  $r \leftarrow \text{False}$ 52: for each  $i \in V$  do *▶iterate over vertices* 53 **if**  $d_i == 1$  and  $c_i == 0$  **then**  $\triangleright$  *if* degree-1 and not a cone 54: for each  $ij \in E$  do  $\Gamma_{ij} = 0$ ▷remove edge 55:  $r \leftarrow \text{True}$ 56: 57: end while 58: return ( $\Gamma$ ,  $\zeta$ , s,  $\varphi$ ,  $\omega^0$ )

# **Algorithm 3** SolveOptimizationProblem( $M, A, \alpha, \varphi, \omega^0, s$ )

**Input:** A manifold, orientable, simply-connected triangle mesh *M* = (V, E, F), areas  $A_{ijk}$  for each triangle  $ijk \in F$ , angles  $\alpha_i^{jk}$  at each corner  $j_i^k \in C$ , angles  $\varphi_{ij}$  describing the angle of the reference frame relative to each halfedge  $\overrightarrow{ij} \in H$ , angles  $\omega_{ii}^0$ describing the rotation of the reference frame across each edge  $ij \in E$ , and values  $s_i^{jk} \in \{\pm 1\}$  giving the sign of the value at each corner  $\frac{jk}{i}$  relative to a value stored at vertex *i*. **Output:** Scale factors  $u_i, v_i \in \mathbb{R}$  at each vertex  $i \in V$ , and angles  $\theta_{iik} \in \mathbb{R}$  for each triangle  $ijk \in F$ . 1:  $u, v \leftarrow 0 \in \mathbb{R}^{|V|}$ *▶initialize scale factors* 2:  $\theta \leftarrow 0 \in \mathbb{R}^{|F|}$ *▶initialize frame angles* 3:  $\lambda \leftarrow 0 \in \mathbb{R}^{|E|}$ ▷initialize Lagrange multiplier 4: while TRUE do  $H \leftarrow \text{HessObjective}(M, A, u, v, \theta)$ 5:  $q \leftarrow \text{gradObjective}(M, A, u, v, \theta)$ 6:  $D \leftarrow \text{BuildHessian}(M, \alpha, u, v, \lambda, s, \varphi, \theta)$ 7: *⊳constraint*  $J \leftarrow \text{BuildJacobian}(M, \alpha, u, v, s, \varphi, \theta) \triangleright constraint \exists a cobian$ 8:  $F \leftarrow \text{BuildSystem}(M, \alpha, u, v, s, \varphi, \theta, \omega^0) \triangleright constraint value$ 9:  $\mathbf{r} \leftarrow (\mathbf{u} \ \mathbf{v} \ \mathbf{\theta} \ \lambda)$ 

12:	$b \leftarrow \begin{vmatrix} g + J^{\top} \lambda \\ F \end{vmatrix}$	right hand side
13:	$y \leftarrow \tilde{\text{Solve}}(A, b)$	
14:	$\tau \leftarrow \text{LineSearch}(M, \alpha, s, \varphi, \omega^0, J, F, g, x, y)$	
15:	$(u, v, \theta, \lambda) \leftarrow x + \tau y$	⊳update guess
16:	<b>if</b> $  g + J^{\top}\lambda   +   F   < \varepsilon$ <b>then</b>	
17:	break	
18:	end while	
19:	return $(u, v, \theta)$	

### **Algorithm 4** LINESEARCH( $M, \alpha, s, \varphi, \omega^0, J, F, q, x, y$ )

**Input:** A manifold, orientable triangle mesh M = (V, E, F) with angles  $\alpha_i^{jk}$  at each corner *i* of each triangle  $ijk \in F$ , sign bits  $s: C \to \pm 1$  giving the sign of v at each corner relative to the values stored at vertices, angles  $\varphi : H \to \mathbb{R}$  giving the angle of the reference frame relative to each halfedge, angles  $\omega^0$ :  $E \to \mathbb{R}$  describing the rotation of the reference frame  $X^0$ across each edge, a sparse matrix  $J \in \mathbb{R}^{|E| \times (2|V| + |F|)}$  giving the Jacobian of the constraint, values  $F : E \to \mathbb{R}$  giving the discrete residual of Equation 17, a vector  $g \in \mathbb{R}^{2|V| + |F|}$ giving the gradient of the objective function, two vectors  $x, y \in \mathbb{R}^{2|V| + |F|}$  describing the current estimate and the descent direction

**Output:** Step size  $\tau > 0$  decreasing the Lagrangian gradient norm 1:  $(u, v, \theta, \lambda) \leftarrow x$ 

2:	$\mathcal{E} \leftarrow \ g + J^{\top}\lambda\  + \ F\ $	⊳norm of gradient Lagrangian
3:	$\tau \leftarrow 1$	⊳step size
4:	while True do	
5:	$(u, v, \theta, \lambda) \leftarrow x + \tau y$	<i>⊳update estimate</i>
6:	$g \leftarrow \text{gradObjective}(M, A, u)$	$(u, v, \theta) $ > gradient of objective
7:	$J \leftarrow \text{BuildJacobian}(M, \alpha, u)$	$, v, s, \varphi, \theta) \triangleright constraint Jacobian$
8:	$F \leftarrow \text{BuildSystem}(M, \alpha, u, z)$	$(\varphi, s, \varphi, \theta, \omega^0) \triangleright constraint value$
9:	$\mathcal{E}_{\text{new}} \leftarrow \ g + J^{\top}\lambda\  + \ F\ $	<i>▶update gradient norm</i>
10:	if $\mathcal{E}_{\text{new}} \leq \left(1 - \frac{\tau}{2}\right) \mathcal{E}$ then	
11:	break	▶ stop when sufficient decrease
12:	else	
13:	$\tau \leftarrow 0.9\tau$	⊳decrease step size
14:	end while	
15:	return $\tau$	

**Algorithm 5** BUILDSYSTEM( $M, \alpha, u, v, s, \varphi, \theta, \omega^0$ )

**Input:** A manifold, orientable triangle mesh M = (V, E, F) with angles  $\alpha_i^{jk}$  at each corner *i* of each triangle  $ijk \in F$ , values  $u, v : V \to \mathbb{R}$  at each vertex giving the log scale factors, sign bits  $s : C \to \pm 1$  giving the sign of v at each corner relative to the values stored at vertices, angles  $\varphi : H \to \mathbb{R}$  giving the angle of the reference frame relative to each halfedge, angles  $\theta : F \to \mathbb{R}$  describing the current frame  $X_1^{\theta}$  as a rotation of the reference frame, angles  $\omega^0 : E \to \mathbb{R}$  describing the rotation of the reference frame  $X^0$  across each edge.

**Output:** Values  $\rho : E \to \mathbb{R}$  giving the residual of our discrete equation (Equation 17) on each edge.

1: **for each**  $ij \in E$  **do** 2:  $\rho_{ij} \leftarrow (\theta_{ijk} - \theta_{jil}) - \omega_{ii}^0$   $\triangleright$  left-hand side of Equation 10  $\triangleright$ k and l are left/right of ij

- 3: for each  $ijk \in F$  do  $\triangleright$  right-hand side of Equation 10  $\triangleright$  add contribution to residual from dual edges inside this triangle
- 4:  $\rho_{ij} \leftarrow \rho_{ij} + \text{Residual}(i, j, k, \alpha, u, v, s, \varphi_{ij} + \theta_{ijk})$
- 5:  $\rho_{jk} \leftarrow \rho_{jk} + \text{Residual}(j, k, i, \alpha, u, v, s, \varphi_{jk} + \theta_{ijk})$
- 6:  $\rho_{ki} \leftarrow \rho_{ki} + \text{Residual}(k, i, j, \alpha, u, v, s, \varphi_{ki} + \theta_{ijk})$

7: return  $\rho$ 

### Algorithm 6 RESIDUAL $(i, j, k, \alpha, u, v, s, \eta)$

- **Input:** Three vertices *i*, *j*, *k* of a triangle, the angles  $\alpha$  at triangle corners, the log scale factors u, v at vertices, the signs s at triangle corners, and the angle  $\eta$  of the current frame  $X_1^{\theta}$ relative to a coordinate system aligned with edge ij.
- Output: The contribution to the residual of equation Equation 17 due to the piece of the edge dual to *ij* within triangle *ijk*.

1: $(v_i^{jk}, v_j^{ki}, v_k^{ij}) \leftarrow (s_i^{jk}v_i, s_j^{ki}v_j, s_k^{ij})$	v <sub>k</sub> ) ⊳values at corners
2: $\rho \leftarrow u_j - u_i$	⊳conformal part
3: $\rho \leftarrow \rho + \cos(2\eta)(v_j^{ki} - v_i^{jk})$	⊳non-conformal part
4: $\rho \leftarrow \rho + \sin(2\eta) \cot \alpha_j^{ki} (v_k^{ij} - v_i^{jk})$	<sup>2</sup> )
5: $\rho \leftarrow \rho + \sin(2\eta) \cot \alpha_i^{jk} (v_k^{ij} - v_j^{k})$	<sup>i</sup> )
6: $\rho \leftarrow \frac{1}{2} \cot \alpha_k^{ij} \rho$	
7: <b>if</b> $i > j$ <b>then</b>	
8: $\rho \leftarrow -\rho$	<i>▶account for edge orientation</i>
9: return $\rho$	

**Algorithm 7** BUILDJACOBIAN( $M, \alpha, u, v, s, \varphi, \theta$ )

- **Input:** A manifold, orientable triangle mesh M = (V, E, F) with angles  $\alpha_i^{j\kappa}$  at each corner *i* of each triangle  $ijk \in F$ , values  $u, v : V \to \mathbb{R}$  at each vertex giving the log scale factors, signs  $s : C \to \pm 1$  describing the values of v at corners relative to the values stored at vertices, angles  $\varphi : F \to \mathbb{R}$ describing the reference frame  $X_1^0$  in each triangle  $ijk \in F$ as a rotation relative to the direction of the first edge *ij*, and angles  $\theta$  :  $F \to \mathbb{R}$  describing the current frame  $X_1^{\theta}$  as a rotation of the reference frame.
- **Output:** A sparse matrix  $J \in \mathbb{R}^{|E| \times (2|V|+|F|)}$  giving the Jacobian of the constraint equations computed by BUILDSYSTEM.
- 1:  $L_v \leftarrow ()$  $\triangleright$  list of nonzeros of Jacobian with respect to v2:  $L_{\theta} \leftarrow ()$  $\triangleright$  list of nonzeros of Jacobian with respect to  $\theta$
- 3: for each  $ij \in E$  do  $\triangleright$  Jacobian of left-hand side of Equation 10  $L_{\theta} \leftarrow \operatorname{Append}(L_{\theta}, (ij, ijk, +1))$ 4:
- $L_{\theta} \leftarrow \operatorname{Append}(L_{\theta}, (ij, jil, -1))$ 5:
- 6: for each  $ijk \in F$  do  $\triangleright$  Jacobian of right-hand side of Equation 10
- $\eta_{ij} \leftarrow \varphi_{ijk} + \theta_{ijk} \qquad \triangleright angle \ of X_1^{\theta} \ relative \ to \ each \ edge$  $n_{ij} \leftarrow n_{ii} (\pi \alpha^{ki})$ 7:

8: 
$$\eta_{jk} \leftarrow \eta_{ij} - (\pi - \alpha_{j_i}^{\kappa})$$

- 9:  $\eta_{ki} \leftarrow \eta_{jk} (\pi \alpha_k^{ij})$ 10:  $\triangleright$  append contribution to Jacobian from dual edges in this triangle  $(L_u, L_v, L_\theta) \leftarrow \text{Jacobian}(i, j, k, \alpha, u, v, s, \eta_{ij}, L_u, L_v, L_\theta)$ 11:
- $(L_u, L_v, L_\theta) \leftarrow \text{Jacobian}(j, k, i, \alpha, u, v, s, \eta_{jk}, L_u, L_v, L_\theta)$ 12:
- $(L_u, L_v, L_\theta) \leftarrow \text{Jacobian}(k, i, j, \alpha, u, v, s, \eta_{ki}, L_u, L_v, L_\theta)$ 13:
- Build Jacobian blocks from lists of nonzeros 14:
- 15:  $J_u \leftarrow \text{SparseFromTriplets}(L_u, |E|, |V|)$

- 16:  $J_{v} \leftarrow \text{SparseFromTriplets}(L_{v}, |E|, |V|)$
- 17:  $J_{\theta} \leftarrow \text{SparseFromTriplets}(L_{\theta}, |E|, |F|)$

18: return  $[J_u J_v J_{\theta}]$  >assemble overall  $|E| \times (2|V| + |F|)$  Jacobian

**Algorithm 8** JACOBIAN $(i, j, k, \alpha, u, v, s, \eta, L_u, L_v, L_\theta)$ 

Input: The same inputs as RESIDUAL, plus the cumulative lists of entries  $L_u, L_v, L_\theta$  of the three Jacobian blocks built so far.

**Output:** Modified lists  $L_u, L_v, L_\theta$  that account for the derivatives (with respect to u, v and  $\theta$ ) of the residual due to the piece of the edge dual to *ij* within triangle *ijk*. 1:  $w \leftarrow \frac{1}{2} \cot \alpha_{l}^{ij}$ *⊳*cotan weight 2: if i > j then  $w \leftarrow -w$ *▶account for edge orientation* 3: *⊳conformal part w.r.t. u* 4 5:  $L_u \leftarrow \text{Append}(L_u, (ij, i, -w))$ 6:  $L_u \leftarrow \text{Append}(L_u, (ij, j, +w))$ ⊳non-conformal part w.r.t ט 8:  $L_{\upsilon} \leftarrow \operatorname{Append}(L_{\upsilon}, (ij, i, -ws_i^{jk}(\cos(2\eta) + \sin(2\eta)\cot\alpha_i^{ki})))$ 9:  $L_{\upsilon} \leftarrow \operatorname{Append}(L_{\upsilon}, (ij, j, +ws_j^{ki}(\cos(2\eta) - \sin(2\eta)\cot\alpha_i^{jk})))$ 10:  $L_{\upsilon} \leftarrow \operatorname{Append}(L_{\upsilon}, (ij, k, +ws_k^{ij} \sin(2\eta)(\cot \alpha_j^{ki} + \cot \alpha_i^{jk})))$ 11:  $\triangleright non-conformal part w.r.t \theta$ 12:  $(v_i^{jk}, v_i^{ki}, v_k^{ij}) \leftarrow (s_i^{jk}v_i, s_i^{ki}v_j, s_k^{ij}v_k)$ ▷values at corners 13:  $L_{\theta} \leftarrow \text{Append}(L_{\theta}, (ij, ijk, -2w \sin(2\eta)(v_j^{ki} - v_i^{jk})))$ 14:  $L_{\theta} \leftarrow \operatorname{Append}(L_{\theta}, (ij, ijk, +2w\cos(2\eta)\cot \alpha_{j}^{ki}(v_{k}^{ij} - v_{i}^{jk})))$ 15:  $L_{\theta} \leftarrow \text{Append}(L_{\theta}, (ij, ijk, +2w\cos(2\eta)\cot\alpha_{i}^{jk}(v_{k}^{ij} - v_{i}^{ki})))$ 16: return  $(L_u, L_v, L_\theta)$ 

**Algorithm 9** BUILDHESSIAN( $M, \alpha, u, v, \lambda, s, \varphi, \theta$ )

- **Input:** A manifold, orientable triangle mesh M = (V, E, F) with angles  $\alpha_i^{jk}$  at each corner *i* of each triangle  $ijk \in F$ , values  $u, v : V \to \mathbb{R}$  at each vertex giving the log scale factors, values  $\lambda : E \to \mathbb{R}$  at each dual-edge giving the Lagrange multiplier associated to the constraint, signs  $s : C \rightarrow \pm 1$ describing the values of v at corners relative to the values stored at vertices, angles  $\varphi : F \to \mathbb{R}$  describing the reference frame  $X_1^0$  in each triangle  $ijk \in F$  as a rotation relative to the direction of the first edge ij, and angles  $\theta: F \rightarrow \mathbb{R}$  describing the current frame  $X_1^{\theta}$  as a rotation of the reference frame.
- **Output:** A sparse matrix  $H \in \mathbb{R}^{(2|V|+|F|) \times (2|V|+|F|)}$  giving the Jacobian of the matrix computed by JACOBIAN multiplied by the Lagrange multiplier  $\lambda$ .

1: 
$$L_u \leftarrow ()$$
  $\triangleright$  list of nonzeros of Jacobian with respect to u

- 2:  $L_{v} \leftarrow ()$  $\triangleright$  list of nonzeros of Jacobian with respect to v
- 3:  $L_{\theta} \leftarrow ()$  $\triangleright$  list of nonzeros of Jacobian with respect to  $\theta$
- 4: for each  $ijk \in F$  do > Jacobian of right-hand side of Equation 10

- 5:  $\eta_{ij} \leftarrow \varphi_{ijk} + \theta_{ijk}$  > angle of  $X_1^{\theta}$  relative to each edge 6:  $\eta_{jk} \leftarrow \eta_{ij} (\pi \alpha_k^{ji})$ 7:  $\eta_{ki} \leftarrow \eta_{jk} (\pi \alpha_k^{ij})$ 8: > append contribution to Jacobian from dual edges in this triangle
- $(L_{v}, L_{\theta}) \leftarrow \text{Hessian}(i, j, k, \alpha, u, v, \lambda, s, \eta_{ij}, L_{v}, L_{\theta})$ 9:
- $(L_{\upsilon}, L_{\theta}) \leftarrow \text{Hessian}(j, k, i, \alpha, u, \upsilon, \lambda, s, \eta_{ik}, L_{\upsilon}, L_{\theta})$ 10:

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# **Algorithm 10** HESSIAN $(i, j, k, \alpha, u, v, \lambda, s, \eta, L_v, L_\theta)$

**Input:** The same inputs as BUILDHESSIAN, plus the cumulative lists of entries  $L_v, L_\theta$  of the two Hessian blocks built so far.

**Output:** Modified lists  $L_v$ ,  $L_\theta$  that account for the derivatives (with respect to v and  $\theta$ ) of the Jacobian due to the piece of the edge dual to *ij* within triangle *ijk*.

1:  $w \leftarrow \frac{1}{2} \cot \alpha_{i}^{ij}$ *⊳*cotan weight 2: if i > j then  $w \leftarrow -w$ *▶account for edge orientation* 3: ⊳non-conformal part w.r.t v 4: 5:  $L_{\upsilon} \leftarrow \operatorname{Append}(L_{\upsilon}, (i, ijk, 2s_i^{jk}\lambda_{ij}w(\sin(2\eta) - \cos(2\eta)\cot\alpha_j^{ki})))$ 6:  $L_{\upsilon} \leftarrow \operatorname{Append}(L_{\upsilon}, (j, ijk, 2s_j^{ki}\lambda_{ij}w(-\sin(2\eta) - \cos(2\eta)\cot\alpha_i^{jk})))$ 7:  $L_{\upsilon} \leftarrow \operatorname{Append}(L_{\upsilon}, (k, ijk, 2s_k^{ij}\lambda_{ij}w\cos(2\eta)(\cot\alpha_j^{ki} + \cot\alpha_i^{jk})))$ 8:  $\triangleright non-conformal \ part \ w.r.t \ \theta$ 9:  $(v_i^{jk}, v_i^{ki}, v_k^{ij}) \leftarrow (s_i^{jk}v_i, s_i^{ki}v_j, s_k^{ij}v_k)$ ⊳values at corners 10:  $L_{\theta} \leftarrow \operatorname{Append}(L_{\theta}, (ijk, ijk, -4\lambda_{ij}w\cos(2\eta)(v_i^{ki} - v_i^{jk})))$ 11:  $L_{\theta} \leftarrow \operatorname{Append}(L_{\theta}, (ijk, ijk, -4\lambda_{ij}w\sin(2\eta)\cot\alpha_{i}^{ki}(v_{k}^{ij} - v_{i}^{jk})))$ 12:  $L_{\theta} \leftarrow \operatorname{Append}(L_{\theta}, (ijk, ijk, -4\lambda_{ij}w\sin(2\eta)\cot\alpha_i^{jk}(v_k^{ij} - v_i^{ki})))$ 13: return  $(L_v, L_\theta)$ 

### **Algorithm 11** RecoverParameterization( $M, \Gamma, \zeta, \ell, \alpha, \varphi, \theta, u, v$ )

**Input:** The mesh connectivity M = (V, E, F), the cut curve  $\Gamma : E \rightarrow \{0, 1\}$  cutting the mesh to disk topology, the closest quarterrotation  $\zeta : E \rightarrow \mathbb{R}$  between two frame in adjacent triangle, the length  $\ell_{ij} \in \mathbb{R}_{\geq 0}$  of each edge  $ij \in E$ , angles  $\alpha_i^{jk}$  at each triangle corner, the angle  $\varphi_{ijk} \in \mathbb{R}$  of the reference frame relative to edge ij in each triangle ijk, the angle  $\theta_{ijk}$  of the optimized frame relative to the reference frame, and the log scale/stretch factors  $u_i, v_i$  at each vertex  $i \in V$ .

**Output:**  $f : C \to \mathbb{R}^2$  giving the uv-coordinates per triangle corner. 1: for each  $ij \in E$  do  $\triangleright$  edge scale from log-scale

2:  $a_{ii} \leftarrow e^{(u_i + u_j + v_i + v_j)/2}$ 

3: 
$$b_{ij} \leftarrow e^{(u_i + u_j - \upsilon_i - \upsilon_j)/2}$$

```
4: for each ijk \in F do
```

5:  $\eta_{ij} \leftarrow \varphi_{ijk} + \theta_{ijk}$  > angle of  $X_1^{\theta}$  relative to each edge 6:  $\eta_{jk} \leftarrow \eta_{ij} - (\pi - \alpha_{j_{ij}}^{ki})$ 7:  $\eta_{ki} \leftarrow \eta_{jk} - (\pi - \alpha_{k}^{j})$ 8: > target edge vector 9:  $\mu_{ij}^k \leftarrow \ell_{ij}(+a_{ij}\cos(\eta_{ijk}), +b_{ij}\sin(\eta_{ijk}))$ 10:  $\mu_{jk}^i \leftarrow \ell_{jk}(-a_{jk}\cos(\eta_{ijk} + \alpha_{j}^{ki}), -b_{jk}\sin(\eta_{ijk} + \alpha_{j}^{ki}))$ 

11: 
$$\mu_{ki}^{j} \leftarrow \ell_{ki}(-a_{ki}\cos(\eta_{ijk} - \alpha_{i}^{jk}), -b_{ki}\sin(\eta_{ijk} - \alpha_{i}^{jk}))$$

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12: *▶right hand-side in Equation ??* 13: for each  $_{k}^{ij} \in C$  do  $b_{k}^{ij} \leftarrow \frac{1}{2}(R_{\zeta_{ij}}\mu_{ij}^{k} + \mu_{ij}^{l})$ 14: 15:  $L_A \leftarrow ()$ 16:  $L_U \leftarrow ()$ 17:  $p \leftarrow 0$ 18: for each  $_{k}^{ij} \in C$  do 19 *▶ build matrix A in Equation 6*  $\begin{array}{l} L_A \leftarrow \operatorname{Append}(L_A, (^{ij}_k, ^{jk}_k, +1)) \\ L_A \leftarrow \operatorname{Append}(L_A, (^{ij}_k, ^{ki}_j, -1)) \end{array}$ 20: 21: 22. *▶ build matrix U in Equation 6* **if**  $\Gamma_{ij} == 0$  **then**  $\triangleright$  *if ij not in the cut, add corner constraints* 23:  $L_U \leftarrow \operatorname{Append}(L_U, (p, \frac{jk}{i}, +1))$ 24:  $L_U \leftarrow \operatorname{Append}(L_U, (p, \overset{l_j}{i}, -1))$ 25:  $\begin{array}{l} L_U \leftarrow \operatorname{Append}(L_U, (p+1, \overset{i}{j}{}^i, +1)) \\ L_U \leftarrow \operatorname{Append}(L_U, (p+1, \overset{i}{j}{}^i, -1)) \end{array}$ 26: 27:  $p \leftarrow p + 2$ 28 Build Jacobian blocks from lists of nonzeros 29:  $A \leftarrow \text{SparseFromTriplets}(L_A, |C|, |C|)$ 30:  $U \leftarrow \text{SparseFromTriplets}(L_U, p, |C|)$ 

31:  $f \leftarrow \text{SolveQP}(A^\top WA, -A^\top Wb, U, 0) \triangleright \text{corner coordinates}$