1 - 2



- Pencils of circles
- Definition, empty circle property
- Motivations: reconstruction, meshing
- Properties: EMST, max-min angle
- Lower bound
- In-sphere predicate
- Diagonal flipping
- Incremental algorithm
- Sweep line algorithm
- Miscellaneous

Imagine moving circles

fixed center

Imagine moving circles

fixed center



Imagine moving circles



fixed center

Imagine moving circles



fixed center

Imagine moving circles



fixed center

increasing radius

Cocentric pencil

Imagine moving circles







Imagine moving circles



Imagine moving circles



Imagine moving circles



Imagine moving circles





Imagine moving circles



Imagine moving circles



Imagine moving circles







Imagine moving circles

Imagine moving circles



# Delaunay Triangulation: pencils of circles Imagine moving circles two fixed points Pencil with limit points radical axis 4 - 9



Imagine moving circles

Imagine moving circles

Imagine moving circles



Imagine moving circles
Imagine moving circles

a point on a line

Pencil with tangent point

radical axis/



Circle equation

$$x^2 + y^2 - 2ax - 2by + c = 0$$

Circle equation

$$x^2 + y^2 - 2ax - 2by + c = 0$$

Another circle equation

$$x^2 + y^2 - 2a'x - 2b'y + c' = 0$$

Circle equation

$$x^2 + y^2 - 2ax - 2by + c = 0$$

Another circle equation

$$x^2 + y^2 - 2a'x - 2b'y + c' = 0$$

Pencil of circles

$$\lambda \cdot (x^2 + y^2 - 2ax - 2by + c) + (1 - \lambda) \cdot (x^2 + y^2 - 2a'x - 2b'y + c') = 0$$

Circle equation

$$x^2 + y^2 - 2ax - 2by + c = 0$$

Another circle equation

$$x^2 + y^2 - 2a'x - 2b'y + c' = 0$$

Pencil of circles

$$\lambda \cdot (x^2 + y^2 - 2ax - 2by + c) + (1 - \lambda) \cdot (x^2 + y^2 - 2a'x - 2b'y + c') = 0$$

A special "circle: the radical axis

$$6-4$$
  $\lambda = \infty$ 

Power of a point w.r.t a circle

$$x^2 + y^2 - 2ax - 2by + c$$

Power of a point w.r.t a circle

$$x^2 + y^2 - 2ax - 2by + c$$

- = 0 on the circle
- < 0 inside the circle
- > 0 outside the circle

# **Delaunay Triangulation:** pencils of circles Power of a point w.r.t a circle

 $\lambda (x^2 + y^2 - 2a'x - 2b'y + c')$  $+(1-\lambda)(x^2+y^2-2ax-2by+c) = 0$ blue yields smaller power equal power black yields smaller power 7 - 3





























Output: the curve (the points in order along the curve)



Output: the curve (the points in order along the curve)





Output: the curve (the points in order along the curve)





Output: the curve (the points in order along the curve)



If good sampling, ouput  $\in$  Delaunay





Output: the surface (a triangulation of the points approximating the surface)



Output: the surface (a triangulation of the points approximating the surface)





Output: the surface (a triangulation of the points approximating the surface)



If good sampling, ouput  $\in$  Delaunay














11 - 6



#### 11 - 7

#### Delaunay Triangulation.

Teaser meshing lecture

#### sharp features



12 - 1

A spanning tree



Another spanning tree



The Euclidean Minimum-length Spanning Tree



The Euclidean Minimum-length Spanning Tree is included in Delaunay



The Euclidean Minimum-length Spanning Tree is included in Delaunay



Choose an edge of EMST

The Euclidean Minimum-length Spanning Tree





Is diametral circle empty ?



Is diametral circle empty ? assume ∃ blue point inside







Is diametral circle empty ?



Algorithm



Algorithm



Algorithm



Algorithm



Algorithm



Algorithm



Algorithm



Algorithm



Algorithm



Algorithm



Algorithm



Algorithm



Algorithm



Algorithm

choose shorter purple edge



13 - 14

Algorithm

choose shorter purple edge



 $O(n \log n)$  after Delaunay

13 - 15


# Delaunay Triangulation: size



14

# Delaunay Triangulation: size

Vertices Edges Faces

n - e + f = 2









n - e + t + 1 = 2





3t + k = 2e 2n - 3t - k + 2t = 2



# Delaunay Triangulation: size

$$\sum_{p \in S} d^{\circ}(p) = 2e = 6n - 2k - 6$$
$$\mathbb{E}(d^{\circ}(p)) = \frac{1}{n} \sum_{p \in S} d^{\circ}(p) < 6$$
average on the choice of point  $p$  in set of points  $S$ 

$$n - e + t + 1 = 2$$

$$\begin{array}{ll} 3t+k=2e & 2n-3t-k+2t=2 \\ & t=2n-k-2<2n \\ & e=3n-k-3<3n \end{array}$$





#### Triangulation

Delaunay





#### Triangulation

Delaunay

#### smallest angle



#### Triangulation

Delaunay

#### smallest angle





#### Triangulation

Delaunay

#### smallest angle





#### Triangulation

Delaunay

#### smallest angle

second smallest angle

Proof

Definition

Delaunay edge





Definition

locally Delaunay edge w.r.t. a triangulation



Definition

locally Delaunay edge w.r.t. a triangulation



neighbor = visible from the edge 16 - 5

Lemma ( $\forall$  edge: locally Delaunay)  $\iff$  Delaunay

### **Delaunay Triangulation:** max-min angleLemma ( $\forall$ edge: locally Delaunay) $\iff$ Delaunay

Proof:

choose an edge



Delaunay Triangulation:max-min angleLemma $(\forall edge: locally Delaunay) \iff Delaunay$ 

Proof:

- choose an edge
- edges of the quadrilateral are locally Delaunay



Delaunay Triangulation:max-min angleLemma $(\forall edge: locally Delaunay) \iff Delaunay$ 

Proof:

- choose an edge
- edges of the quadrilateral are locally Delaunay



Delaunay Triangulation:<br/>max-min angleLemma(∀ edge: locally Delaunay) ⇔ Delaunay

Proof:

- choose an edge
- edges of the quadrilateral are locally Delaunay

Vertices visible through one edge are outside circle

Delaunay Triangulation:max-min angleLemma $(\forall edge: locally Delaunay) \iff Delaunay$ 

Proof:

- choose an edge
- edges of the quadrilateral are locally Delaunay

Vertices visible through one edge are outside circle

Induction  $\rightarrow$  all vertices outside circle

Lemma For four points in convex position

Delaunay  $\iff$  maximize the smallest angle

Two possible triangulation





Lemma For four points in convex position

Delaunay  $\iff$  maximize the smallest angle



Lemma For four points in convex position

Delaunay  $\iff$  maximize the smallest angle



 $\exists$  a smaller angle  $\in$  other triangulation

Lemma For four points in convex position

 $Delaunay \iff maximize \ the \ smallest \ angle$ 



Lemma For four points in convex position

 $Delaunay \iff maximize \ the \ smallest \ angle$ 





 $\exists$  a smaller angle  $\in$  other triangulation

Map: Triangulations  $\longrightarrow \mathbb{R}^{6n-3k-4}$  smallest angle  $\alpha_1$ 



Map: Triangulations  $\longrightarrow \mathbb{R}^{6n-3k-4}$ 

smallest angle  $\alpha_1$ second smallest angle  $\alpha_2$ 



Map: Triangulations  $\longrightarrow \mathbb{R}^{6n-3k-4}$ 

smallest angle  $\alpha_1$ second smallest angle  $\alpha_2$ third smallest angle  $\alpha_3$ 



Map: Triangulations  $\longrightarrow \mathbb{R}^{6n-3k-4}$ 

 $(\alpha_1, \alpha_2, \alpha_3, \ldots, \alpha_{6n-3k-4})$ 

smallest angle  $\alpha_1$ second smallest angle  $\alpha_2$ third smallest angle  $\alpha_3$ 




Theorem:

Delaunay maximizes minimum angles (in lexicographic order)

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Let T be the triangulation maximizing angles

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Let T be the triangulation maximizing angles

 $\implies \forall \text{ convex quadrilateral (from 2 triangles } \in T)$ 

the diagonal maximizes smallest angle (in quad)

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Delaunay maximizes minimum angles (in lexicographic order) Proof:

Let T be the triangulation maximizing angles

 $\implies \forall \text{ convex quadrilateral (from 2 triangles} \in T)$ the diagonal maximizes smallest angle (in quad)

 $\Longrightarrow \forall$  edge, it is locally Delaunay

Theorem:

Delaunay maximizes minimum angles (in lexicographic order) Proof:

Let T be the triangulation maximizing angles

 $\implies \forall \text{ convex quadrilateral (from 2 triangles} \in T)$ the diagonal maximizes smallest angle (in quad)

 $\Longrightarrow \forall$  edge, it is locally Delaunay

 $\implies T = \mathsf{Delaunay}$ 

#### Delaunay Triangulation: lower bound



#### Delaunay Triangulation: lower bound





22

#### 22 - 2

#### query s

#### pqr ccw triangle



#### Delaunay Triangulation: incircle predicate



pqr ccw triangle

query s inside circumcircle

#### 22 - 4

query s cocircular

pqr ccw triangle



#### Delaunay Triangulation: incircle predicate



pqr ccw triangle

query s outside circumcircle





Space of circles



 $p = (x, y) \rightsquigarrow p^{\star} = (x, y, x^2 + y^2)$ 

Space of circles



 $p = (x, y) \rightsquigarrow p^{\star} = (x, y, x^2 + y^2)$ 

$$C: x^2 + y^2 - 2ax - 2by + a^2 + b^2 - r^2 = 0$$
$$\rightsquigarrow C^* = (a, b, a^2 + b^2 - r^2)$$

Space of circles



 $p = (x, y) \rightsquigarrow p^{\star} = (x, y, x^2 + y^2)$ 

$$C: x^{2} + y^{2} - 2ax - 2by + a^{2} + b^{2} - r^{2} = 0$$
  

$$\rightsquigarrow C^{\star} = (a, b, a^{2} + b^{2} - r^{2})$$
  

$$\rightsquigarrow C^{\dagger}: z - 2ax - 2by + a^{2} + b^{2} - r^{2} = 0$$

Space of circles

$$p \in C \Longleftrightarrow p^{\star} \in C^{\dagger}$$



 $p = (x, y) \rightsquigarrow p^{\star} = (x, y, x^2 + y^2)$ 

$$C: x^{2} + y^{2} - 2ax - 2by + a^{2} + b^{2} - r^{2} = 0$$
  

$$\rightsquigarrow C^{\star} = (a, b, a^{2} + b^{2} - r^{2})$$
  

$$\rightsquigarrow C^{\dagger}: z - 2ax - 2by + a^{2} + b^{2} - r^{2} = 0$$

Delaunay Triangulation: incircle predicate



22 - 11  $\rightsquigarrow C^{\dagger}: z - 2ax - 2by + a^2 + b^2 - r^2 = 0$ 





incircle predicate

 $\rightsquigarrow$  3D orientation predicate



incircle predicate

 $\rightsquigarrow$  3D orientation predicate

sign 
$$\begin{vmatrix} 1 & 1 & 1 & 1 \\ x_p & x_q & x_r & x_s \\ y_p & y_q & y_r & y_s \\ x_p^2 + y_p^2 & x_q^2 + y_q^2 & x_r^2 + y_r^2 & x_s^2 + y_s^2 \end{vmatrix}$$





Degree 4 vertex in Voronoi diagram



Degree vertex in Voronoi diagram d



#### Degree 4 vertex in Voronoi diagram Delaunay quad ? random diagonal ? 23 - 4



Degree 4 vertex in Voronoi diagram

Delaunay quad ? random diagonal ? 23 - 5

# Delaunay Triangulation: Teaser CAL lecture

#### Data structure for (Delaunay) triangulation

Representing incidences

Representing hull boundary

Representing user's data

put colors in triangles

• • •















check edges of quadrilateral












































Delaunay is obtained

Complexity ?

Complexity ?





Locally Delaunay

Complexity ?

Convex



Delaunay

Complexity ?

Non convex



Delaunay Triangulation: Diagonal flipping Complexity ? Non convex edge















Complexity ?





Complexity ?




27 - 3











Complexity ?

Delaunay











New point



New point

Locate



New point

Locate



e.g.: straight walk

28 - 4

New point

Locate



New point

Locate



New point

Locate



New point

Locate



New point

Locate



New point

Locate









New point

Locate



e.g.: visibility walk

28 - 14

New point

Locate



New point

Locate



New point

Locate







#### Delaunay Triangulation: incremental algorithm Visibility walk terminates ?

#### Delaunay Triangulation: incremental algorithm Visibility walk terminates



#### Delaunay Triangulation: incremental algorithm Visibility walk terminates ?



# Delaunay Triangulation: incremental algorithm Visibility walk terminates



May loop

### Delaunay Triangulation: incremental algorithm Visibility walk terminates ?



28 - 24

Visibility walk terminates



Visibility walk terminates

Green power < Red power

28 - 26
Visibility walk terminates

Green power < Red power

Power decreases

28 - 27

Visibility walk terminates

Green power < Red power

Power decreases

Visibility walk terminates

28 - 28



Locate







Locate





Locate





Locate





Locate





Locate





Locate





Locate





Locate





New point

Locate



New point

Locate



New point



New point



Complexity

Locate

Complexity

Locate

Search conflicts

# triangles in conflict

# triangles neighboring triangles in conflict

Complexity

Locate

Search conflicts

‡ triangles in conflict

# triangles neighboring triangles in conflict

degree of new point in new triangulation

< n

Complexity

Locate Walk may visit all triangles < 2n

Search conflicts

#### degree of new point in new triangulation

< n

Complexity

Locate O(n) per insertion

Complexity

Locate O(n) per insertion

Search conflicts

 ${\cal O}(n^2)$  for the whole construction

Complexity

Locate

Search conflicts

half-parabola and circle

Complexity

Locate

Search conflicts

half-parabola and circle

Delaunay triangle

Complexity

Locate

Complexity

Locate

Search conflicts

Insertion:  $\Omega(n)$ 

Whole construction:  $\Omega(n^2)$ 

Complexity

In practice

Locate Many possibilities (walk, Delaunay hierarchy)

Search conflicts Randomized



Discover the points from left to right



Discover the points from left to right



Discover the points from left to right



Certified Delaunay triangles

Discover the points from left to right



Certified Delaunay triangles

Discover the points from left to right



Discover the points from left to right



Certified Delaunay triangles Certified Delaunay edges

Discover the points from left to right



Boundary edges

Discover the points from left to right



Boundary edges
Discover the points from left to right

Boundary edges

Empty circles tangent to sweep line













Discover the points from left to right



Closing a triangle ?







Discover the points from left to right



Closing a triangle ?

#### Next circle event

Discover the points from left to right



Next circle event Close triangle

31 - 21



Discover the points from left to right



31 - 23

Discover the points from left to right



31 - 24

Delaunay Triangulation: sweep-line algorithm		
Complexity	Circle events	Point events
Number		
Triangulation		
List of events ( $x$ sorted)		
List of boundary edges (ccw sorted)		
32 - 1		

Delaunay Triangulation: sweep-line algorithm		
Complexity	Circle events processed	Point events
Number		
Triangulation		
List of events $(x \text{ sorted})$		
List of boundary edges (ccw sorted)		
32 - 2		

Delaunay Triangulation: sweep-line algorithm		
Complexity	Circle events processed	Point events
Number	2n	n
Triangulation		
List of events $(x \text{ sorted})$		
List of boundary edges (ccw sorted)		
32 - 3		

Delaunay Triangulation: sweep-line algorithm		
Complexity	Circle events processed	Point events
Number	2n	n
Triangulation	create 2 triangles per event	create one edge per event
List of events $(x \text{ sorted})$		
List of boundary edges (ccw sorted)		
32 - 4		

Delaunay Triangulation: sweep-line algorithm		
Complexity	Circle events processed	Point events
Number	2n	n
Triangulation	create 2 triangles per event	create one edge per event
List of events $(x \text{ sorted})$	$\leq 3$ deletions $\leq 2$ insertions per event	$\leq 2$ deletions $\leq 2$ insertions per event
List of boundary edges (ccw sorted)		
32 - 5		

Delaunay Triangulation: sweep-line algorithm		
Complexity	Circle events processed	Point events
Number	2n	n
Triangulation	create 2 triangles per event	create one edge per event
List of events $(x \text{ sorted})$	$\leq 3$ deletions $\leq 2$ insertions per event	$\leq 2$ deletions $\leq 2$ insertions per event
List of boundary edges (ccw sorted) 32 - 6	replace 2 edges by 1 per event	locate, then insert 2 edges per event

Delaunay Triangulation: sweep-line algorithm		
Complexity	Circle events processed	Point events
Number	2n	n
TriQ(1) per operation	create 2 triangles per event	create one edge per event
$O(\log n)$ per operation sorted	$\leq 3$ deletions $\leq 2$ insertions per event	$\leq 2$ deletions $\leq 2$ insertions per event
List of boundary operation $O(\log n)$ per operation (ccw sorted) 32 - 7	replace 2 edges by 1 per event	locate, then insert 2 edges per event

Delaunay Triangulation: sweep-line algorithm		
Complexity	Circle events processed	Point events
Number	2n	n
f(0(1)) per operation	create 2 triangles	create one edge
nangulation	O(n)	$\log n$
$o(\log n)$ per operation		
10(105-)	per event	per event
List of boundary operation $O(\log n)$ per operation (CCW sorted) 32 - 8	replace 2 edges by 1 per event	locate, then insert 2 edges per event













# Delaunay Triangulation: incremental algorithm

New point









#### Delaunay Triangulation: deletion algorithm (sketch) Extract hole


Extract hole

Triangulate







# Delaunay Triangulation: deletion algorithm (sketch) Extract hole Triangulate and sew 34 - 10







Ear queue

Ear with largest power is added



Ear queue

Ear with largest power is added



Ear queue

Ear with largest power is added

Iterate



Ear queue

Ear with largest power is added

Iterate



#### Delaunay Triangulation: deletion algorithm (sketch) Triangulate and flip



#### Delaunay Triangulation: deletion algorithm (sketch) Triangulate and flip



#### Delaunay Triangulation: deletion algorithm (sketch) Triangulate and flip



for degree  $\geq 8$ 



#### Delaunay Triangulation: deletion algorithm (sketch) Decision tree for small holes



for degree  $\leq 7$ 









#### Delaunay Triangulation: deletion algorithm (sketch) Decision tree for small holes





for degree  $\leq 7$ 

#### Delaunay Triangulation: deletion algorithm (sketch) Decision tree for small holes





ensions
m <i>i</i> )
-1 + 1
$_{d} = 1$

Duality with 4D convex hull

Incremental algorithm (find the hole and star)

Delaunay	Convex hull	Higher dimensions	
	Dehn Sommerville relations	$f_i = \sharp(faces \ of \ dim \ i)$	
Same as 2[	Euler: $f_0 - f_1 + f_2$ -	$-\dots f_{d-1} = (-1)^{d-1} + 1$	
Dual	$\sum_{j} = k^{d-1} - 1^{j} \begin{pmatrix} j+1\\ k+1 \end{pmatrix}$	$\int f_j = (-1)^{d-1} f_k$	
Empt	$-1 \le k \le d-2$	$f_{-1} = f_d = 1$	
Tri	$\left\lfloor \frac{d+1}{2} \right\rfloor$ independent	equations	
Duality with 4D convex hull quadratic .			
Incremental algorithm (find the hole and star)			













Algorithms

4D convex hull duality



Incremental

Algorithms

4D convex hull duality

Flip

Incremental

 $O(f\log n + n^{rac{4}{3}})$  or  $\Theta(n^2)$ 



# Delaunay Triangulation: higher dimensions

d+1 convex hull duality

 $O\left(n^{\lfloor \frac{d+1}{2} \rfloor}\right)$ 

Incremental practical O(n) for random points

coeff exponential in  $\boldsymbol{d}$ 

