

Basics of Database Security

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“Basics” of Database Security

- Technical and Organizational point of view:
 1. Protecting data against
 1. Unauthorized access
 2. Unauthorized operations
 3. Watching the database activities (Auditing)
 2. Operational Security: tend toward service continuity:
 1. Replication
 2. Recovery

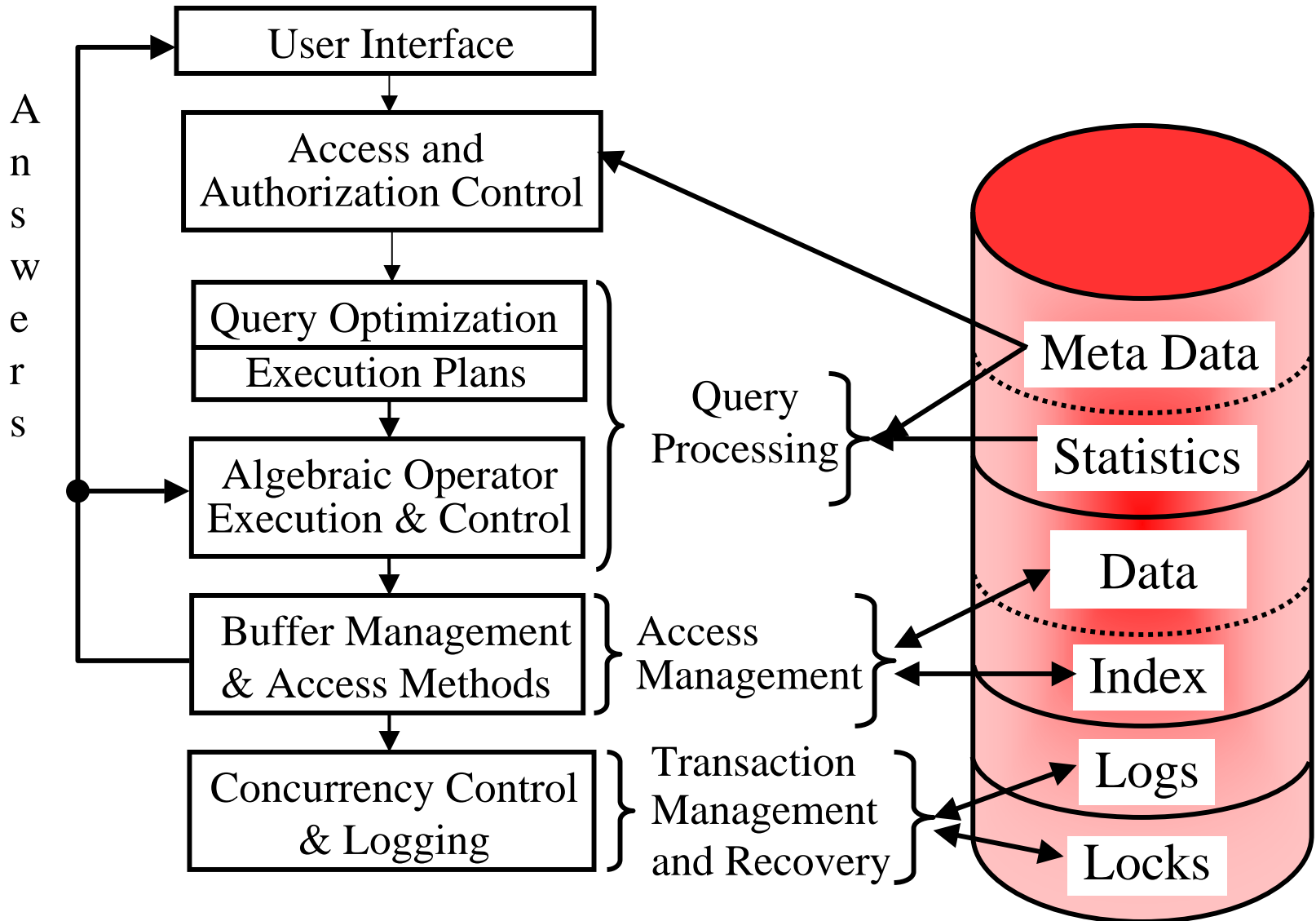
“Basics” of Database Security: Outline

- I. Preamble:
 - i. DBMS
 - ii. DB applications in C/S Architectures
- II. Protecting Database Access
- III. Operational Security and Recovery
- IV. Operational Security Thanks to Replication
- V. Database Auditing
- VI. Concluding Remarks

Database Management Systems Features

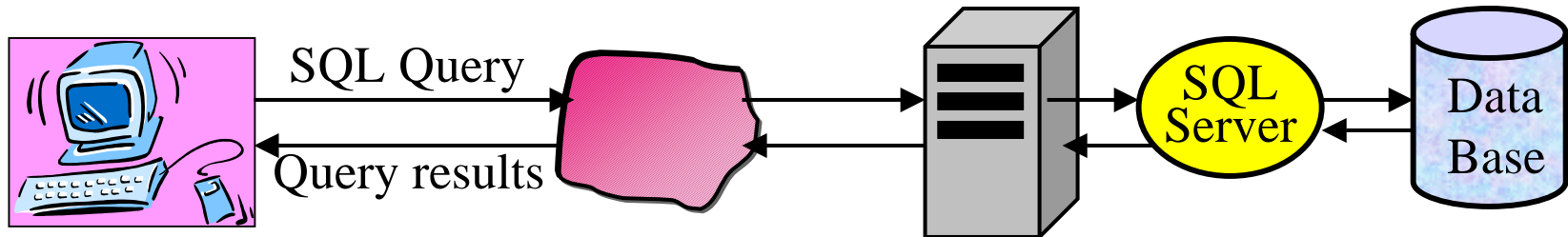
1. Data Description (“objects”, attributes, relationships, constraints)
2. (High Level, Set-Oriented) Data Manipulation
3. Consistency and Privacy
4. Concurrency Control
5. Security (and Recovery from failure)

Database Management System Architecture

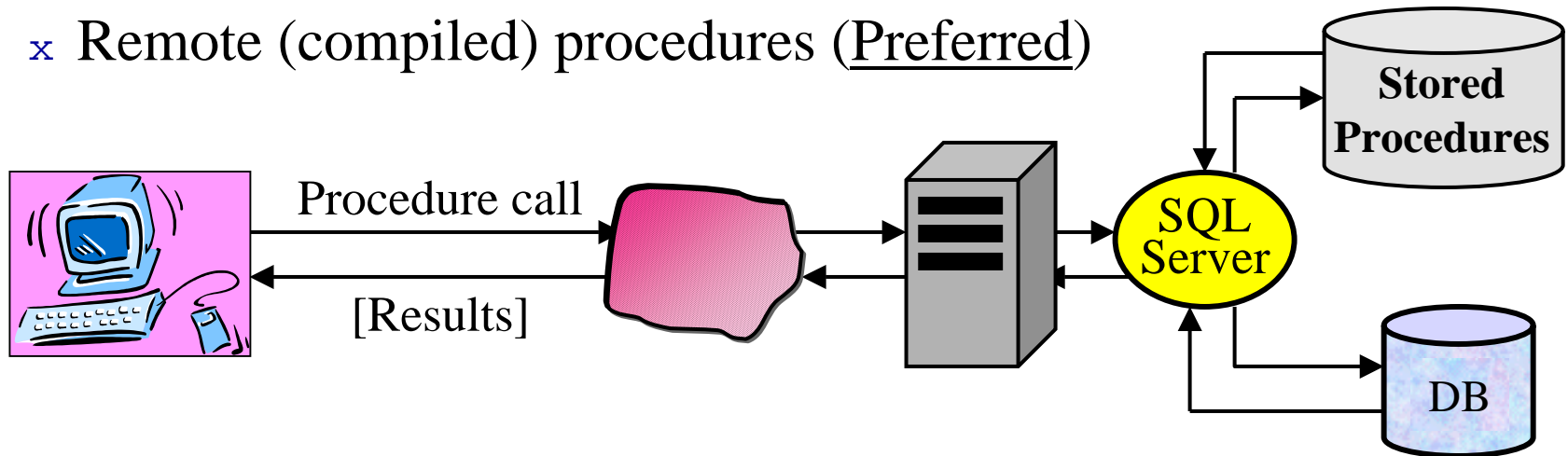


○ Accessing Data/Object Servers

x SQL Queries

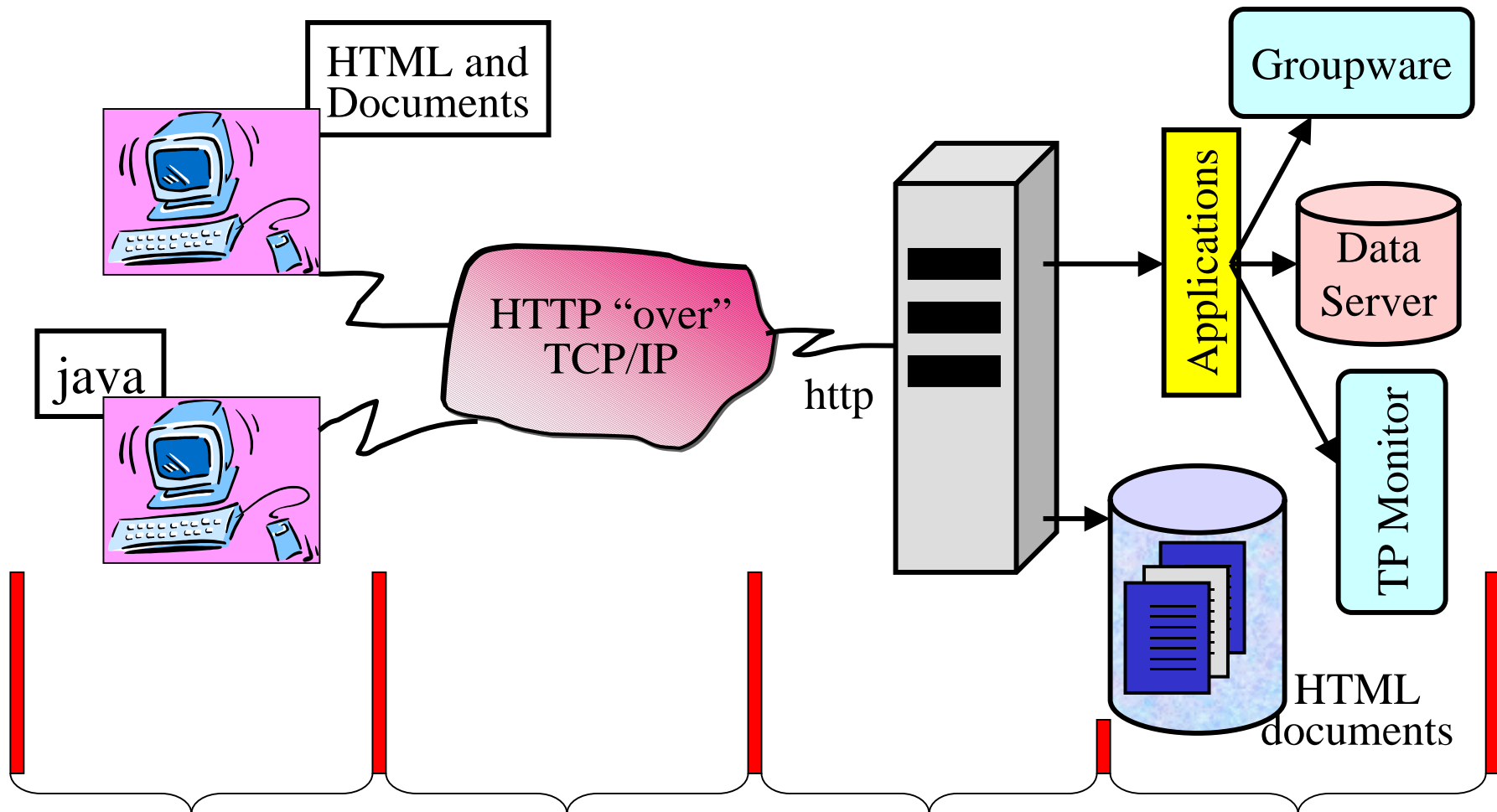


x Remote (compiled) procedures (Preferred)



DBMSs in C/S Architectures

○ Web Servers and Data/Object Servers



DBMSs in C/S Architectures: Middleware

- Middleware : the Client-Server “glue”
- General Middleware includes:
 - x Communication stacks
 - x Authentication services
 - x Remote procedure calls
 - x Message queues
 - x Distributed naming services
 - x etc.

DBMSs in C/S Architectures: Middleware

- Specific Middleware:
 - × Database: ODBC, JDBC, DRDA, EDA/SQL
 - × Transaction: TxRPC, Transactional RPC
 - × Distributed Objects: CORBA, OLE/DCOM
 - × Internet: HTTP, S(ecure)-HTTP

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Data Integrity and Privacy

- Privacy:
 - x Encryption, Delusions
 - x Login, Password,
 - x Views
 - x Priviledges, Roles and Profiles

- Identification Levels
 1. Operating/Networking System
 2. DBMS
 3. [Database(s)]

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Operational Security and Recovery

- Ensure fast recovery from
 - × Program failure
 - × DBMS failure
 - × Operating System failure
 - × etc.
- Tightly Coupled with Concurrency Control
- Technical as well as Organizational Problem

Concurrency Control & Transaction Management

- Transaction: Atomic sequence of actions
- ACID Transaction Properties:
 - x Atomicity
 - x Consistency
 - x Isolation
 - x Durability (Skip details ?)
- Recovery: Transaction logs (Transaction history)
 - x Undo: “Play” the transaction backward (Rollback)
 - x Redo: “Play” the transaction forward

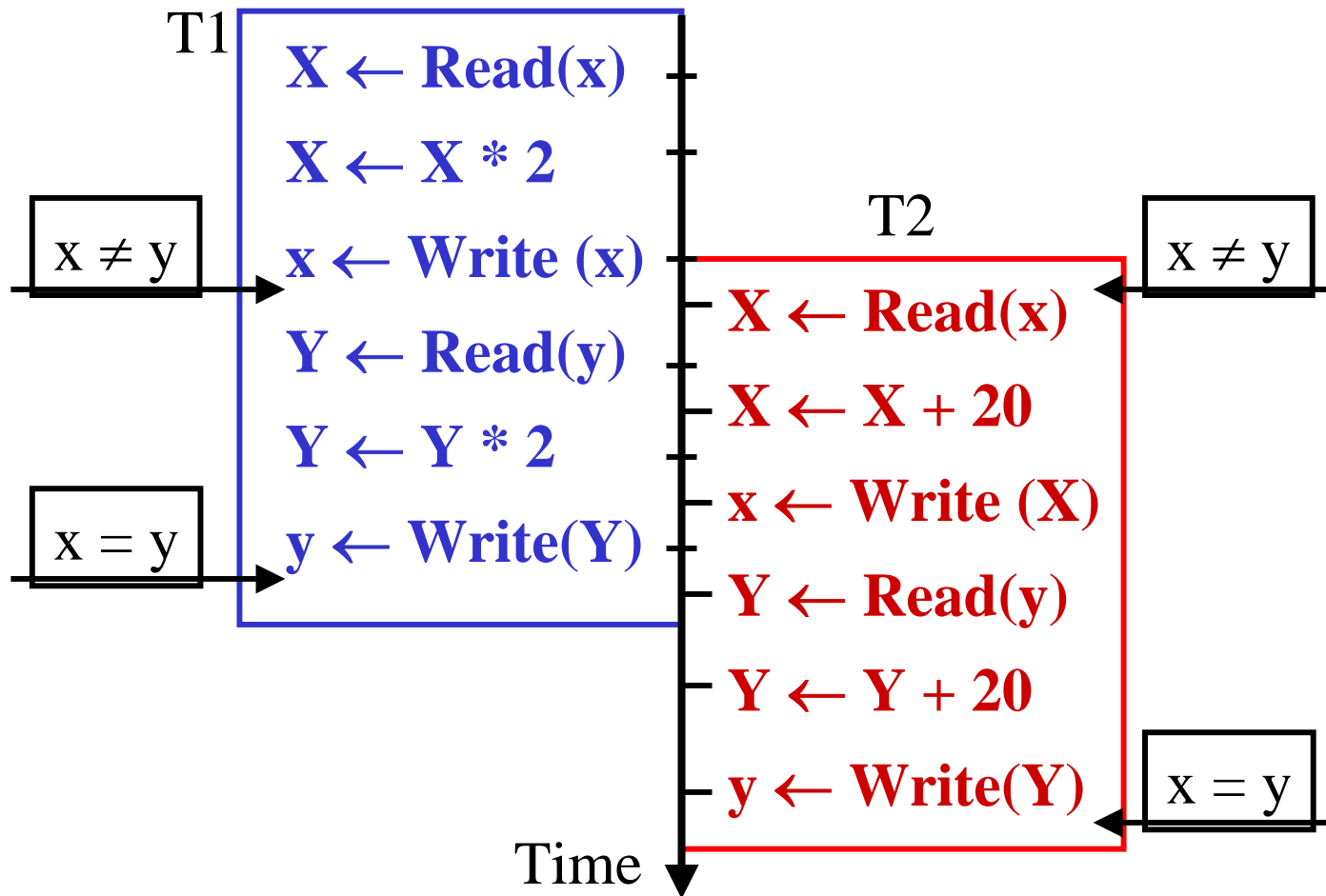
Transaction Management (cont'd): Isolation Levels

- **Atomicity**: Money transfer from C1 to C2
 - × Withdraw from C1 ($C1 \leftarrow C1 - 100$)
 - × Deposit into C2 ($C2 \leftarrow C2 + 100$)
 - × Problem after Withdrawal and before Deposit

Transaction	Account C1	C2
Read(C1) ←	3100	5000
$C1 \leftarrow C1 - 100$		
Write C1 →	3000	
Read(C2) ←		
$C2 \leftarrow C2 + 100$		
Write(C2) →		5100

Transactions (cont'd): Isolation Levels

- Consistency: Constraint ($x = y$)



Transactions (cont'd): Isolation Levels

- Isolation: Withdrawal and Deposit (same account)
- T1: $\{ X \leftarrow \text{Read}(C); X \leftarrow X - 200; C \leftarrow \text{Write}(X) \}$
- T2: $\{ Y \leftarrow \text{Read}(C); Y \leftarrow Y + 600; C \leftarrow \text{Write}(Y) \}$
- Concurrent execution of T1 and T2 (Dirty Reads)
 - × **T1**: $X \leftarrow \text{Read}(C): X \leftarrow 5000$
 - × **T1**: $X \leftarrow X - 200: X \leftarrow 4800$
 - × T2: $Y \leftarrow \text{Read}(C): Y \leftarrow 5000$
 - × T2: $Y \leftarrow Y + 600: Y \leftarrow 5600$
 - × T2: $C \leftarrow \text{Write}(Y): C \leftarrow 5600$
 - × **T1**: $C \leftarrow \text{Write}(X): C \leftarrow 4800$ (*instead of 5400*)

Transaction (cont'd): Isolation Levels

- Durability: Permanent database updates

1. (Non) Repeatable Reads

x C in Database = 5000

x **T1**: Read(C) → X: X ← 5000

x T2: Read(C) → Y: Y ← **5000**

x **T1**: X ← X + 5000: X ← 10 000

x **T1**: Write(X) → C: C ← 10 000

x T2: Read(C) → Y: Y ← **10 000**

- Different values !!!

Transaction (cont'd): Isolation Levels

- **Durability:** Permanent database updates (cont'd)

2. (Non) Shadow tuples

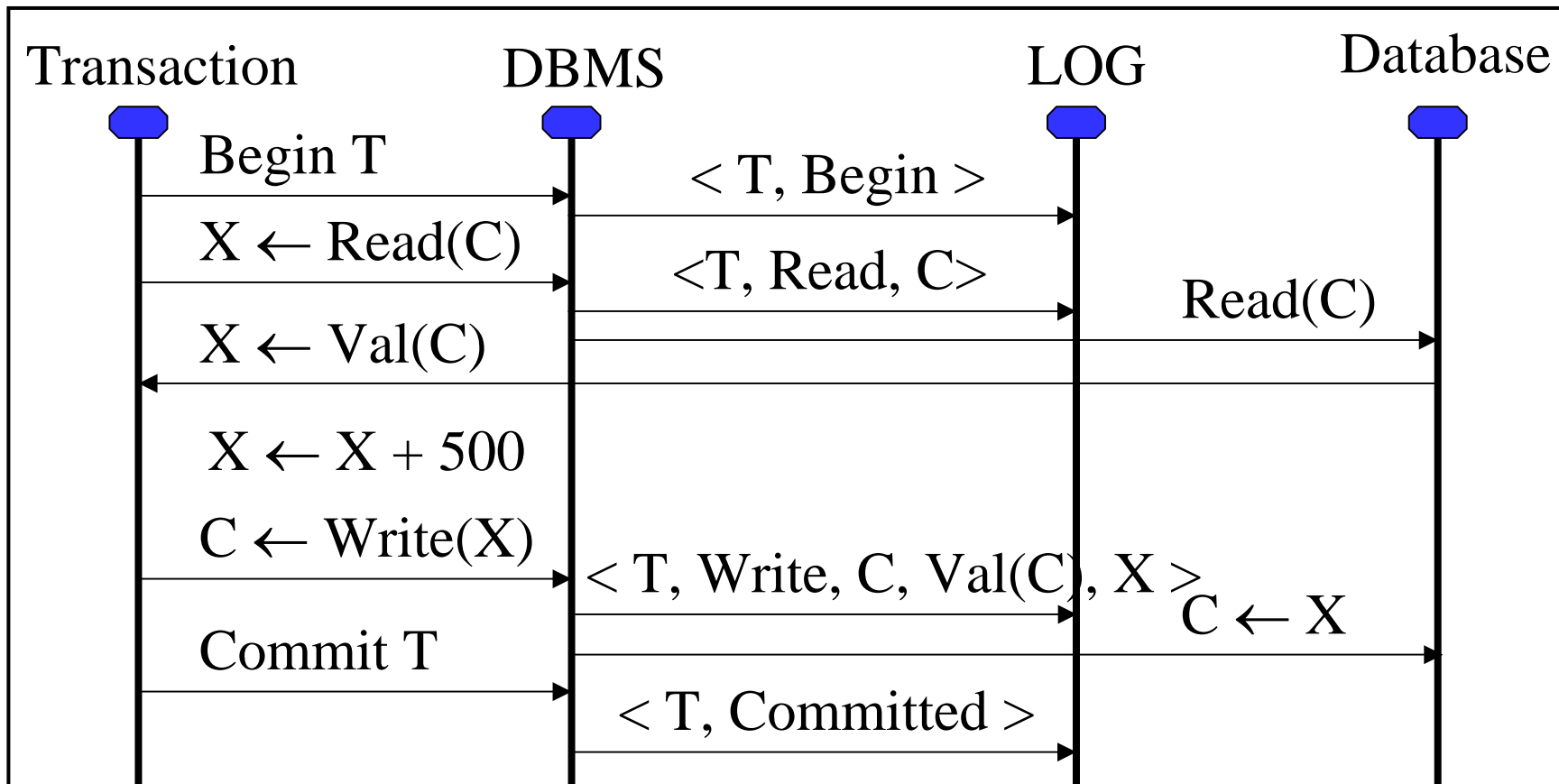
- x **T1:** Read(X) such that $X = 500$
 - x T2: Read(X) such that $X = 500$
 - x **T1:** Delete(X)
 - x T2: Read(X) such that $X = 500$
- No X , such that $X = 500$, exists anymore

Transaction (cont'd): Isolation Levels (end)

- Introduced in SQL'92
- Level 0: Dirty reads permitted
 - × While a data item D is being updated by $T1$
 - × Any other transaction can read D but it CANNOT modify it
- Level 1: Dirty reads not allowed
- Level 2: Reads repeatability not allowed before commit
- Level 3: Inhibits shadow tuples

Transaction Logging and Recovery

- Write Ahead Logging Protocol (WAL)
- Rollback: 'Execute' Transaction Log backward



Transaction Management and Recovery (cont'd)

- Concurrency Control and Serialisability:
Concurrent execution is equivalent to sequential (serial) execution
- Foundations: partial ordering of actions
- Techniques
 - x Time stamping (not usually implemented)
 - x Two-Phase Locking (ensures Serialisability)
 - 7 Phase1: Acquire Locks
 - 7 Phase2: Free Locks
 - 7 Does not prevent from deadlocks or livelocks

Transaction Management and Recovery (cont'd)

- Most Usual Types of locks:
 - × Shared: multiple reads and no update
 - × **E**Xclusive: single update
 - × Intention locks: **R**ead, **W**rite, **U**ppdate

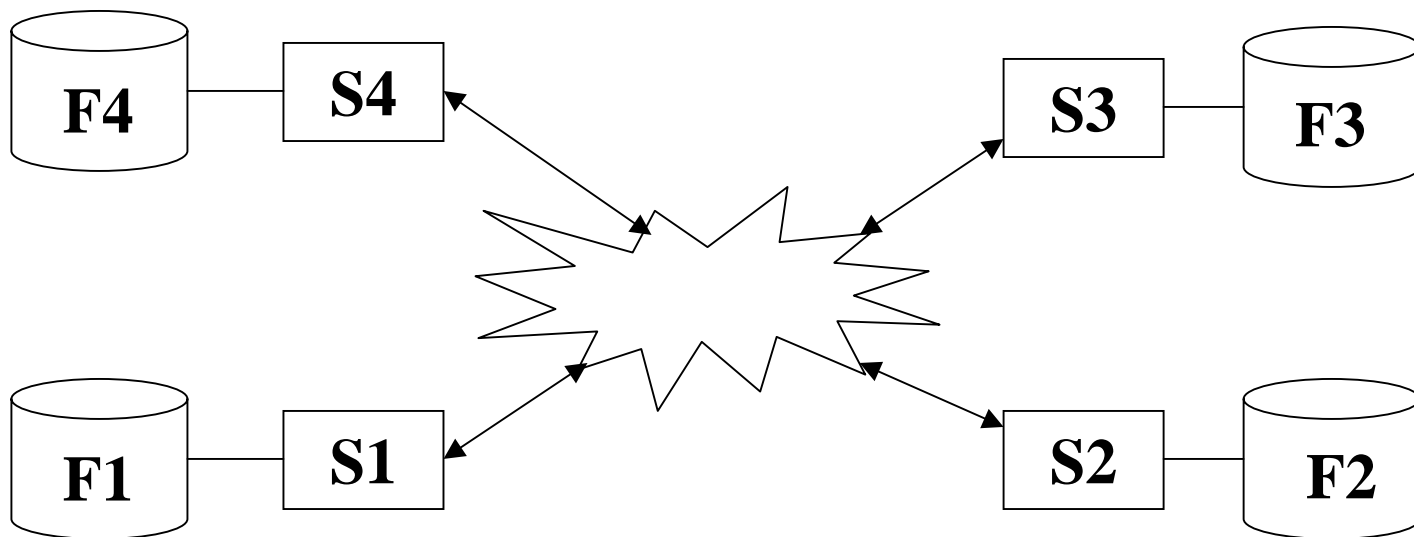
Lock mode		S	X
Requested Lock Mode	S	Yes	No
	X	No	No

	R	U	W
R	Yes	No	No
U	Yes	No	No
W	No	No	No

- cf. Select ... for update/at isolation level serializable/holdlock

Note: Distributed Transaction Management

- Distributed DB:
 - × DB = « Union » of « Fragments »

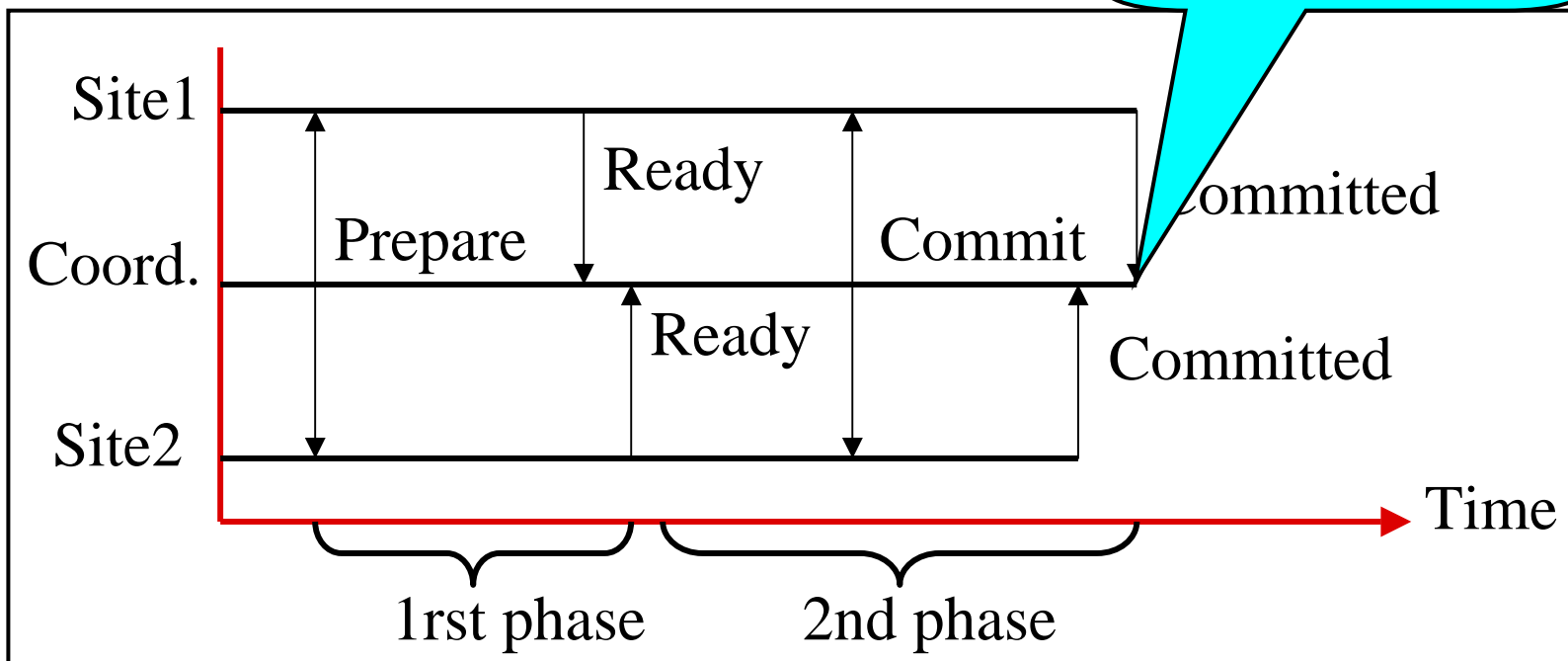


Note: Distributed Transaction Management

- Context:
 - x Multiple sites cooperate in (sub-)transactions
 - x Commit/Rollback on every site
- *Two-Phase Commit Protocol*
 - x Lampson [1976], Gray [1978]
 - x Most commonly implemented
 - x Ensures transaction Atomicity [Baer & al., ICSE81]
 - x Requires a coordination site

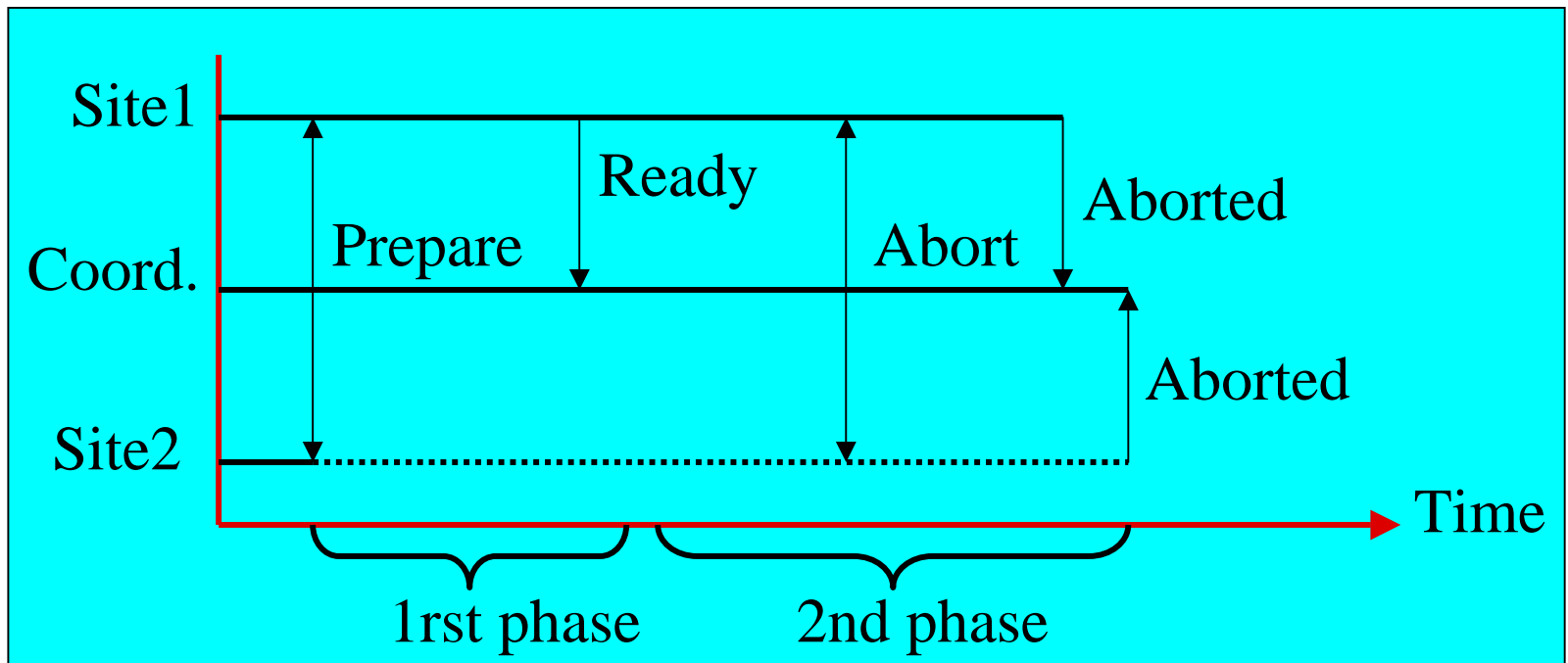
Distributed Transaction Management (cont'd)

- Coordination site:
 - × Decomposes a transaction into sub-transactions depending on data location (cf. Query processing)
 - × Sends the sub-transactions



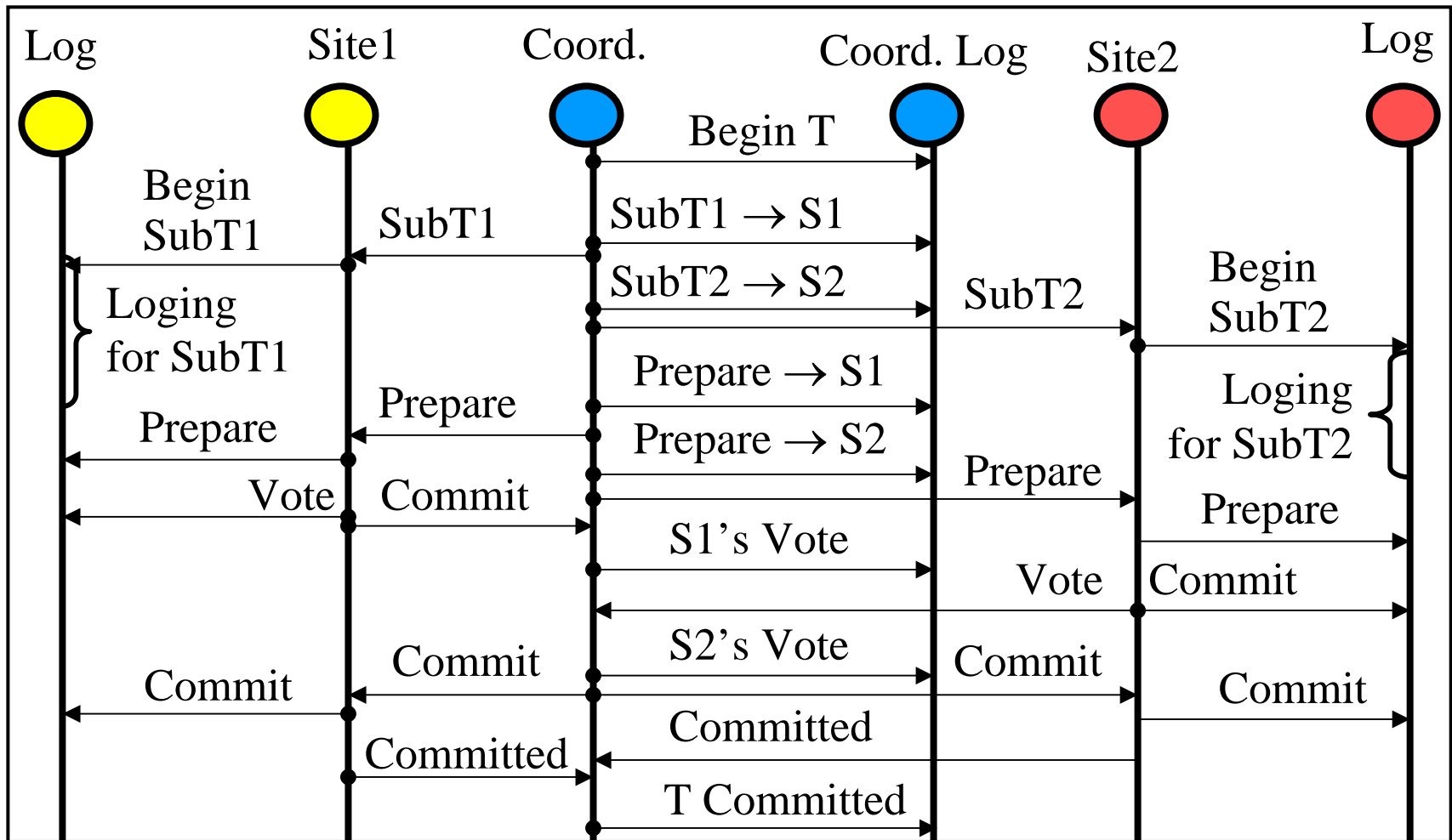
Distributed Transaction Management (cont'd)

- Transaction failure:
 - ✗ One or more sites cannot commit
 - ✗ One or more answers to the prepare message are missing...



Distributed Transaction Management (cont'd)

Transaction Logs: (Usually) WAL Protocol

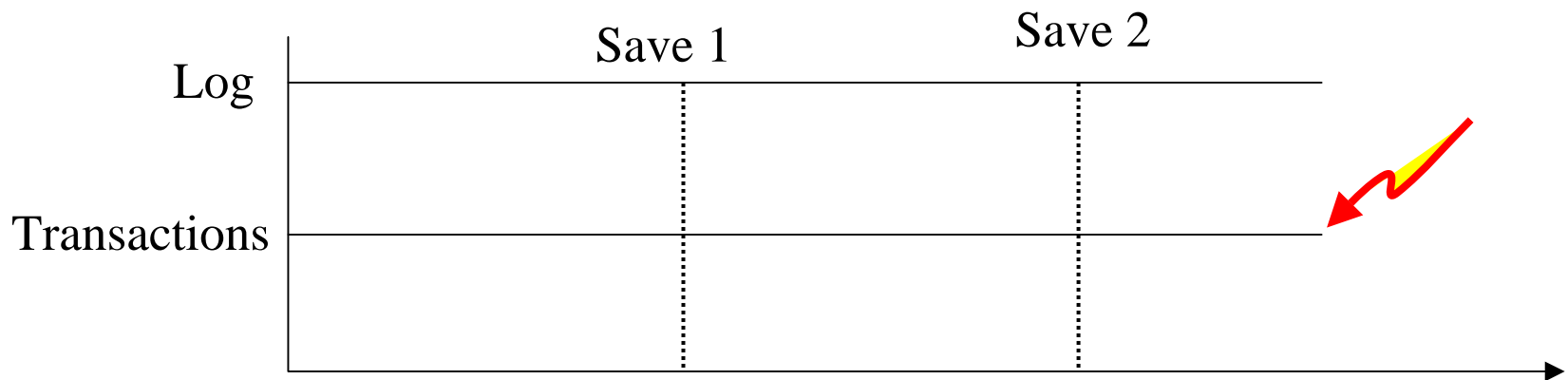


Transaction Management (end)

- Serialisability: Serial execution equivalent to concurrent execution
- ACID properties
- Mechanisms:
 - x Concurrency Manager
 - x Integrity Manager
 - x Recovery Manager
 - x Programmers
 - x Security Officer [+ DB Administrator]

Log Management and Recovery

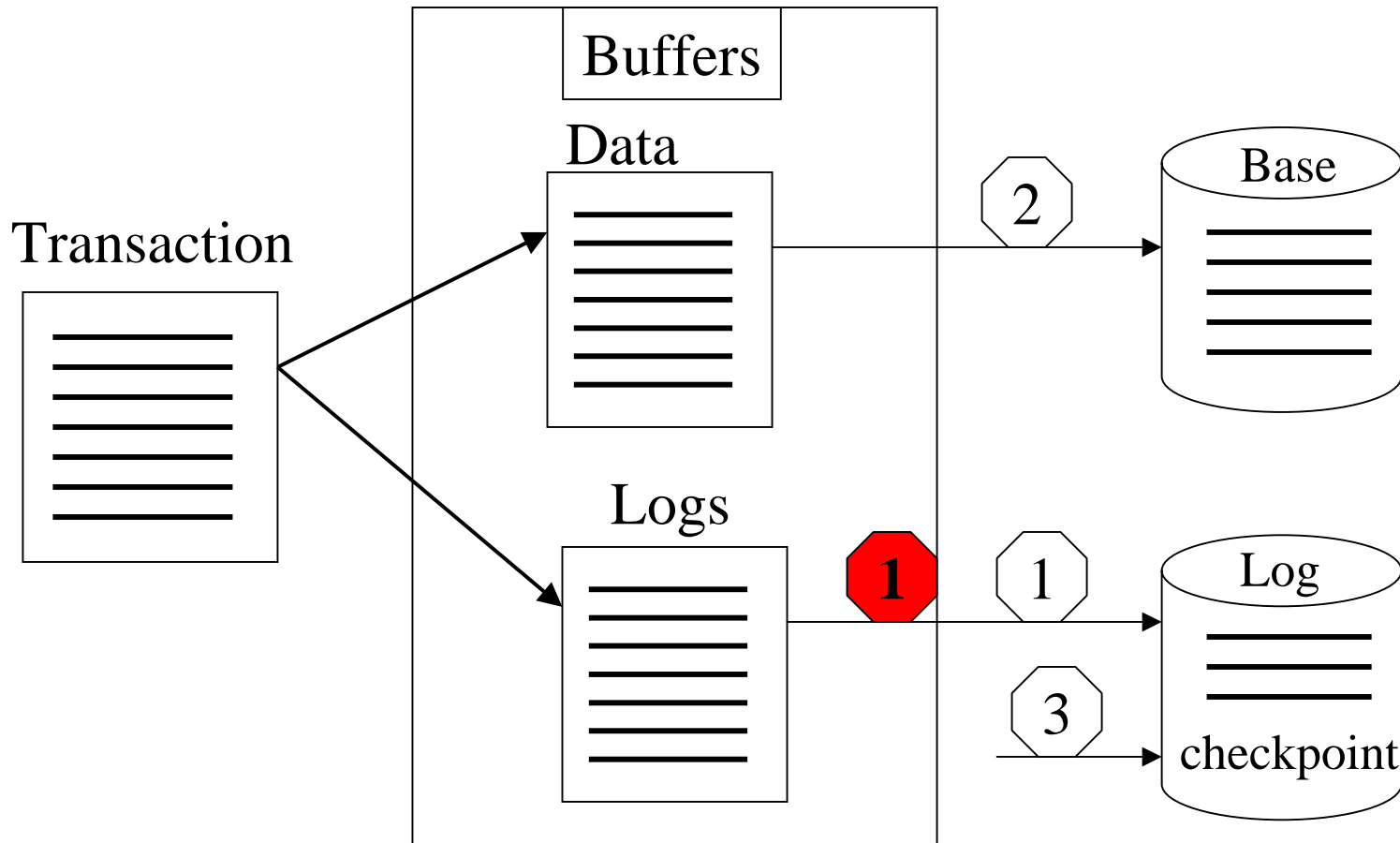
- Log management:
 - × Configuration
 - 7 Database creation
 - 7 During database « life »
 - × Periodic Save
 - × [Restore, when needed]



Log Management and Recovery

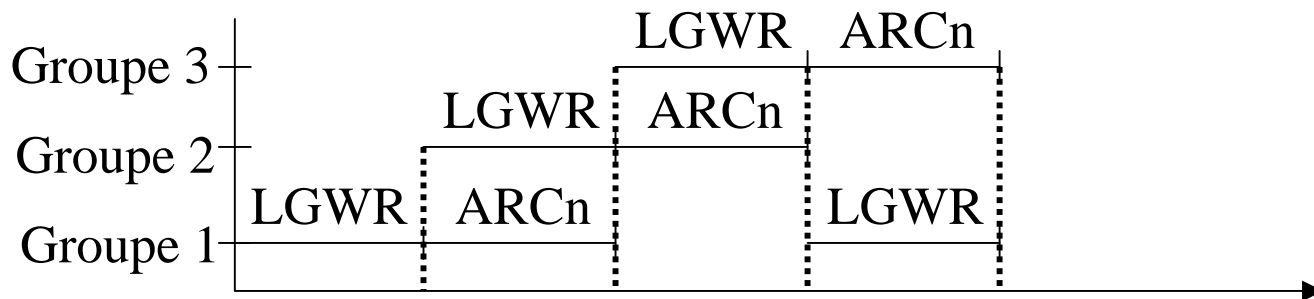
○ Commit vs Checkpoint

x Data Server configuration (*checkpoint frequency*)



Log Management and Recovery in Oracle

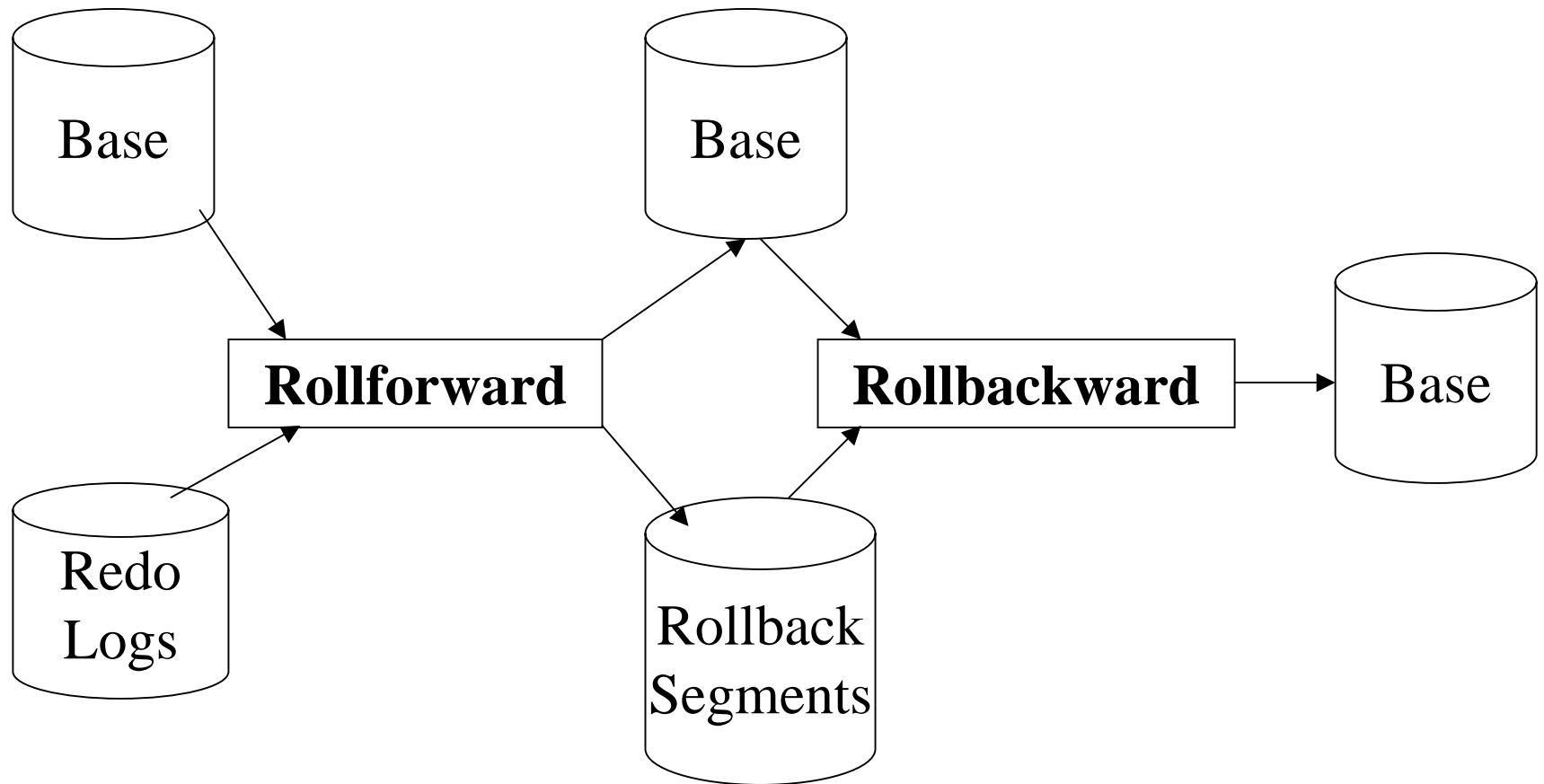
- Checkpoint: configuration parameters
 - x *Log-checkpoint_interval* (number of blocks)
 - x *Log_checkpoint_timeout* (in seconds)
- Logging:
 1. Redo logs
 2. Rollback segments/Undo Tablespaces
- Automatic Archive Logs



- Back-Up and Recovery
 - × Total/Partial

- Recovery
 - × Rollforward (*cache recovery*)
 - × Rollbackward (*transaction recovery*)

Log Management and Recovery in Oracle



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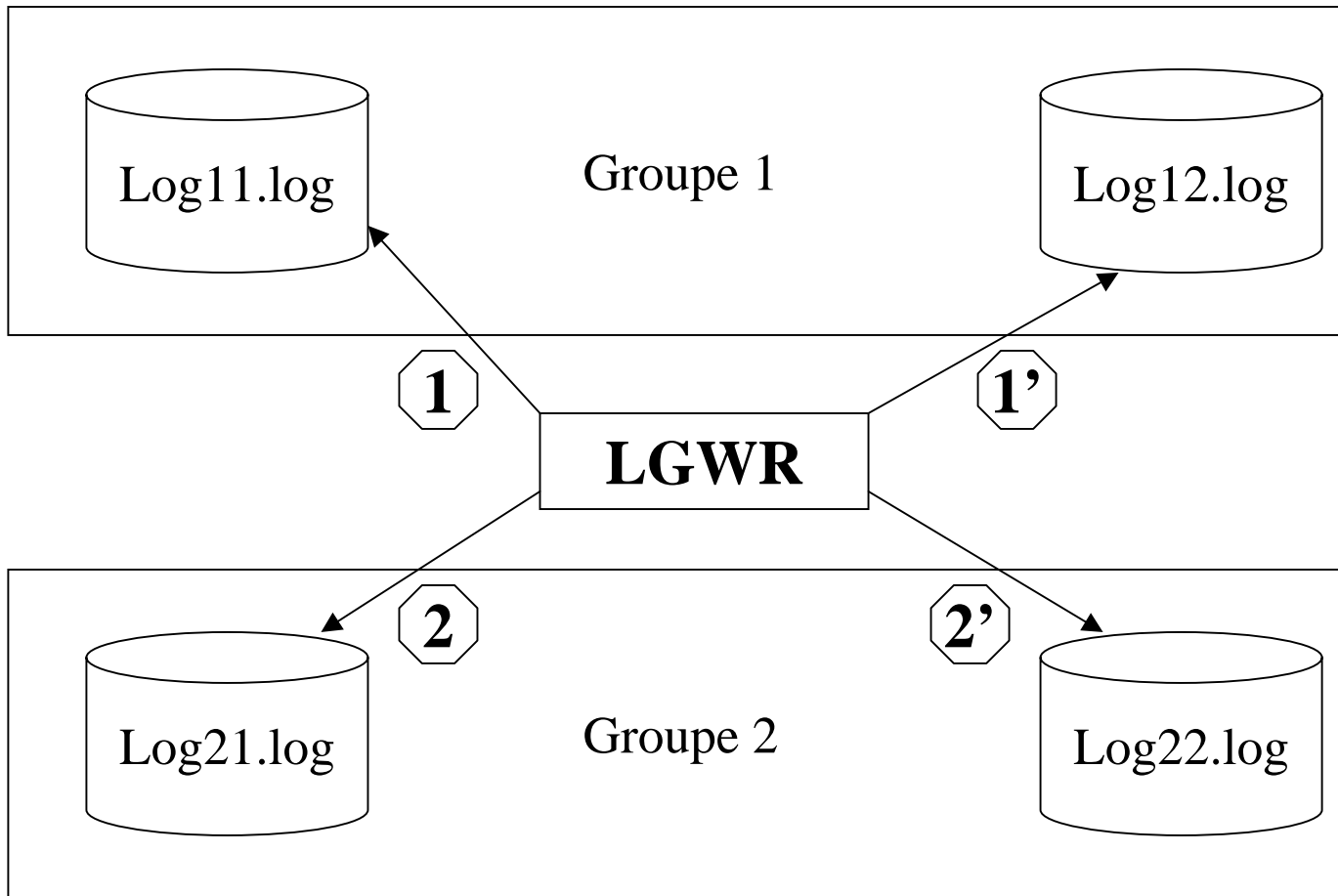
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5. Database Auditing
6. Concluding Remarks

Security thanks to Replication

- Synchronize the content of a storage unit with its replicate
- Switch in case of failure
- Replicate
 - × At least the logs
 - × And/or sensitive data
- Requires more space
- Activity overhead (may conflict performance issues)

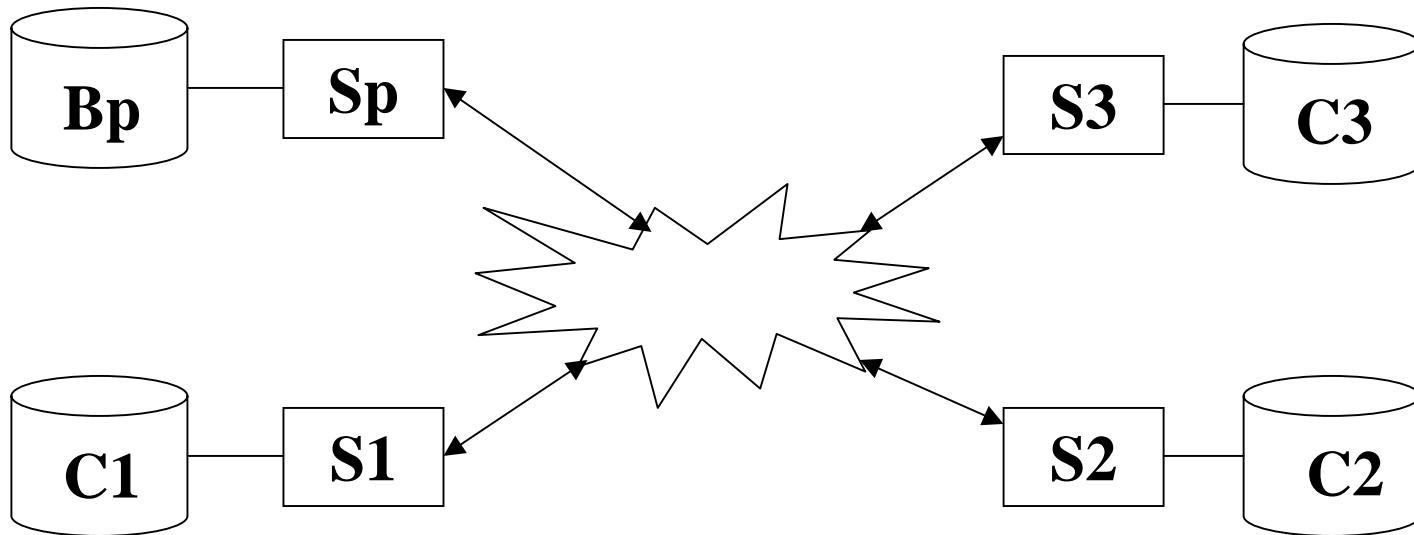
Security thanks to Replication

- Example: Oracle Multiplexing



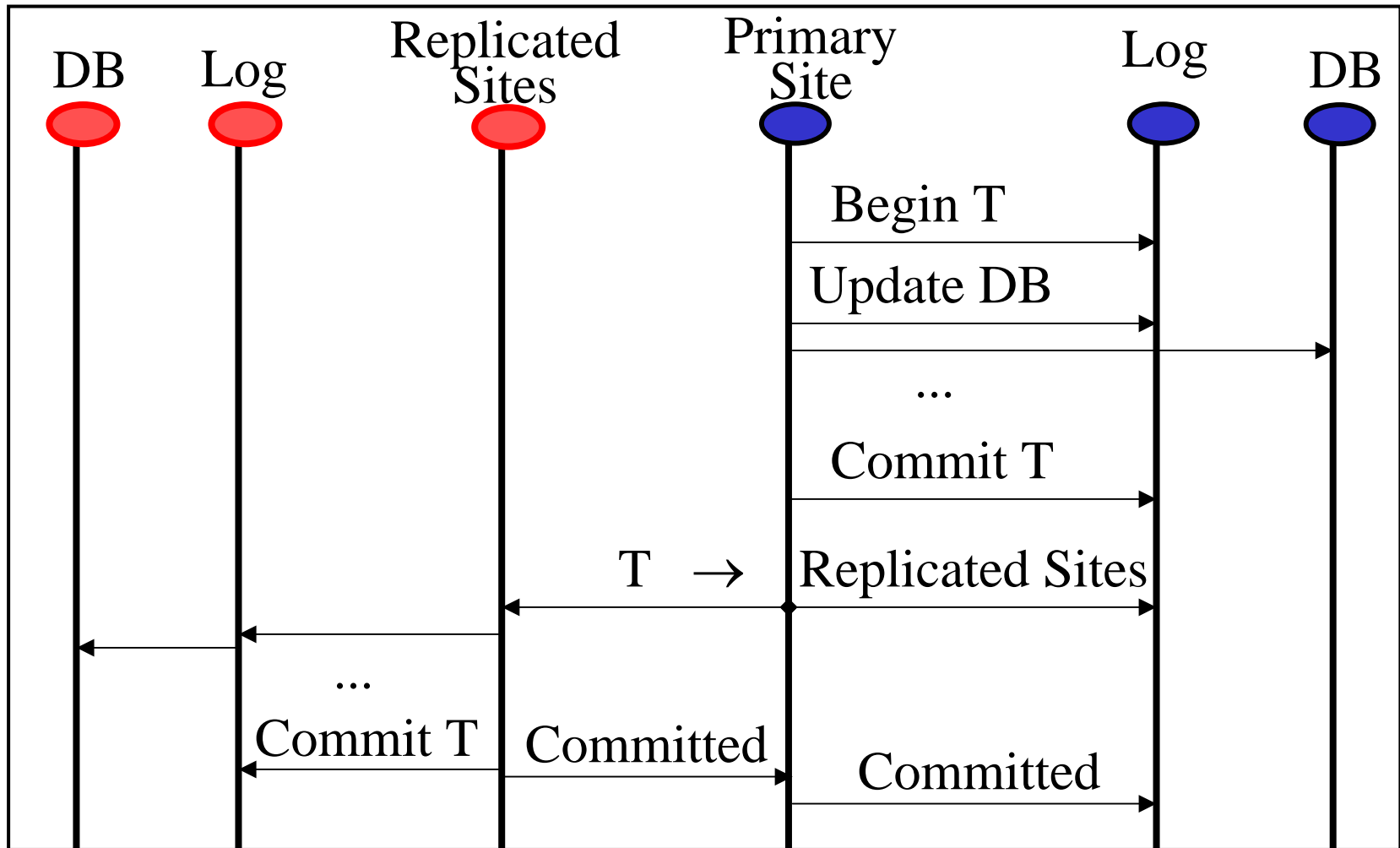
Replication and Distribution

- Case of Replication Servers with primary copy



Replication and Distribution

○ Case of Replication Servers with primary copy



Replication and Distribution

- Case of Replication Servers with primary copy
 - × Case of primary site failure
 - 7 Elect an other one

 - × Case of secondary site failure
 - 7 Exclude it during the failure
 - 7 Re-synchronization (Logs)

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Auditing Services

- Record events (*audit options*) about
 - x Database(s)
 - x User(s)
 - x Data Server(s)

- Process:
 1. Install or Configure the Audit Mechanism
 2. Define the audit options
 3. Manage the audit trail

Auditing Services in Oracle

1. Install or Configure the Audit Mechanism

- `audit_trail = true` (init file)
- `sys.aud$` exists

2. Define the audit options

1. Per session/per access
2. Whenever successful/not successful
3. Per type of SQL statements
4. Per type of system privilege commands
5. Per type of action on database objects

Auditing Services in Oracle: Examples

- *audit* select on U1.MaTable whenever not successful
- *audit* select table, update table by U1, U2
- *audit* role whenever successful
- *audit* all privileges by U1, U2
- *audit* insert, delete, update on sys.aud\$ per access
- *noaudit* select table by U1, U2

Auditing Services in Oracle

3. Manage the audit trail
 - a) Fix the right size (create/alter table)
 - b) Explore
 - Dictionary tables and views
 - *dba_stmt_opts*, *all_def_audit_opts*, etc.
 - c) Save or purge

Concluding Remarks

- DB security:
 - × Tools and mechanisms exist
 - × But, it's not only a matter of DB technology
 - 7 Platform
 - 7 Network
 - 7 Global enterprise security: policy and rules
- New challenges:
 - × Ubiquitous/mobile computing
 - × Replicated Mobile Data
 - × Web-based computing

The (Very) Basics of Database Security

That's all Folks !

Thank you for paying attention.

Any questions ?