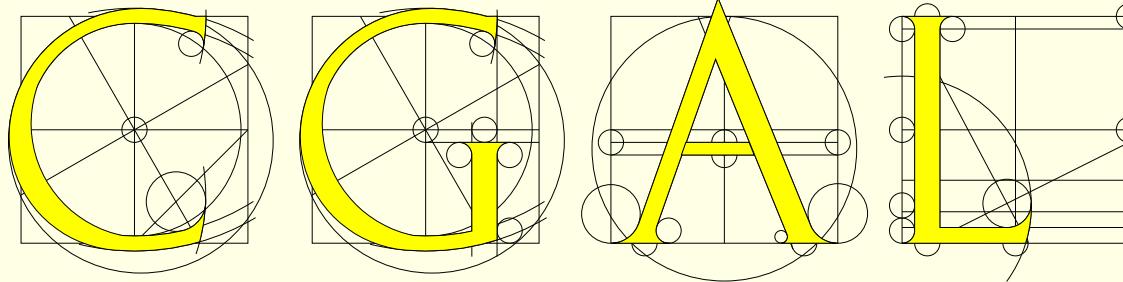


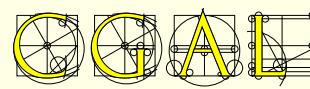
Robustness in



Monique Teillaud



Robustness in



Robustness issues

- Algorithms —> explicit treatment of **degenerate cases**

Symbolic perturbation for 3D dynamic Delaunay triangulations
[Devillers Teillaud SODA'03]

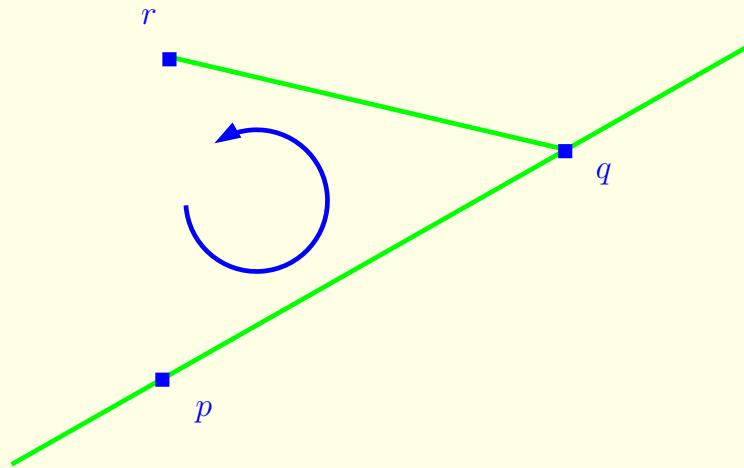
- Kernel and arithmetics —> **Numerical robustness**

Numerical robustness issues

```
typedef CGAL::Cartesian<NT> Kernel;  
NT sqrt2 = sqrt( NT(2) );  
  
Kernel::Point_2 p(0,0), q(sqrt2,sqrt2);  
Kernel::Circle_2 C(p,2);  
  
assert( C.has_on_boundary(q) );
```

OK if NT gives exact sqrt
assertion violation otherwise

Orientation of 2D points

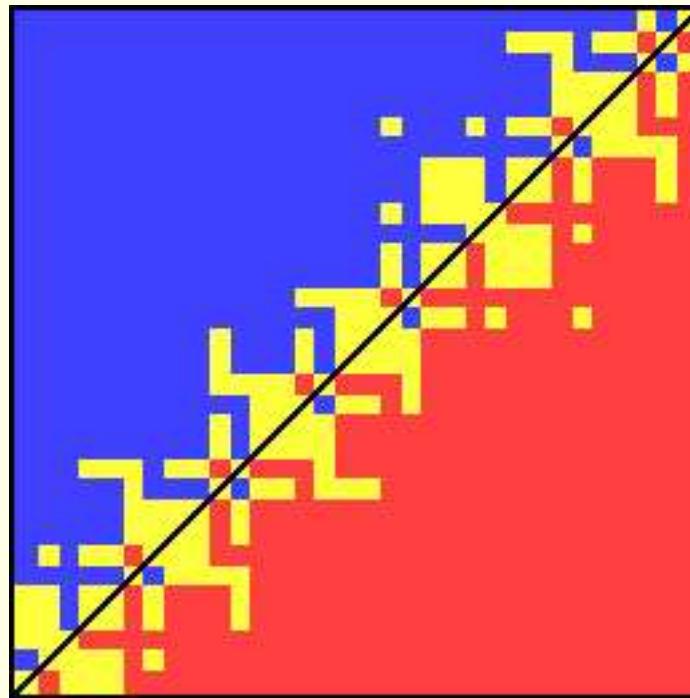


$$\begin{aligned} \text{orientation}(p, q, r) &= \text{sign} \left(\det \begin{bmatrix} p_x & p_y & 1 \\ q_x & q_y & 1 \\ r_x & r_y & 1 \end{bmatrix} \right) \\ &= \text{sign}((q_x - p_x)(r_y - p_y) - (q_y - p_y)(r_x - p_x)) \end{aligned}$$

$p = (0.5 + x.u, 0.5 + y.u)$
 $0 \leq x, y < 256, u = 2^{-53}$
 $q = (12, 12)$
 $r = (24, 24)$

orientation(p, q, r)
evaluated with double

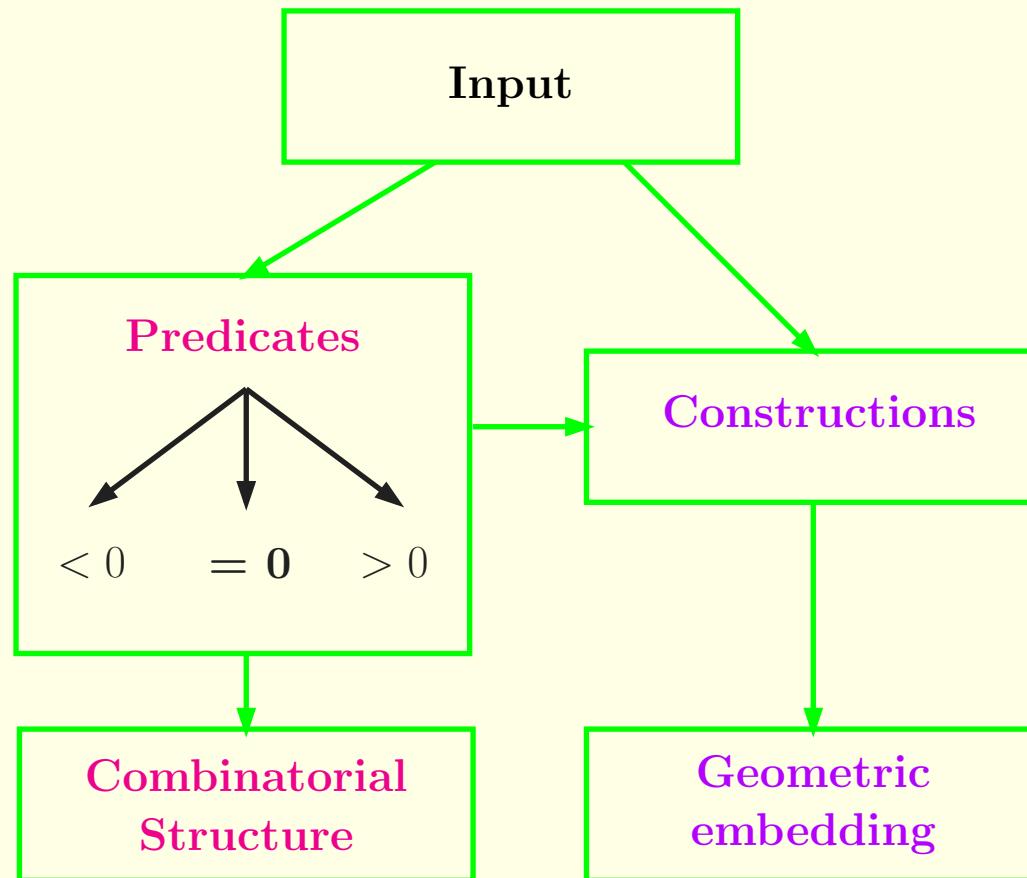
256 × 256 pixel image
■ > 0 , ■ = 0 , ■ < 0



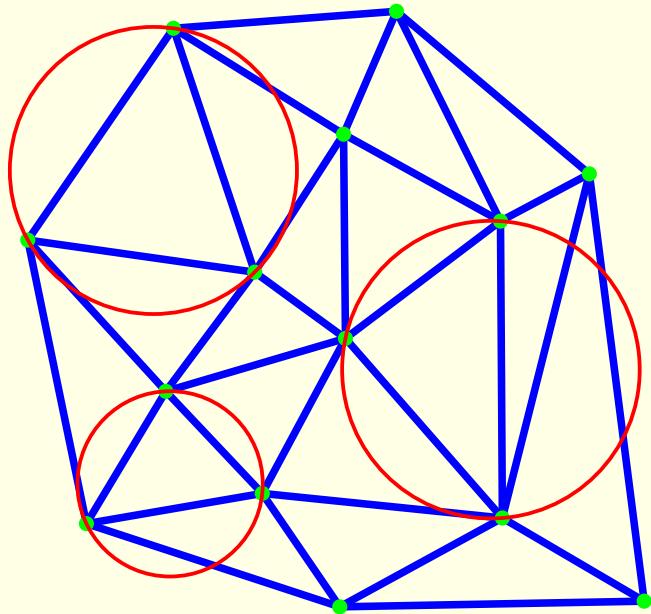
→ **inconsistencies** in predicate evaluations

[Kettner, Mehlhorn, Pion, Schirra, Yap, ESA'04]

Predicates and Constructions

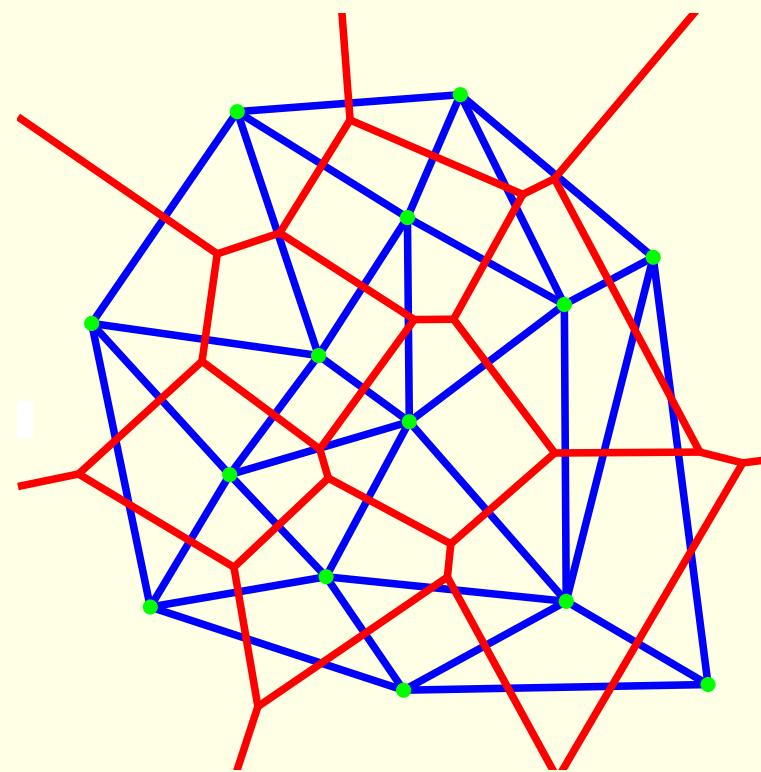


Delaunay triangulation



only **predicates** are used
orientation, in_sphere

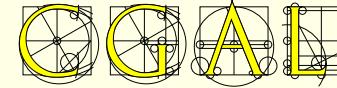
Voronoi diagram



constructions are needed
circumcenter

Arithmetic filters

Numerical Robustness in CGAL



imprecise numerical evaluations
combinatorial result

→ non-robustness

Exact Geometric Computation

\neq
exact arithmetics



Optimize easy cases

Most expected cases: easy, to be optimized first

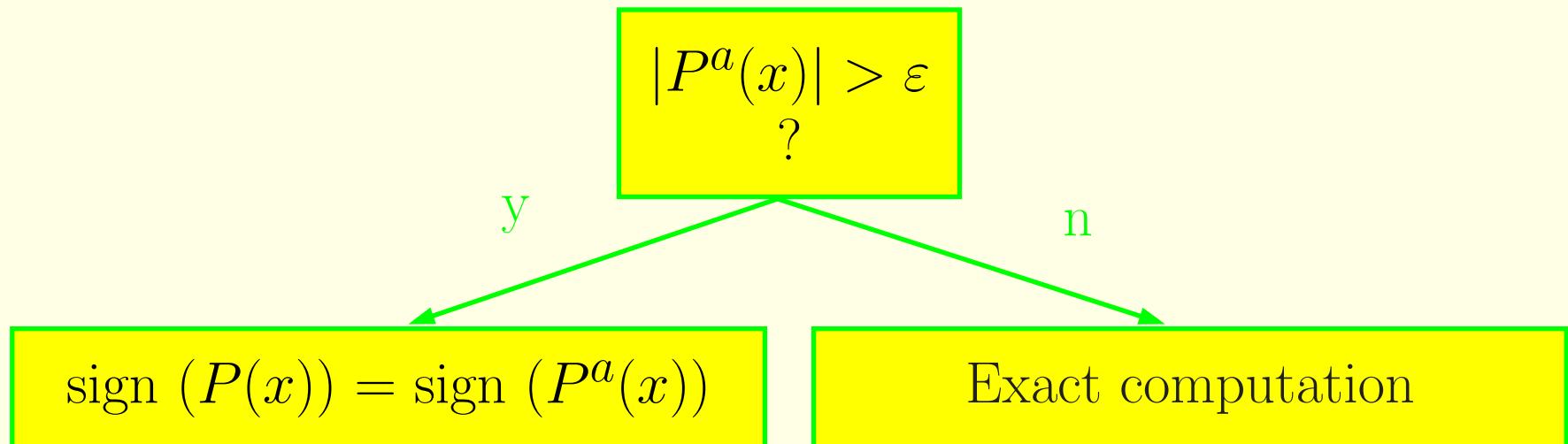
Control rounding errors of floating point computation
⇒ exact computation, expensive, not often used

In good cases, exact geometric computation
but cost \simeq cost of floating point computation.

Filtering Predicates

sign ($P(x)$) ?

Approximate evaluation $P^a(x)$
+ Error ε



Dynamic filters: interval arithmetic

Floating point operation replaced by
operations on **intervals** of floating point values $[\underline{x}; \bar{x}]$
encoding rounding errors.

Inclusion property:
at each operation, the interval contains the exact value of X .

Operations on intervals

Rounding modes IEEE 754

Addition / subtraction

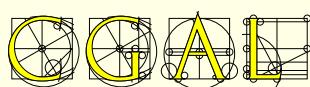
$$X + Y \longrightarrow [\underline{x+y}; \bar{x}+\bar{y}]$$

$$X - Y \longrightarrow [\underline{x-\bar{y}}; \bar{x}-\underline{y}]$$

Optimization:

$$X + Y \longrightarrow [-((-x)\bar{-}y); \bar{x}+\bar{y}]$$

(fewer changes of rounding modes)



Operations on intervals

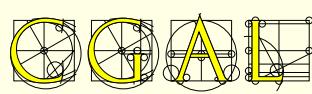
Multiplication :

$$X \times Y \longrightarrow [\min(\underline{x} \times \underline{y}, \underline{x} \times \bar{y}, \bar{x} \times \underline{y}, \bar{x} \times \bar{y}); \max(\underline{x} \times \bar{y}, \underline{x} \times \bar{y}, \bar{x} \times \underline{y}, \bar{x} \times \bar{y})]$$

In practice: comparisions for different cases before performing multiplications.

Division : similar

Handling of division by 0.



Comparisons

Inclusion property

if

$$[\underline{x}; \bar{x}] \cap [\underline{y}; \bar{y}] = \emptyset$$

then

we can decide whether $X < Y$ or $X > Y$

else

we cannot decide.

→ Filter failure

Static filters

Static analysis of error propagation on evaluation of a polynomial expression, assuming **bounds on the input data**.

x being a positive floating point value,
and y the smallest floating point value greater than x

$$\text{ulp}(x) = y - x$$

(Unit in the Last Place).

Remark 1 : $\text{ulp}(x)$ is a power of 2 (or ∞).

Remark 2 : In normal cases : $\text{ulp}(x) \simeq x \cdot 2^{-53}$



x real, x value computed in double,
 e_x and b_x doubles such that

$$\begin{cases} e_x \geq |x - \text{x}| \\ b_x \geq |\text{x}| \end{cases}$$

Initially, value rounded to closest
(if values cannot be represented by a double)

$$\begin{cases} b_x = |\text{x}| \\ e_x = \frac{1}{2}\text{ulp}(\text{x}) \end{cases}$$

For $+, -, \times, \div, \sqrt{}$, rounding error on result r smaller than
- $\frac{1}{2}\text{ulp}(r)$ for rounding to nearest mode
- $\text{ulp}(r)$ otherwise.



Addition and subtraction

Propagation of error on an addition $z = x + y$:

$$\begin{cases} b_z = b_x + b_y \\ e_z = e_x \mp e_y \mp \frac{1}{2}\text{ulp}(z) \end{cases}$$

Indeed:

$$\begin{aligned} |z - z| &= |\underbrace{(z - (x + y))}_{=0} + \underbrace{((x + y) - (x + y))}_{\leq e_x + e_y} + \underbrace{((x + y) - z)}_{\leq \frac{1}{2}\text{ulp}(z)}| \\ &\leq e_x \mp e_y \mp \frac{1}{2}\text{ulp}(z) \end{aligned}$$



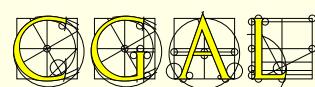
Multiplication

Propagation of error on a multiplication $z = x \times y$:

$$\begin{cases} b_z = b_x \times b_y \\ e_z = e_x \overline{\times} e_y \mp e_y \overline{\times} |x| \mp e_x \overline{\times} |y| \mp \frac{1}{2}\text{ulp}(z) \end{cases}$$

Indeed:

$$\begin{aligned} |z - z| &= |\underbrace{(z - (x \times y))}_{=0} + \underbrace{((x \times y) - (x \times y))}_{=(x-x)(y-y)-(x-x) \times y - (y-y) \times x} + \underbrace{((x \times y) - z)}_{\leq \frac{1}{2}\text{ulp}(z)}| \\ &\leq e_x \overline{\times} e_y \mp e_x \overline{\times} y \mp e_y \overline{\times} x \mp \frac{1}{2}\text{ulp}(z) \end{aligned}$$



Application: *orientation* predicate

Approximate non guaranteed version

```
int orientation(double px, double py,
                double qx, double qy,
                double rx, double ry)
{
    double pqx = qx - px, pqy = qy - py;
    double prx = rx - px, pry = ry - py;

    double det = pqx * pry - pqy * prx;

    if (det > 0) return 1;
    if (det < 0) return -1;
    return 0;
}
```

Application: *orientation* predicate

Code with static filtering (for entries **bounded by 1**):

```
int filtered_orientation(double px, double py,
                         double qx, double qy,
                         double rx, double ry)
{
    double pqx = qx - px, pqy = qy - py;
    double prx = rx - px, pry = ry - py;

    double det = pqx * pry - pqy * prx;

    const double E = 1.33292e-15;

    if (det > E) return 1;
    if (det < -E) return -1;

    ... // can't decide => call the exact version
}
```

Variants - Ex : compute the bound at running time

```
int filtered_orientation(double px, double py,
                        double qx, double qy,
                        double rx, double ry)
{
    double b = max_abs(px, py, qx, qy, rx, ry);

    double pqx = qx - px, pqy = qy - py;
    double prx = rx - px, pry = ry - py;

    double det = pqx * pry - pqy * prx;

    const double E = 1.33292e-15;

    if (det > E*b*b) return 1;
    if (det < -E*b*b) return -1;

    ... // can't decide => call the exact version
}
```

Probability of filter failures

Theoretical study: [Devillers-Preparata-99]

Input data **uniformly distributed** in a unit square/cube

static filtering

orientation 2D 10^{-15}

orientation 3D 5.10^{-14}

in_circle 2D 10^{-11}

in_sphere 3D 7.10^{-10}

More degenerate cases

	Dynamic	Semi-static
Random	0	870
$\varepsilon = 2^{-5}$	0	1942
$\varepsilon = 2^{-10}$	0	662
$\varepsilon = 2^{-15}$	0	8833
$\varepsilon = 2^{-20}$	0	132153
$\varepsilon = 2^{-25}$	10	192011
$\varepsilon = 2^{-30}$	19536	308522
Grid	49756	299505

Number of filter failures for dynamic and static filters during the computation of a Delaunay triangulation on 10^5 points).

Data on an integer grid with precision of 30 bits, with relative perturbation.

Comparaison : dynamic vs static filters

static filtering

- **fails more often** than more precise interval arithmetic filtering
- **faster**
- **harder to write**: needs analysis of each predicate.

Fastest method: **Cascading filters**

Implementation in

Arithmetic tools

- **Multiprecision integers**

Exact evaluation of signs / values of polynomial expressions with integer coefficients

CGAL::MP_Float, GMP::mpz_t, LEDA::integer, ...

- **Multiprecision floats**

idem, with float coefficients ($n2^m, n, m \in \mathbb{Z}$)

CGAL::MP_Float, GMP::mpf_t, LEDA::bigfloat, ...

- **Multiprecision rationals**

Exact evaluation of signs / values of rational expressions

CGAL::Quotient< · >, GMP::mpq_t, LEDA::rational, ...

- **Algebraic numbers**

Exact comparison of roots of polynomials

LEDA::real, Core::Expr (work in progress in CGAL)



Dynamic filtering

Number types: **CGAL::Interval_nt**, **MPFR/MPFI**, **boost::interval**

CGAL::Filtered_kernel < K > kernel wrapper

[Pion]

Replaces predicates of **K** by filtered and exact predicates.
(exact predicates computed with MP_Float)

Static + Dynamic filtering in CGAL 3.1

→ more generic generator also available for user's predicates

CGAL::Filtered_predicate

Filtering Constructions

Number type **CGAL::Lazy_exact_nt < Exact_NT >**

[Pion]

Delays exact evaluation with **Exact_NT**:

- stores a **DAG** of the expression
- computes first an approximation with **Interval_nt**
- allows to control the relative precision of `to_double`



CGAL::Lazy_kernel in CGAL 3.2

Predefined kernels

Exact_predicates_exact_constructions_kernel

Filtered_kernel< Cartesian< Lazy_exact_nt< Quotient< MP_Float >>>

Exact_predicates_exact_constructions_kernel_with_sqrt

Filtered_kernel< Cartesian< Core::Expr >>

Exact_predicates_inexact_constructions_kernel

Filtered_kernel< Cartesian< double >>



Efficiency

3D Delaunay triangulation

CGAL-3.1-I-124

1.000.000 random points

Simple_Cartesian< double >

Simple_Cartesian< MP_Float >

Filtered_kernel (dynamic filtering)

Filtered_kernel (static + dynamic filtering)

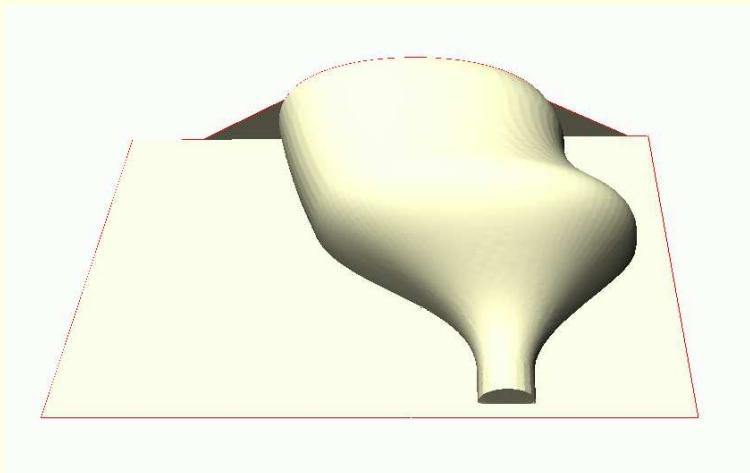
Pentium-M 1.7 GHz, 1GB
g++ 3.3.2, -O2 -DNDEBUG

48.1 sec

2980.2 sec

232.1 sec

58.4 sec



49.787 points (Dassault Systèmes)
double loop !
exact and filtered < 8 sec

Work in progress

- **Automatic generation of code** from a generic version
- filtering of **constructions**
- **Rounding** of constructions
- **Curved objects** (algebraic methods)