Variable-width contouring for additive manufacturing

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Context: 3D printing





Context: Fabricating one layer



In turn, each layer is fabricated by solidifying a **bead** of some material, along a **print path**.

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Overfill



Overfill

Overfill = forbidden...

...except for **closed beads**, a well controlled special case:

(we love closed beads!)

Underfill

Example: two classic ways to fill a square with a **constant-width** bead.





Underfill

Underfill is the existence of areas of the slice **not** covered by a solid bead.



Underfill

Underfill is bad. We want to minimize it.

Our contribution is a new technique for designing print paths that produces

- no overfill (this is somewhat easy)
- a small amount of underfill (almost 10x less than the state of the art)





What to do?

and curvature.

Earlier works suggesting to use variable-width beads:

• Jin, Du, and He. Journal of Manufacturing Systems 44 (2017).

=overfill

=underfill

• Kuipers, Doubrovski, Wu, and Wang. Computer-Aided Design 128 (2020).

We follow suit, use closed, variable-width beads and try to minimize their number

Uniform-width parallel contouring wo. (left) and w. (right) regularization.

Our technique.

Inputs & data structure

A range [2γ, 2Γ] of feasible bead widths — → •
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- A 2γ-fat planar shape S: all the maximal disks inside S have radius ≥ 2γ.¹

 1 In practice, slices are polygons. We process them into $2\gamma\text{-fat}$ shapes.

Inputs & data structure

- A range $[2\gamma, 2\Gamma]$ of feasible bead widths (specific to target 3D printer).
- A 2γ -fat planar shape S: all the maximal disks inside S have radius $\geq 2\gamma$.¹
- An explicit representation of the medial axis of S:
 MA(S) is the closure of the set of centers of maximal disks in S.



 1 In practice, slices are polygons. We process them into 2γ -fat shapes.

Given a shape S, we model a bead that stays in contact with the boundary of S and make the remaining inner shape "rounder."



To do so, we replace parts of the boundary ∂S by inner tangent circular arcs (yellow)...



Then we do a parallel offset of 2γ and obtain a bead of width within the allowed range.



Now we repeat the process



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The circles supporting the tangent circular arcs are chosen as the boundary of maximal disks in S. Hence, their center lies on the medial axis MA(S) of S.



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Replacing by circular arc = **trimming** the medial axis!



Variable-width contouring: basics

1. Trimming the medial axis: removes crescents of width $\leq 2\Gamma - 2\gamma$ from the shape.

2. Parallel offset : removes a band of width exactly 2γ , which together with the crescents, form a bead of width varying within $[2\gamma, 2\Gamma]$.

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If the input is a polygon, then:

• the medial axis is computable (CGAL, BOOST) and

• the two operations above produce shapes with **linear** or **circular** boundary arcs only.

Corollary: in that case, each bead is bounded by linear or circular arcs only.



Trimming



Trimming



The algorithm grows a tree from each leaf (degree-1 vertex) and finds all maximal trimmable trees.

Complete picture with Collapsing





 \mathcal{S}^i (a) \downarrow trimming $\mathcal{S}^i_{\mathbf{tr}}$

11.1

Complete picture with Collapsing





(a) \mathcal{S}^{i} trimming Δ $\mathcal{S}^i_{\mathbf{tr}}$ i+1







Complete picture with Collapsing





 T^i (e) $\operatorname{traj}(T^i) \subset \mathbf{M} \dot{\mathbf{A}} \left(\mathcal{S}^i \setminus (\mathcal{S}^{i+1} \cup K^i) \right)$

(a) \mathcal{S}^{i} trimming S^{i+1} $\mathcal{S}^i_{\mathbf{tr}}$ collapsing Δ offset $(\mathsf{d}), \mathcal{S}^{i+1}$ (C 11.6



Pictures











End of this presentation

See the paper for more, including:

- Less underfill with shaving
- An algorithm for sampling the print path (the center curve of each bead)
- A comparison with the state of the art (almost 10x less underfill)
- A proof of the absence of overfill
- More pictures of fabricated layers

Code: https://github.com/mfx-inria/Variable-width-contouring