

# Cast as Intended in voting protocols

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DE LORRAINE

informatiques mathématiques  
*loria*



What is a good voting system?

# Confidentiality of the votes

## Vote privacy

*"No one should know how I voted"*



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Better: Receipt-free / Coercion-resistant

*"No one should know how I voted,  
even if I am willing to tell my vote!"*



- ▶ vote buying
- ▶ coercion



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**Everlasting privacy:** no one should know my vote, even when the cryptographic keys will be eventually broken.

# Verifiability

**Individual Verifiability:** a voter can check that

- ▶ cast as intended: their ballot contains their intended vote
- ▶ recorded as cast: their ballot is in the ballot box.

**Universal Verifiability:** everyone can check that

- ▶ tallied as recorded: the result corresponds to the ballot box.
- ▶ eligibility: ballots have been casted by legitimate voters.



You should verify the election,  
not the system.

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**Even better: accountability**

- ▶ the system tells whom to blame
- ▶ eases dispute resolution

## And many more properties

- ▶ **Availability**: servers available at any time
- ▶ **Accessibility**: easy to use, adapted to people with various issues
- ▶ ...



## And many more properties

- ▶ **Availability**: servers available at any time
- ▶ **Accessibility**: easy to use, adapted to people with various issues
- ▶ ...

In this talk, focus on verifiability.

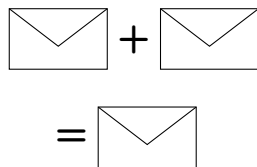
- ▶ cast as intended
- ▶ recorded as cast
- ▶ tallied as recorded
- ▶ eligibility verifiability

# Tallied as recorded

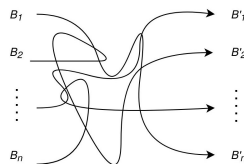
*The result corresponds to the ballot box.*

- ✓ Well studied academically, with two main techniques:

Homomorphic tally



Mixnet

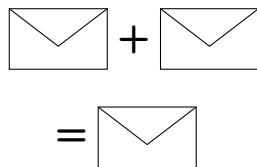


# Tallied as recorded

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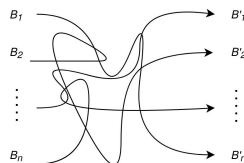
Homomorphic tally



In practice:

- ✓ Many deployed solutions use such techniques: Estonia, France, Switzerland, ...
- Many national voting companies are still behind

Mixnet



## Recorded as cast

- ✓ easy in theory: the voter simply checks that their ballot appear on the bulletin board
- Not so easy in practice
  - ▶ require a **public** bulletin board
  - ▶ voters do not check

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### Alternative approaches

- ▶ delegation: voters send their ballot to a third party (eg French Legislative system, Polyas, ...)
- ▶ Swiss Post: the verification is embedded in the cast-as-intended mechanism, requires **distributed servers**
- ▶ Estonia: the ballot is sent (by the server) to another system component

# Eligibility verifiability

✓academically: just sign but...

- ▶ require a PKI
- ▶ public voter list? everlasting privacy?

# Eligibility verifiability

✓ academically: just sign but...

- ▶ require a PKI
- ▶ public voter list? everlasting privacy?

## In practice

✓ Estonia: voters sign with their id cards

✓ strong and verifiable eligibility

■ no public board

✗ login/password sent by mail, SMS → no eligibility verifiability

■ distributed trust between authorities, eg Belenios (OTP + asymmetric key credential)

# Cast as intended (Cal)

- Few academic protocols

??? but yet a lot of systems in practice!



# Cal in Australia

iVote system in the 2015 state election in New South Wales



What is my vote?



v (in clear!)



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iVote system in the 2015 state election in New South Wales




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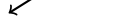
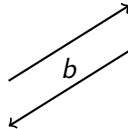
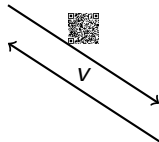
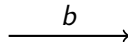
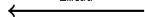
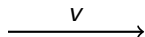


v (in clear!)



- ✓ simple
- ✓ cast-as-intended
- ✗ no vote privacy 
- ✗ no cast-as-recorded

# Cal in Estonia



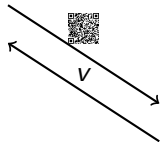
# Cal in Estonia



$v$



$b$



$b$

- ✓ cast-as-intended
- some vote-buying threats (mitigated)
- proxy cast-as-recorded
- heavy infrastructure (two independent servers)

# Cal in Switzerland

Choice Return Code:

Question 1:

YES: 1225

NO: 7092

EMPTY: 2812

Question 2:

YES: 9817

NO: 2111

EMPTY: 6745



Please check that your device displays the correct **choice return codes**.  
If you cannot see the correct codes or in case of doubt, please contact the election authorities (0XX / XXX XX XX).

- ✓ cast-as-intended
- proxy cast-as-recorded
- heavy infrastructure (four independent servers)

# Benaloh's challenge: cast or spoil

$v$



$\xrightarrow{v}$

$h = \text{hash}(b)$

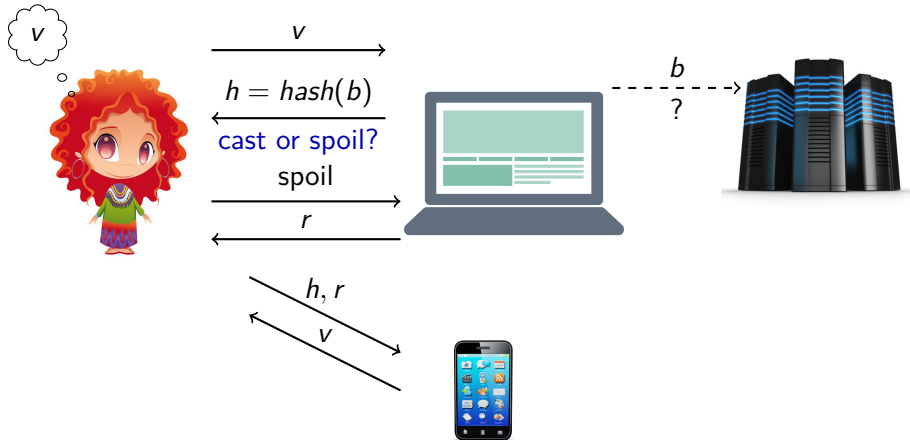
$\xleftarrow{\text{cast or spoil?}}$



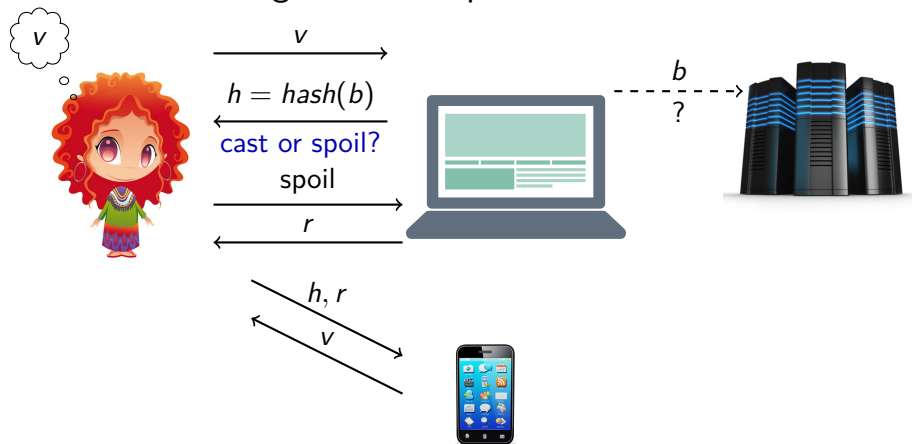
$\xrightarrow{\frac{b}{?}}$



# Benaloh's challenge: cast or spoil



# Benaloh's challenge: cast or spoil



- ✓ simple principle
- ✓ can be adapted to many systems
- requires a second device



## Choice of the EVoteID nicest location

To vote, follow these steps:

1. **Select** your preferred options.
2. **Review** your choices, which are then encrypted.
3. **Submit** your encrypted ballot and authenticate to verify your eligibility.

Start

You can [email for help](#).

Election Fingerprint: **7hBzRD3Am/07fq171xJwYTJ4j0ENV77dDw9NY370UWA**

## Choice of the EVoteID nicest location

(1) Select

(2) Review

(3) Submit

### What is the best location for EVoteID

#1 of 1 — vote for 1

- Bregenz
- Luxembourg

Proceed

Election Fingerprint: 7hBzRD3Am/07fqL71xJwYTJ4j0ENV77dDw9NY370UWA

## Choice of the EVotID nicest location

[\(1\) Select](#)**[\(2\) Review](#)**[\(3\) Submit](#)

### Review your Ballot

Question #1: What is the best location for EVotID

✓ **Bregenz**

[\[change\]](#)

Your ballot tracker is **+JLW+ti+ERZL0jPQNeRIAFi7RD6ZHIakX9a6n6XFwno.**

[Submit this Vote](#)

[Spoil & Audit](#) [optional]

Election Fingerprint: **7hBzRD3Am/07fqL71xJwYTJ4j0ENV77dDw9NY370UWA**

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If you choose, you can spoil this ballot and reveal how your choices were encrypted. This is an optional auditing process.

You will then be guided to re-encrypt your choices for final casting.

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[Spoil & Audit](#)

**Vote privacy issue!**

The voter is likely to use their **true** vote.

# Benaloh: voter strategy

A voter should

1. decide **at random** if they will truly vote or audit
2. → if vote, then vote  
→ if audit, decide **at random** then audit and go to step 1

✗ usability

■ which probabilities to use?

▶ is it truly cast-as-intended? (see e.g. [Jamroga's talk](#))

# Other Cal solutions

Select, Selene, Hyperion: votes appear in clear on the ballot box

- ✓ simple for the voters
- ▶ specific systems
- adversary caught to late  
→ strong accountability needed

## Other Cal solutions

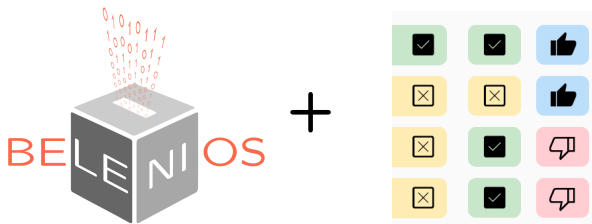
Select, Selene, Hyperion: votes appear in clear on the ballot box

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Two device solutions: Du-Vote, CAISED (this Friday!)



# Our proposal: BeleniosCal



- ▶ based on Belenios
- ▶ could be adapted to other protocols
- ▶ **no second device** (except to read BB), no paper material
- ▶ **on the fly detection**
- ▶ **one server**

# Voting protocol Belenios



- ▶ variant of Helios, designed by Ben Adida
- ▶ developed at Loria, teams Pesto and Caramba (P. Gaudry)  
Developer: Stéphane Glondu
- ▶ used in 2000+ elections, with a total of 100 000+ voters

<http://www.belenios.org/>

- ▶ confidentiality of the votes
- ▶ verifiability of the voting process
  - The ballot box is public at any time.
  - All the operations (tally, ...) can be checked by anyone.

# How Belenios works (simplified)

## Phase 1: vote



$pkE$

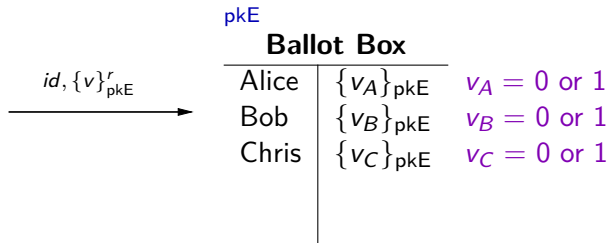
### Ballot Box

Alice	$\{v_A\}_{pkE}$	$v_A = 0 \text{ or } 1$
Bob	$\{v_B\}_{pkE}$	$v_B = 0 \text{ or } 1$
Chris	$\{v_C\}_{pkE}$	$v_C = 0 \text{ or } 1$

$pkE$ : public key, the private keys are shared among the authorities.

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Chris	$\{v_C\}_{pkE}$	$v_C = 0 \text{ or } 1$
David	$\{v_D\}_{pkE}$	$v_D = 0 \text{ or } 1$

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pkE

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Chris	$\{v_C\}_{pkE}$	$v_C = 0 \text{ or } 1$
David	$\{v_D\}_{pkE}$	$v_D = 0 \text{ or } 1$
...	...	

## Phase 2: Tally - homomorphic encryption (El Gamal)

$$\{v_1\}_{pkE} \times \cdots \times \{v_n\}_{pkE} = \{v_1 + \cdots + v_n\}_{pkE} \quad \text{since } g^a \times g^b = g^{a+b}$$

→ Only the final result needs to be decrypted! **And proved.**

pkE: public key, the private keys are shared among the authorities.

# Eligibility



$id, \{v\}_{pkE}^r$

→

$pkE$

## Ballot box

Alice	$\{v_A\}_{pkE}$
Bob	$\{v_B\}_{pkE}$
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...	...
...	...

# Eligibility



$id, \{v\}_{pkE}^r$

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$pkE$

## Ballot box

Alice	$\{v_A\}_{pkE}$
Bob	$\{v_B\}_{pkE}$
Chris	$\{v_C\}_{pkE}$
...	$\{1\}_{pkE}$
...	$\{1\}_{pkE}$

The ballot box could add ballots!



# Eligibility



$id, \{v\}_{pkE}^r$

→

$pkE \quad vk(cred_3), vk(cred_1), vk(cred_2), \dots$

## Ballot box

Alice	$\{v_A\}_{pkE}$
Bob	$\{v_B\}_{pkE}$
Chris	$\{v_C\}_{pkE}$
...	
...	

~~The ballot box could add ballots!~~

1. During the setup phase, a Registrar generates private signing keys, one for each voter

# Eligibility



$id, \{v\}_{pkE}^r$

→

$pkE \quad vk(cred_3), vk(cred_1), vk(cred_2), \dots$

## Ballot box

Alice	$[\{v_A\}_{pkE}]_{sk(cred_1)}$
Bob	$[\{v_B\}_{pkE}]_{sk(cred_2)}$
Chris	$[\{v_C\}_{pkE}]_{sk(cred_3)}$
...	
...	

~~The ballot box could add ballots!~~

1. During the setup phase, a Registrar generates private signing keys, one for each voter
2. The voters sign their ballot with a “credential” they have received (a credential = a right to vote)

# BeleniosCal's principle

Alice' vote

$$bal = \{v\}_{pkE}^{r_v},$$

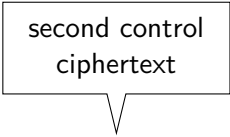
# BeleniosCal's principle

first control  
ciphertext

$$bal = \{v\}_{pkE}^{r_v}, \{a\}_{pkE}^{r_a},$$

# BeleniosCal's principle

second control  
ciphertext



$$bal = \{v\}_{pkE}^{r_v}, \{a\}_{pkE}^{r_a}, \{b\}_{pkE}^{r_b},$$

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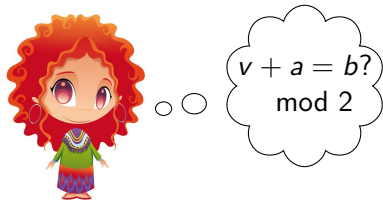
proof that binds  
the three ciphertexts

$$ZKP(v + a = b \pmod{2})$$

# BeleniosCal's principle

$$bal = \{v\}_{pkE}^{r_v}, \{a\}_{pkE}^{r_a}, \{b\}_{pkE}^{r_b},$$

$$ZKP(v + a = b \pmod 2)$$



# BeleniosCal's principle

$$bal = \{v\}_{pkE}^{r_v}, \boxed{\{a\}_{pkE}^{r_a}}, \{b\}_{pkE}^{r_b}, \quad ZKP(v + a = b \pmod 2)$$



Pick one of the two control ciphertexts  
at random



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$$bal = \{v\}_{pkE}^{r_v}, \{a\}_{pkE}^{r_a}, \boxed{\{b\}_{pkE}^{r_b}}, \quad ZKP(v + a = b \pmod{2})$$



Pick one of the two control ciphertexts  
at random

checks that  $bal, (\_, b)$  appears on the ballot box

# BeleniosCal's principle - continued

$v$



$\xrightarrow{v}$

$h = \text{hash}(bal)$

$v + a = b?$



# BeleniosCal's principle - continued

$v$



$v$

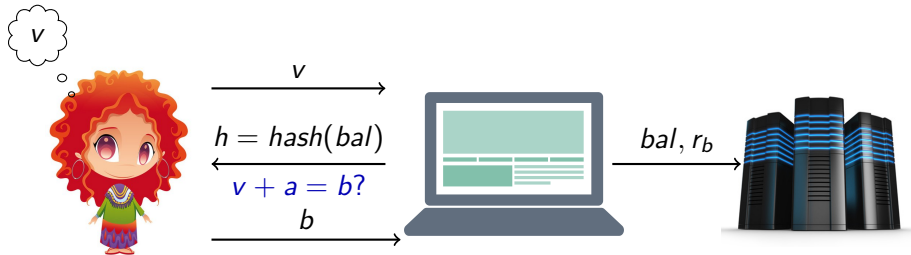
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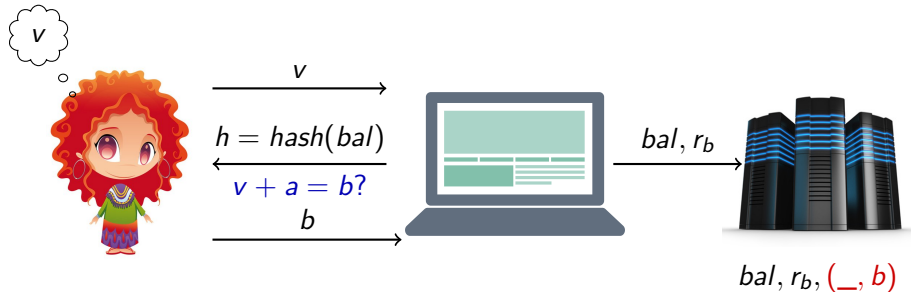
$b$



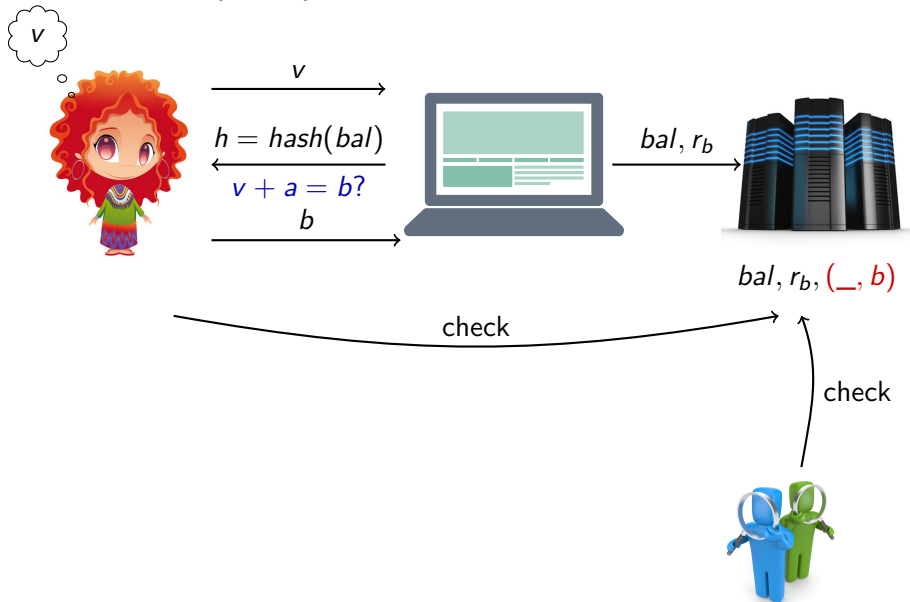
# BeleniosCal's principle - continued



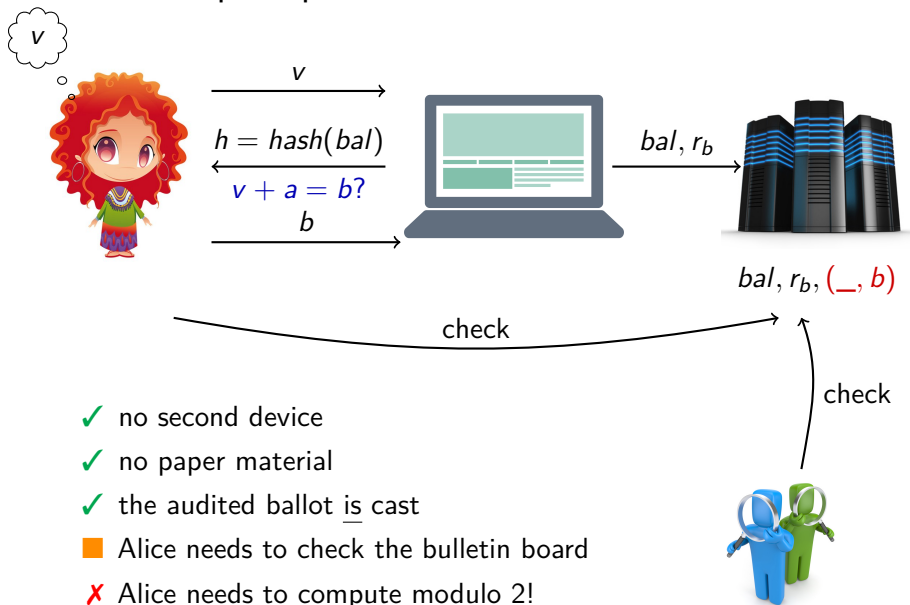
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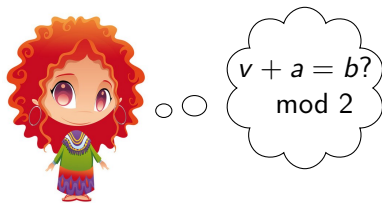
## BeleniosCal's principle - continued



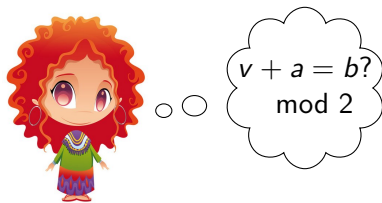
- ✓ no second device
- ✓ no paper material
- ✓ the audited ballot is cast
- Alice needs to check the bulletin board
- ✗ Alice needs to compute modulo 2!



Can voters compute modulo 2 ?!?



Can voters compute modulo 2 ?!?



Let see how we propose to implement it.



## Best dessert

● **Input credential**

○ Answer to questions

○ Review and encrypt

○ Authenticate

○ Security check

○ Confirm

Please enter your credential:

Next

Election UUID: JDwmiDBK8QQK6x

Election fingerprint: djB76kleknKDujNCVzZJ6o5dhStlmmYEIz2TCEIaFe4



# Best dessert



What is your favorite dessert?

Please select 1 answer.

Cheese cake

Tiramisu

Chocolate cake

Blank vote

Question 1 of 1

Next

Election UUID: JDwmiDBK8QQK6x

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# Best dessert



## Review your answers

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Blank vote



Previous

Next

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## Best dessert



### Save your tracking number

Your ballot has been encrypted, but not cast yet!

Your tracking number is:

Please save it to check later that your vote has been taken into account.

Previous

Next

Election UUID: JDwmiDBK8QQK6x

Election fingerprint: djB76kleknKduJNCVzZJ6c5dhSllmnYEIz2TCElaFe4



## Authenticate

A verification code has been sent to veronique.cortier@loria.fr.

Please enter the verification code received by e-mail:

Powered by [Belenios 2.2 \(2.1-288-gd00ef982\)](#). [Get the source code](#) [Privacy policy](#) [Administer elections](#)

Input credential

Answer to questions

Review and encrypt

Authenticate

**Security check**

Confirm

## Security check

- Determine for each line whether the control value is identical or not to your vote.
- Select a control pattern by picking one symbol per line.
- Save your control pattern to compare it later with the one displayed in the ballot box.

[More info](#)

### What is your favorite dessert?

Your vote

Control value

Cheese cake



Are the two symbols identical?

yes no 

Tiramisu



Chocolate cake



Blank vote

[Previous](#)[Next](#)

Election UUID: JDwmiDBK8QQK6x

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Are the two symbols identical?

yes

no

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Previous

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Blank vote

[Previous](#)[Next](#)

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**Security check**

Confirm

## Security check

- Determine for each line whether the control value is identical or not to your vote.
- **Select a control pattern by picking one symbol per line.**
- Save your control pattern to compare it later with the one displayed in the ballot box.

[More info](#)

### What is your favorite dessert?

Your vote

Control pattern

[Previous](#)[Next](#)

Election UUID: JDwmiDBK8QQK6x

Election fingerprint: djB76kleknKDwJNCVzZJ6o5dhStlmmYEIz2TCEIaFe4



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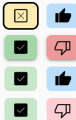
[More info](#)

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Your vote


Control pattern



Pick one of the two symbols

Previous

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Input credential

Answer to questions

Review and encrypt

Authenticate

**Security check**

Confirm

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### What is your favorite dessert?

Your vote

Control pattern

<input checked="" type="checkbox"/>	<input type="checkbox"/>
<input type="checkbox"/>	<input checked="" type="checkbox"/>
<input type="checkbox"/>	<input checked="" type="checkbox"/>
<input checked="" type="checkbox"/>	<input type="checkbox"/>



Take a picture, a screenshot or  
copy your control pattern

[Copy control pattern](#)[Previous](#)[Next](#)



# Best dessert

Input credential

Answer to questions

Review and encrypt

Authenticate

Security check

Confirm

## Thank you for voting!

### Next steps

- Follow the link in your confirmation email
- Verify your control pattern
- If your ballot is missing or the control pattern does not match, contact the administrator: cortier

### About your ballot

Voter	veronique.cortier@loria.fr
Tracking number	HoVF28p19LEUtK0EkoNaitT7RzhM6AM1BJQLnCLic8g
Status	accepted
Revote	yes
Email sent	yes

[Go back to election](#)



## Best dessert - Accepted ballots

Search tracking number

Showing 1 out of 1 ballot.

Tracking number HoVF28p19LEUtK0EkoNaitT7RzhM6AM1BJQLnClc8g

Raw data

Hide



Control pattern



Back to election

Election UUID: JDwmiDBK8QQK6x

Election fingerprint: djB76kleknKDuJNCVzZJ6o5dhStlMnYEIzZTCEIaFe4

How to analyse BeleniosCal ?

# Formal analysis of e-voting systems

Why a formal analysis of an e-voting system?

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—→ Because formal methods can find attacks **before** implementations

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# Formal analysis of e-voting systems

Why a formal analysis of an e-voting system?

→ Because formal methods can find attacks **before** implementations

→ Now a current practice for many protocols (TLS, 5G, ...)

→ Legal requirements in Switzerland to provide **symbolic and cryptographic proofs** of e-voting protocols.

2.14 Proofs of compliance with the cryptographic protocol requirements

2.14.1 A symbolic and a cryptographic proof of compliance must demonstrate that the cryptographic protocol meets the requirements in Numbers 2.1–2.12.

2.14.2 The proofs of compliance must directly refer to the protocol description that forms the basis for system development.

2.14.3 The proofs of compliance relating to basic cryptographic components may be provided according to generally accepted security assumptions and constructions (e.g. «random oracle model», «decisional Diffie-Hellman assumption», «Fiat-Shamir heuristic»).

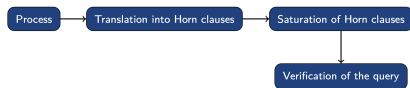


## Two main models for security

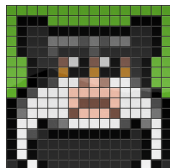
	Formal approach	Computational approach
Messages	<pre>graph TD; Root["{}"] --- Pairs["&lt;, &gt;"]; Root --- Key["k"]; Pairs --- Adversary["A"]; Pairs --- Nonce["N_A"];</pre>	0101000101110101 1101010110101010 0011101011101101
Encryption	terms	bitstrings algorithm
Adversary	idealized	any polynomial algorithm
Guarantees	some attacks missed	stronger
Proof	often automatic	mostly by hand difficult for complex protocols

# Good tools in practice for formal / symbolic models

## ProVerif



## Tamarin



- ▶ fully automatic
  - ▶ axioms, lemmas, and restrictions  
[S&P'22]
  - ▶ framework for verifiability  
[CSF'23]
  - ▶ many voting protocols
    - Swiss Chancellery requirements:  
Swiss Post, CHVote
    - Helios, Belenios, ...
- ▶ semi automatic
  - ▶ exclusive or
  - ▶ voting protocols  
(Belenios, Selene, ...)

# Two major issues for analyzing BeleniosCal

## 1. Addition modulo 2

$$0 + 0 = 0$$

$$0 + 1 = 1$$

$$1 + 0 = 1$$

$$1 + 1 = 0$$

→ state explosion

→ non termination

# Two major issues for analyzing BeleniosCal

## 1. Addition modulo 2

$$0 + 0 = 0$$

$$0 + 1 = 1$$

$$1 + 0 = 1$$

$$1 + 1 = 0$$

→ state explosion

→ non termination

## 2. Probabilistic model

- ▶ Alice checks **either**  $a$  or  $b$  **at random**
- ▶ **Intuition:** An attacker may modify  $k$  votes without been detected with proba  $(\frac{1}{2})^k$ .

# Model for addition modulo - trace properties

*Follows the approach introduced in CCS'22*

## Trace properties (verifiability)

Introduction of two predicates  $\text{isSum}(x, a, b)$  and  $\text{isNotSum}(x, a, b)$

$$\text{isSum}(x, a, b), \quad \text{isSum}(x, a, b') \quad \Rightarrow \quad b = b'$$

$$\text{isSum}(x, a, b), \quad \text{isNotSum}(x, a, b') \quad \Rightarrow \quad b \neq b'$$

- ▶ sound over-approximation
- ▶ another arithmetic operator could be used

# Model for addition modulo - equivalence properties

## Vote secrecy

$$Alice(0) \mid Bob(1) \approx Alice(1) \mid Bob(0)$$

- ▶ over-approximation would be **unsound**

→ Exactly the same tuples  $(x, a, b)$  are created on the left and on the right.

(Lemma)  $isSum(x, a, b) \in fst(tr_b) \Leftrightarrow isSum(x, a, b) \in snd(tr_b)$

→ allow to conclude by hand

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→ allow to conclude by hand

- ▶ privacy relies on the following property:

for all  $x_1, x_2, a_1, a_2$ , there exist  $b_1, b_2, b_3, b_4$  such that

$$x_1 = a_1 + b_1 = a_2 + b_2$$

$$x_2 = a_2 + b_3 = a_1 + b_4$$

→ Encoded in the privacy ProVerif query

# Model for probabilities

- ▶ **to be done!** for a real model with probabilities
- ▶ for the moment, for verifiability, the model assumes that Alice asks for opening **both** ciphertexts.



# To conclude

Many challenges remain! (which is fun 😊)

## Strong demand for Cast as Intended

- ▶ many systems are currently proposed
- ▶ usability (two devices? computation in the head?)
- ▶ trust assumptions?
- ▶ vote secrecy

## Better formal verification

- ▶ decision procedures for larger equational theory classes
- ▶ further improve tools
- ▶ account for probabilities